



ITMO UNIVERSITY

How to Win Coding Competitions: Secrets of Champions

Week 3: Sorting and Search Algorithms **Lecture 6: Mergesort**

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Saint Petersburg 2016

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- ▶ This algorithm needs **extra scratch memory**


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procedure MERGESORT( $A, W, \preceq, s, t$ )  
  if  $s + 1 = t$  then return end if  
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   $i \leftarrow s, j \leftarrow m, k \leftarrow s$   
  while  $i < m$  or  $j < t$  do  
    if  $j = t$  or ( $i < m$  and  $A[i] \preceq A[j]$ ) then  
       $W[k] \leftarrow A[i], i \leftarrow i + 1, k \leftarrow k + 1$   
    else  
       $W[k] \leftarrow A[j], j \leftarrow j + 1, k \leftarrow k + 1$   
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  end while  
  
  for  $i \in [s; t)$  do  $A[i] \leftarrow W[i]$  end for  
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 - ▶ In-place merge is possible, but is either slower or very complicated

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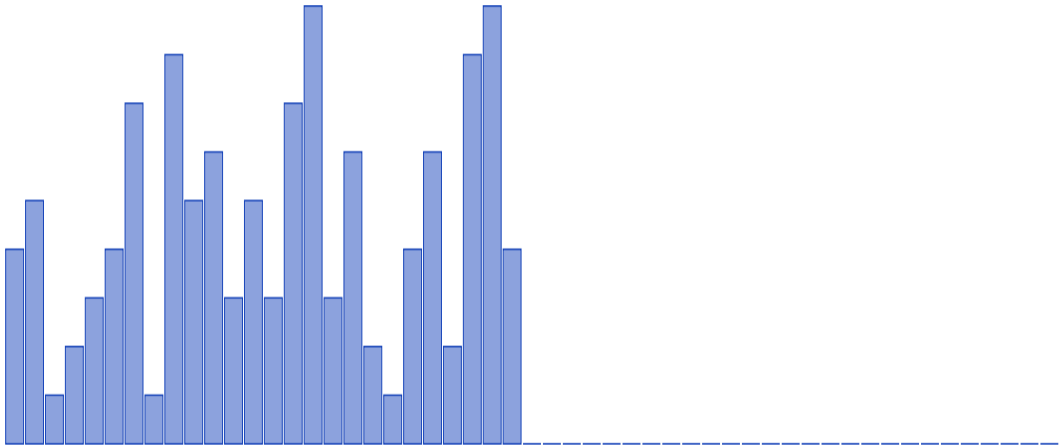
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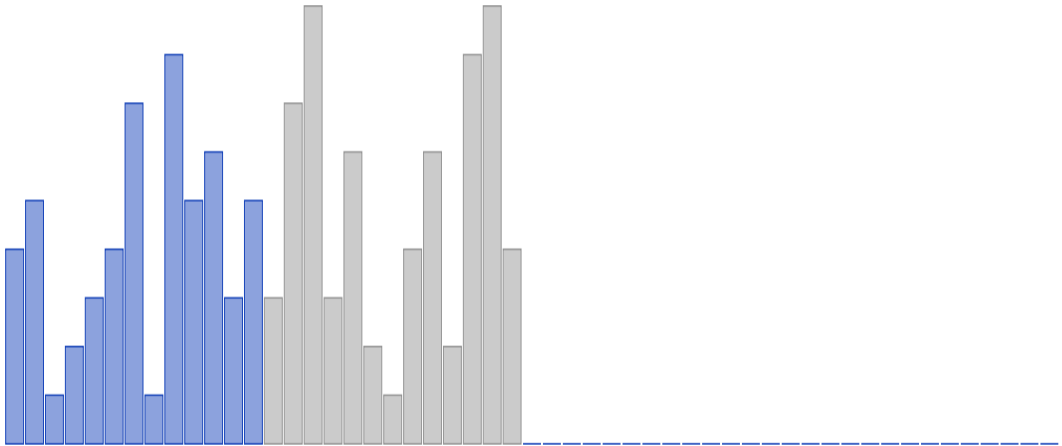
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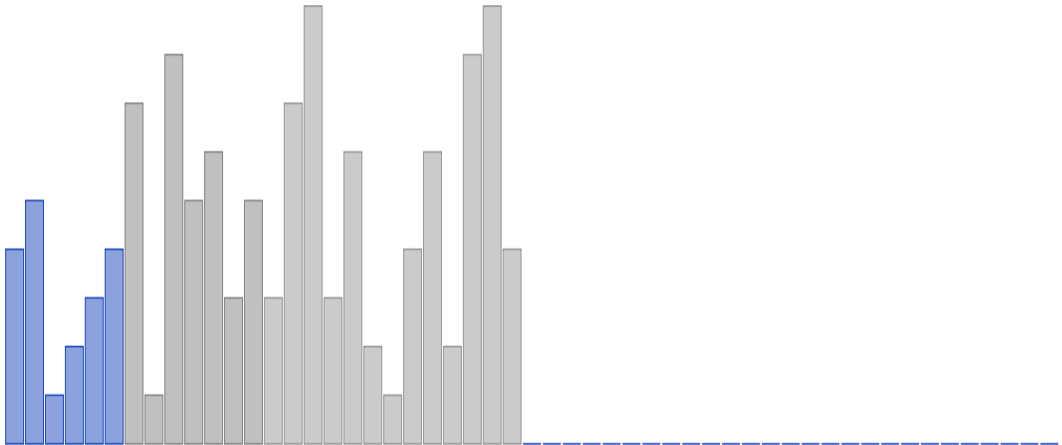
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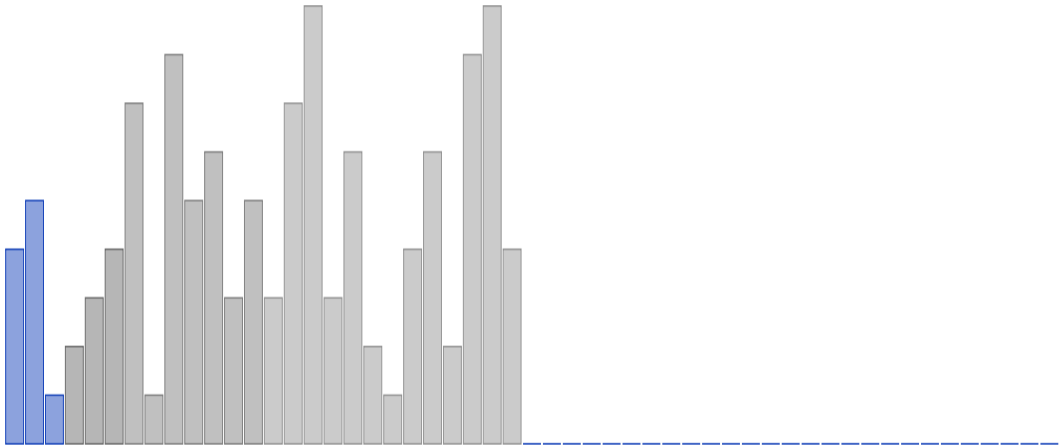
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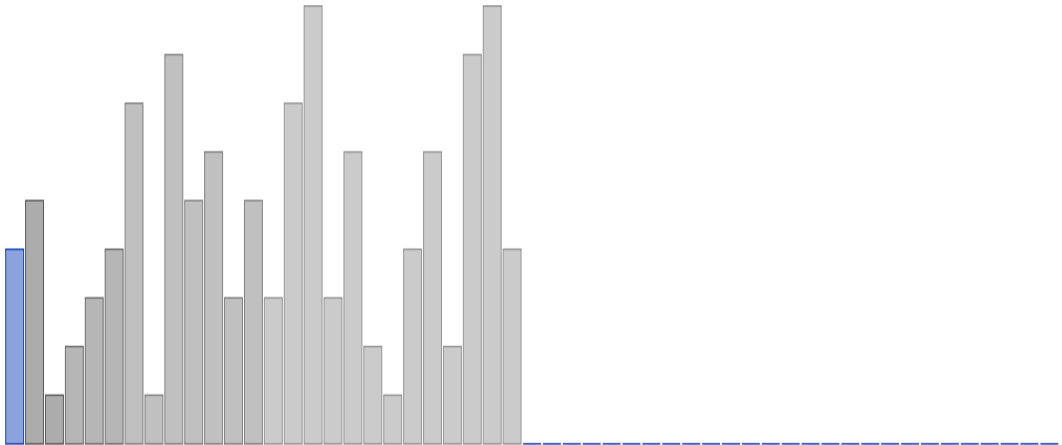
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- ▶ Merge runs in $\Theta(t - s)$
 - ▶ Every scratch element is written exactly once

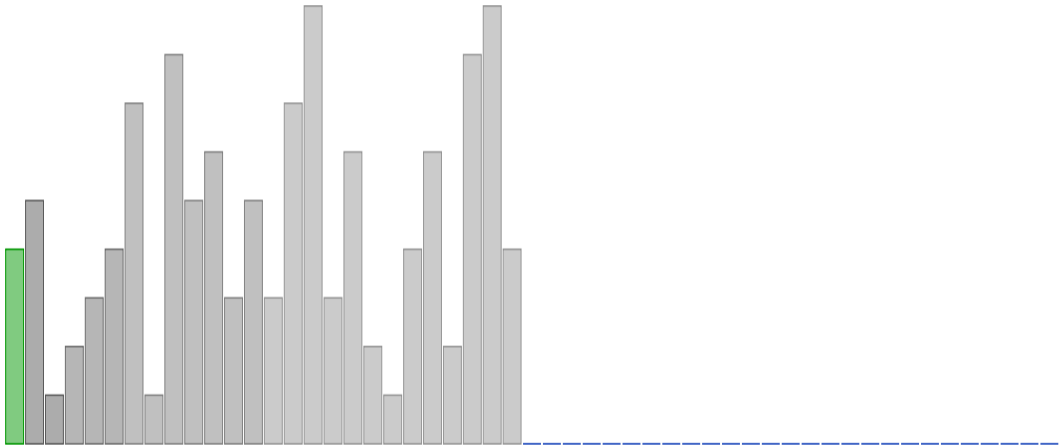


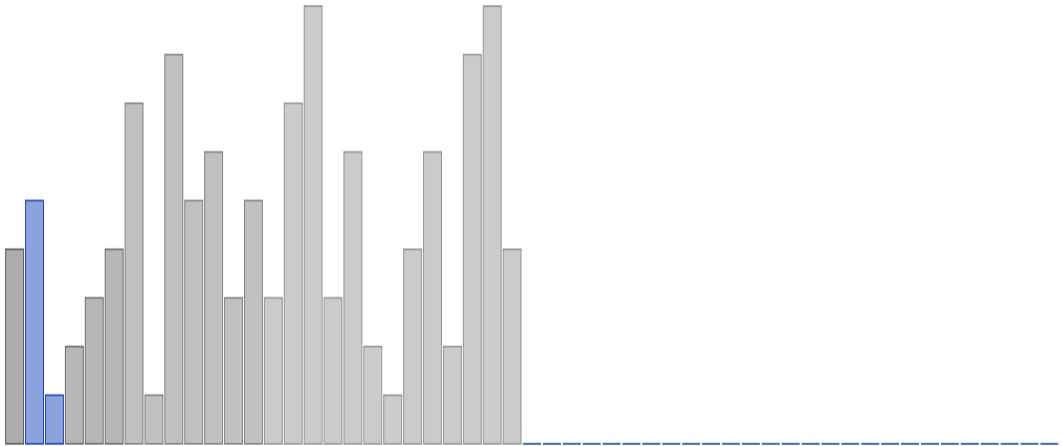


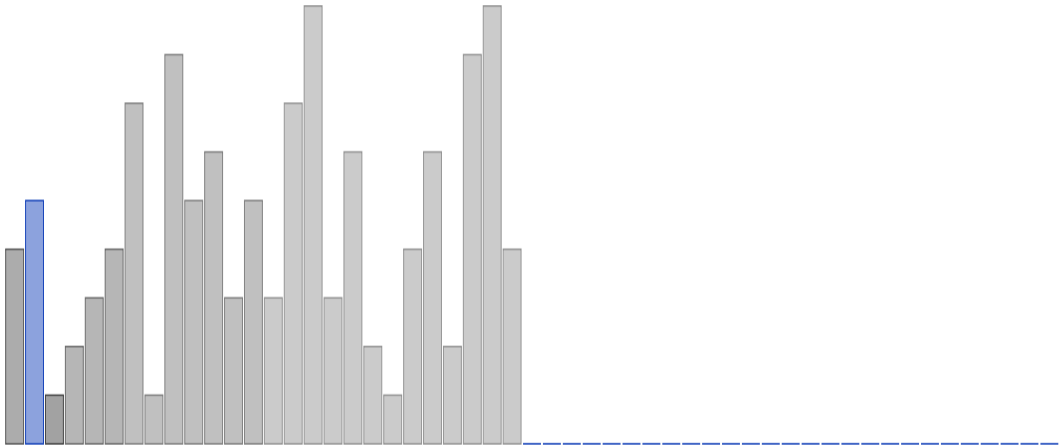


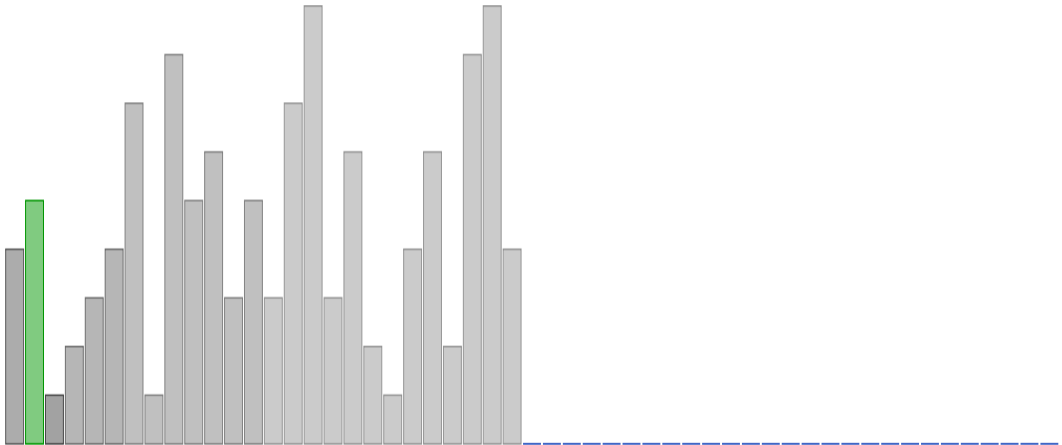


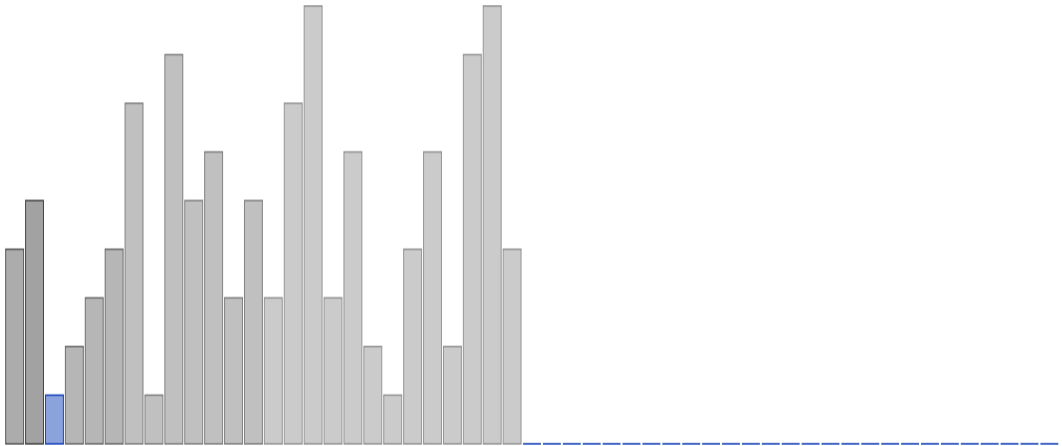


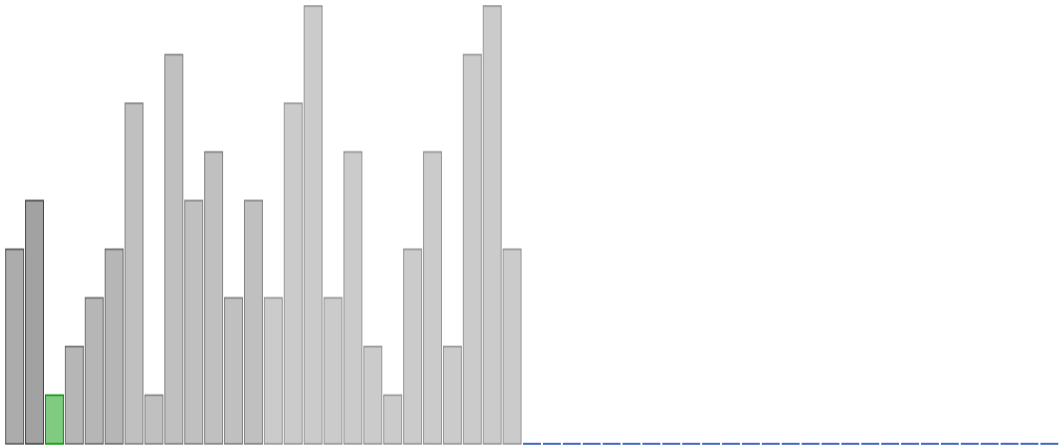


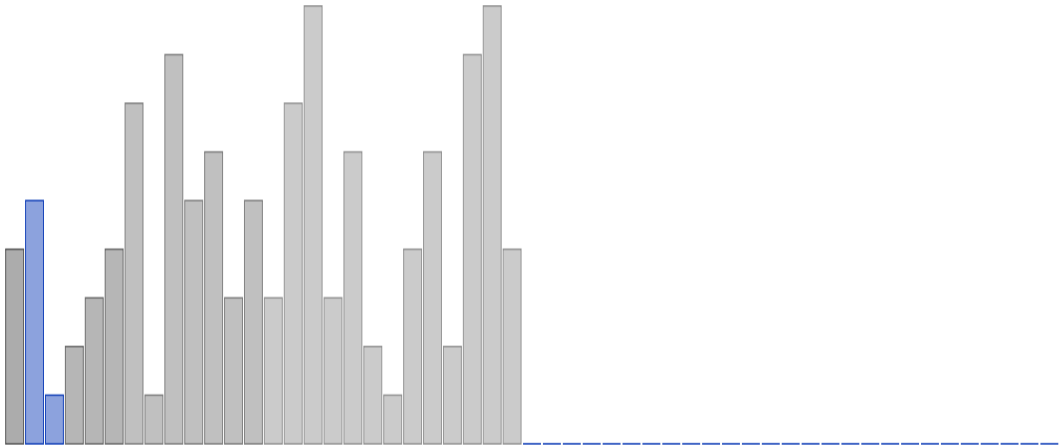


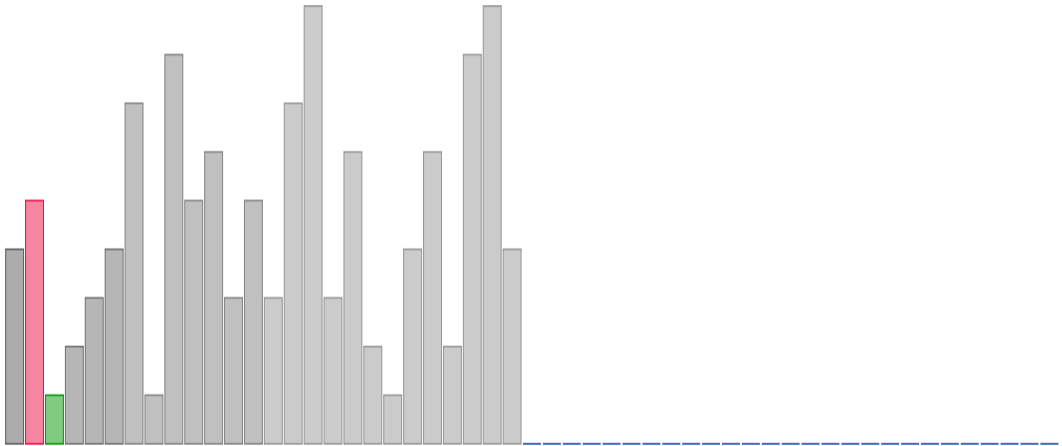


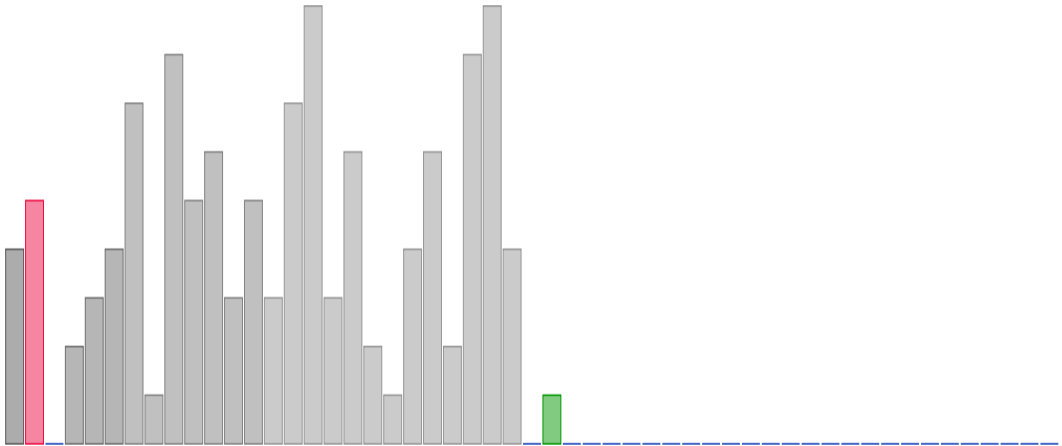


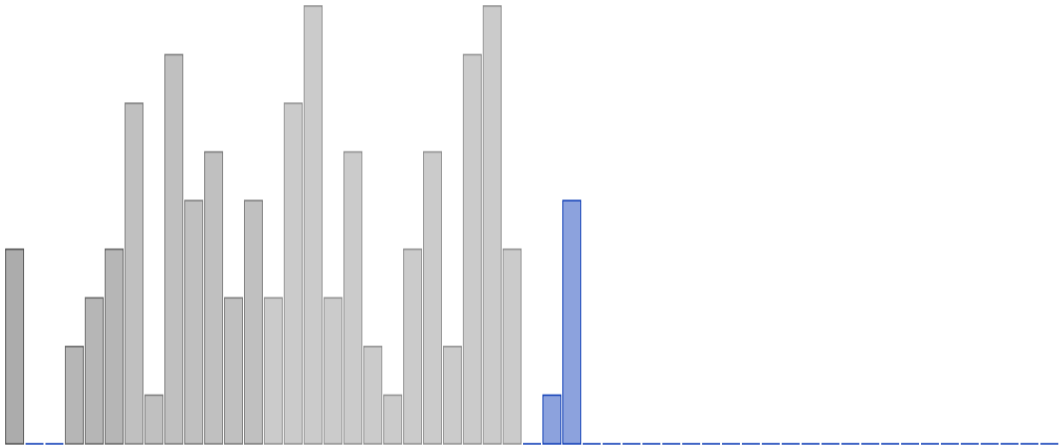


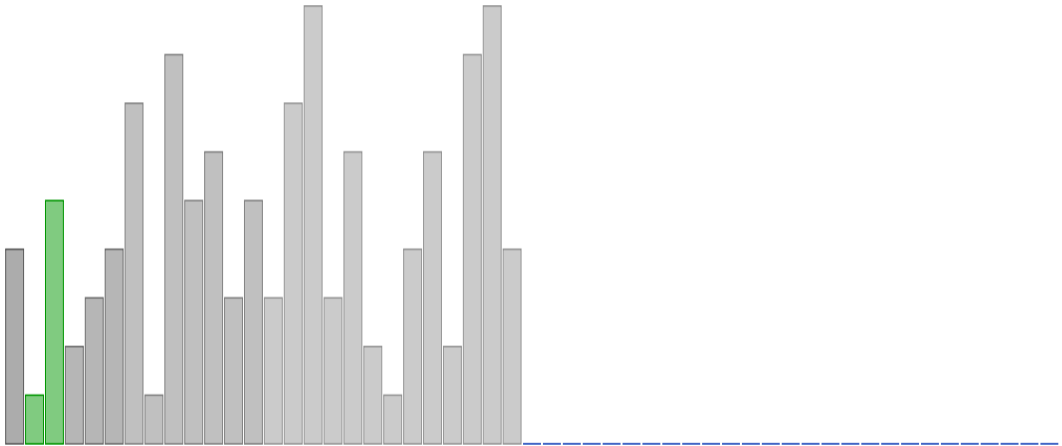


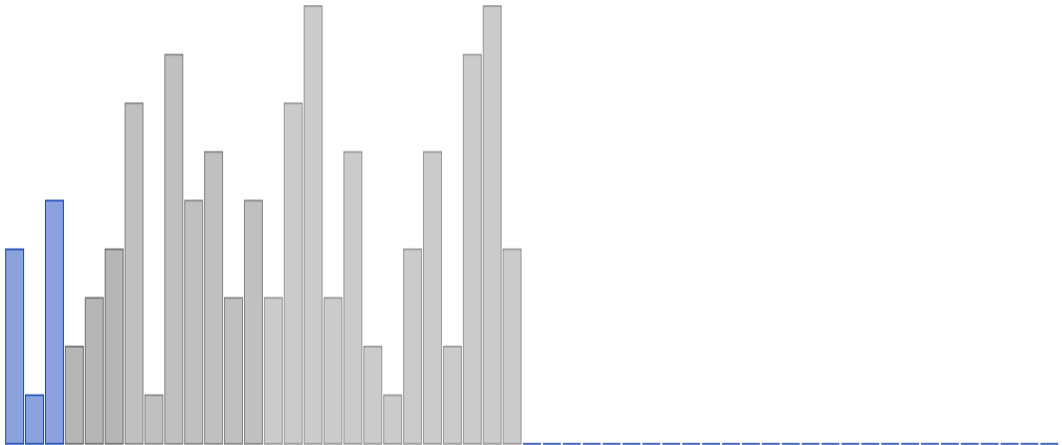


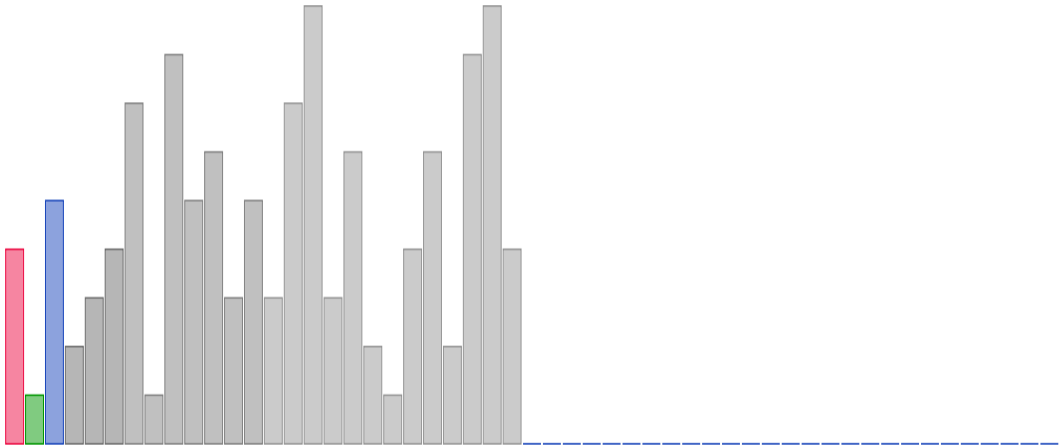


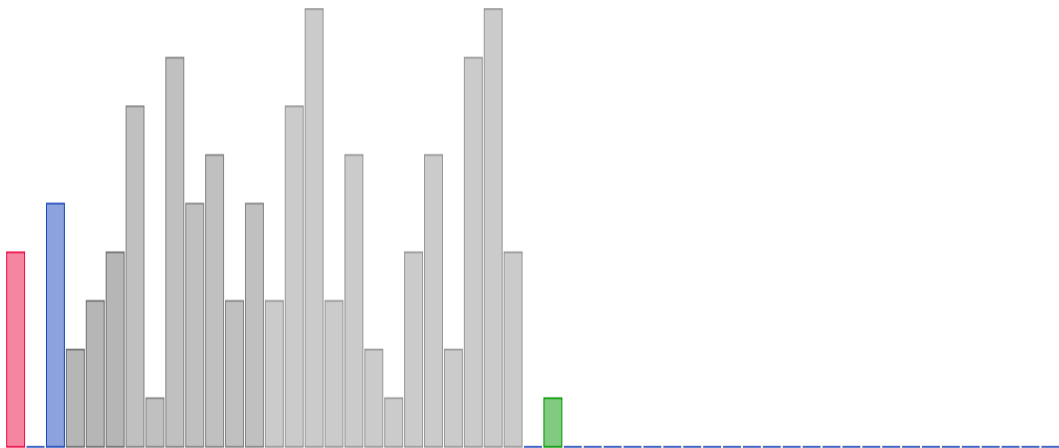


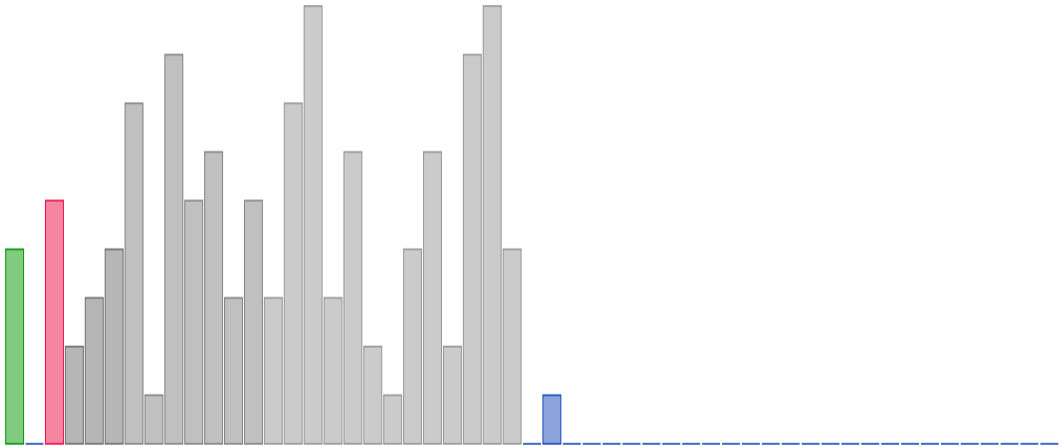


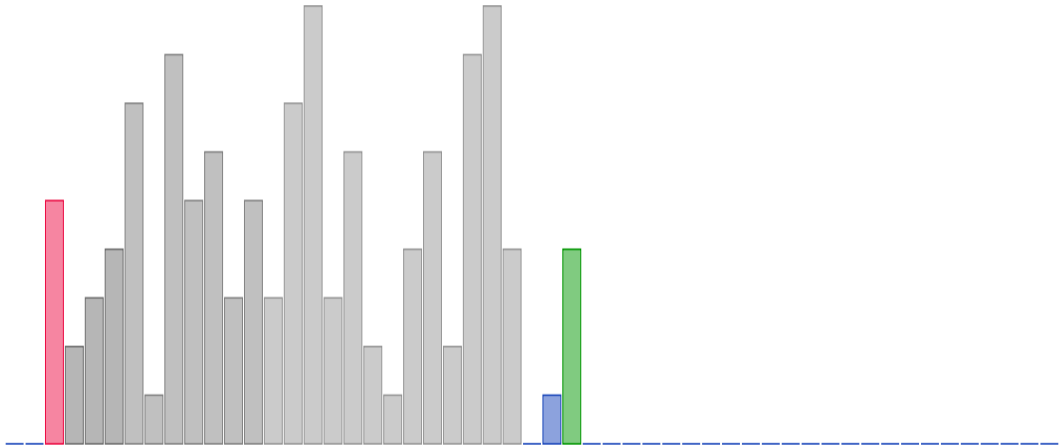


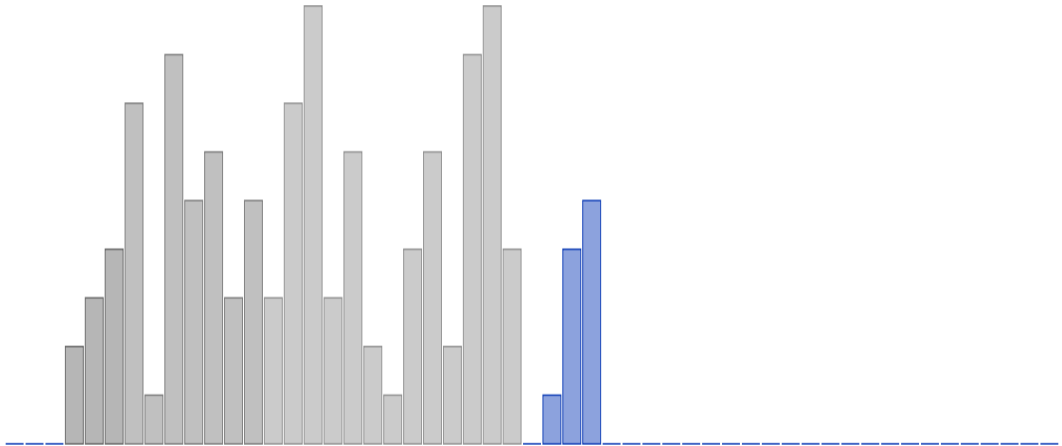


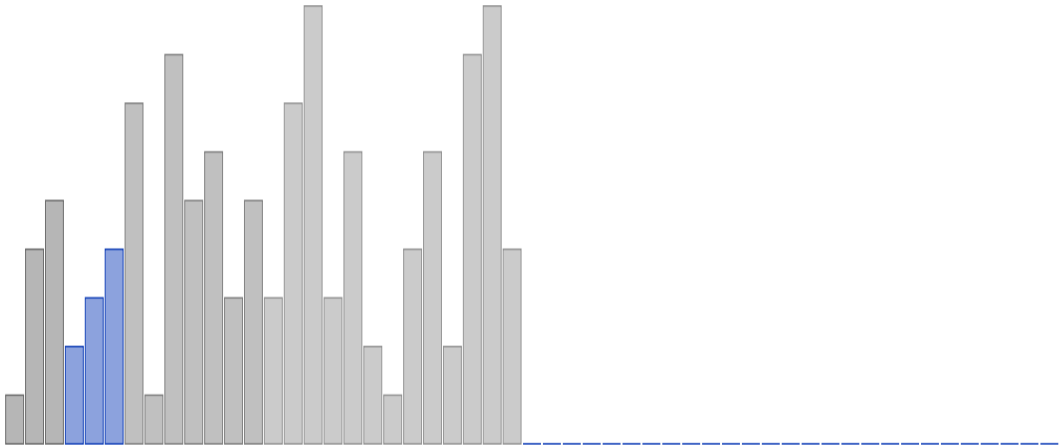


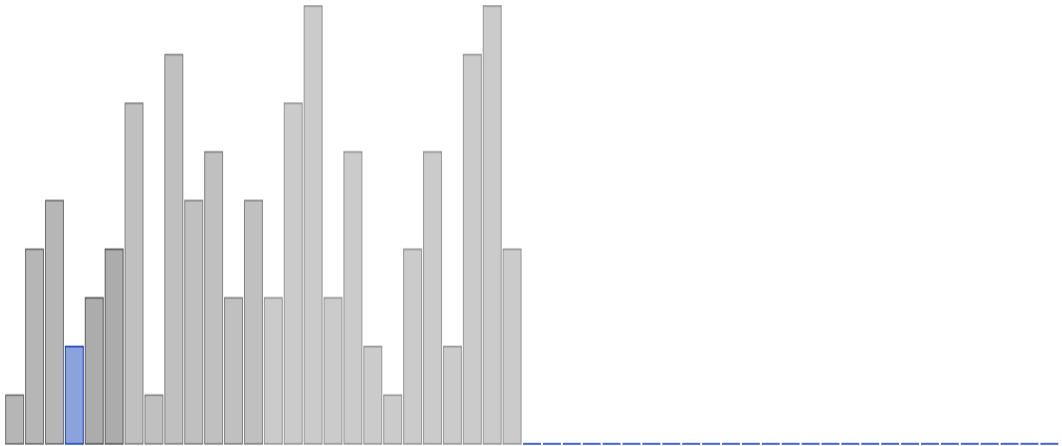


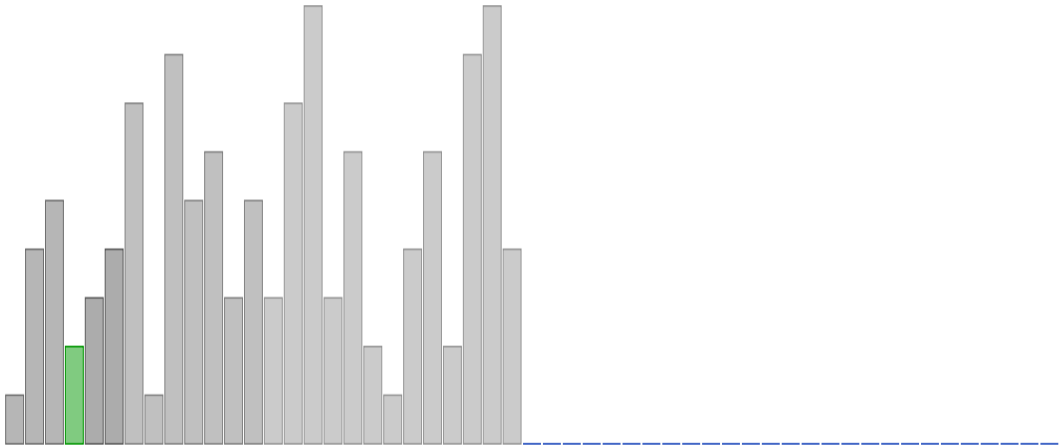


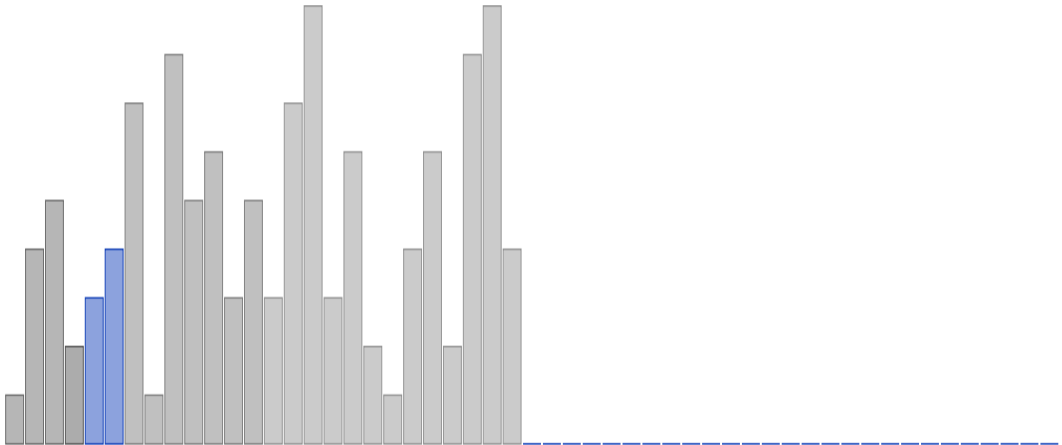


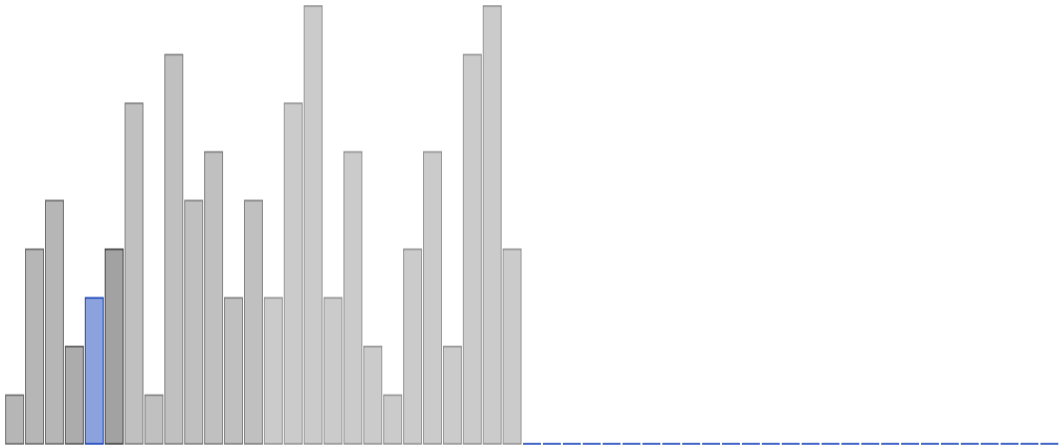


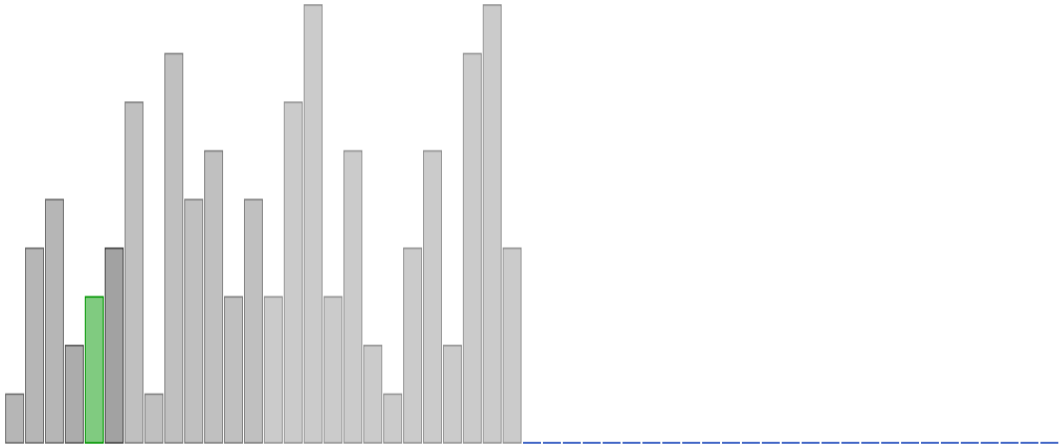


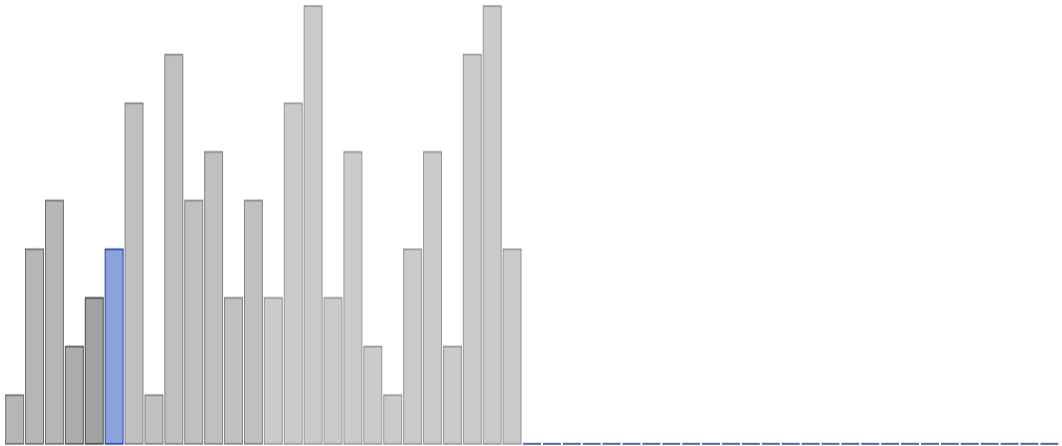


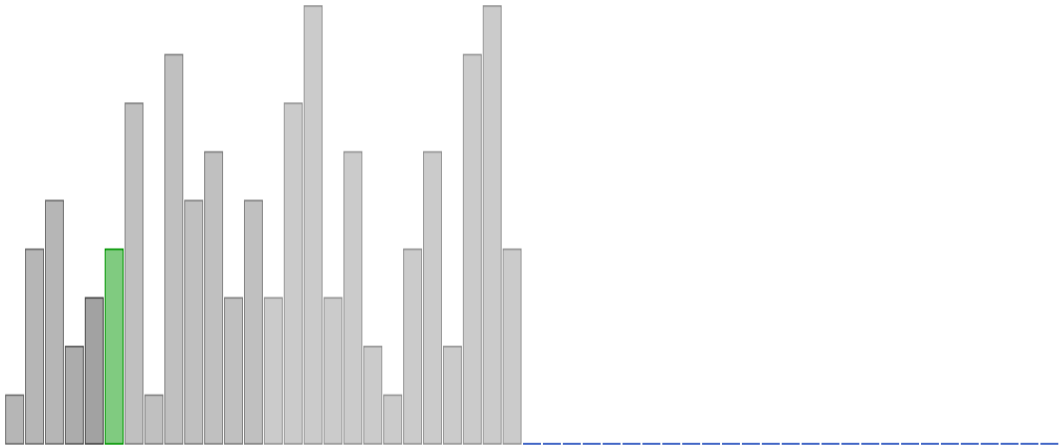


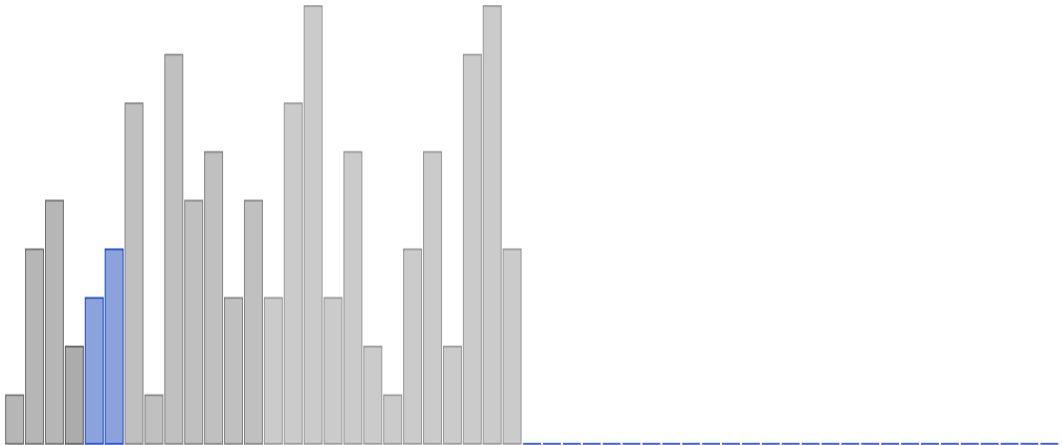


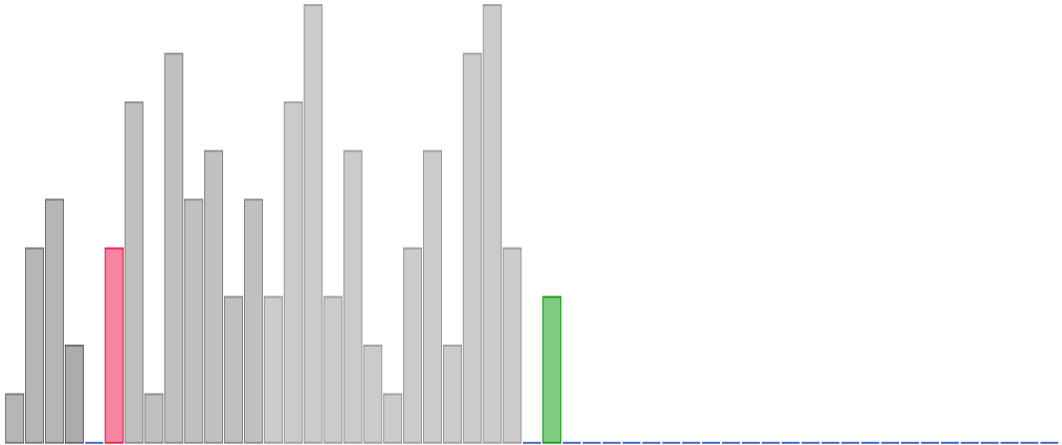


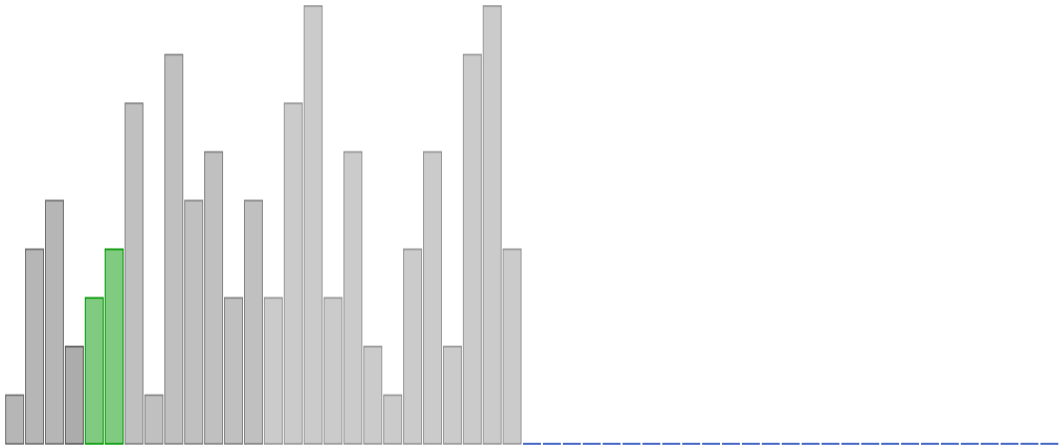


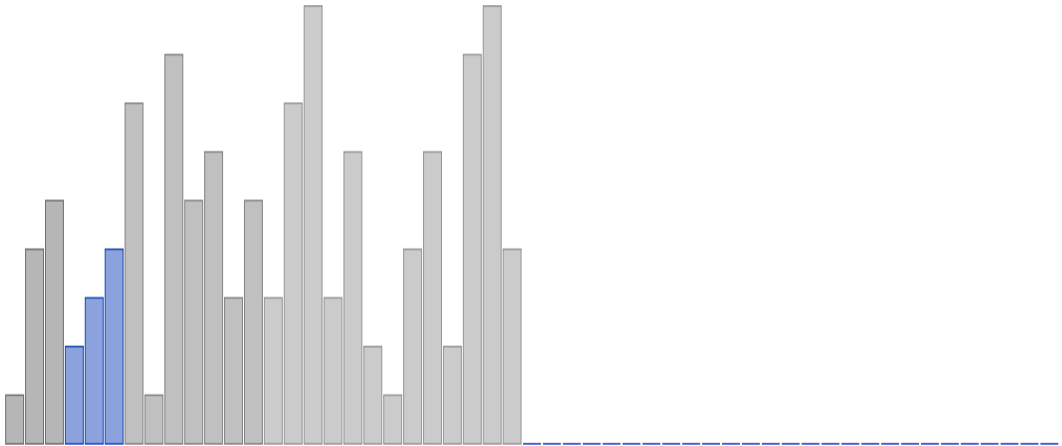


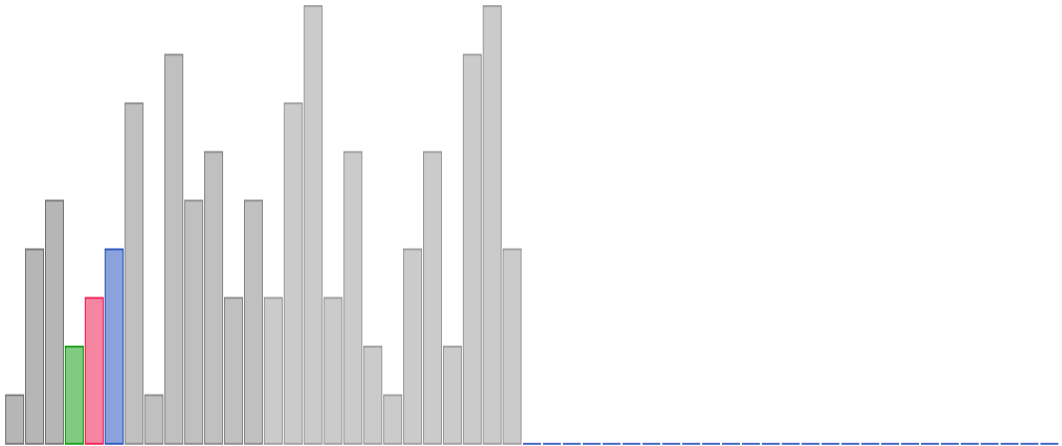


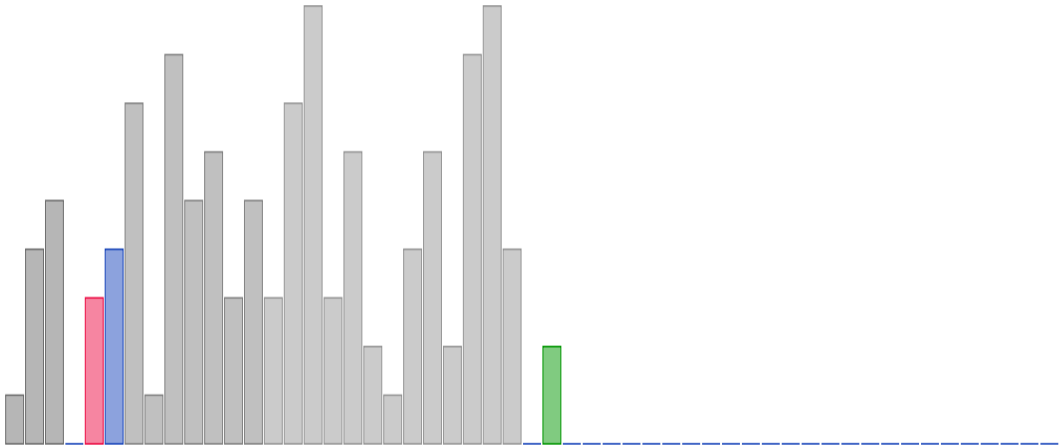


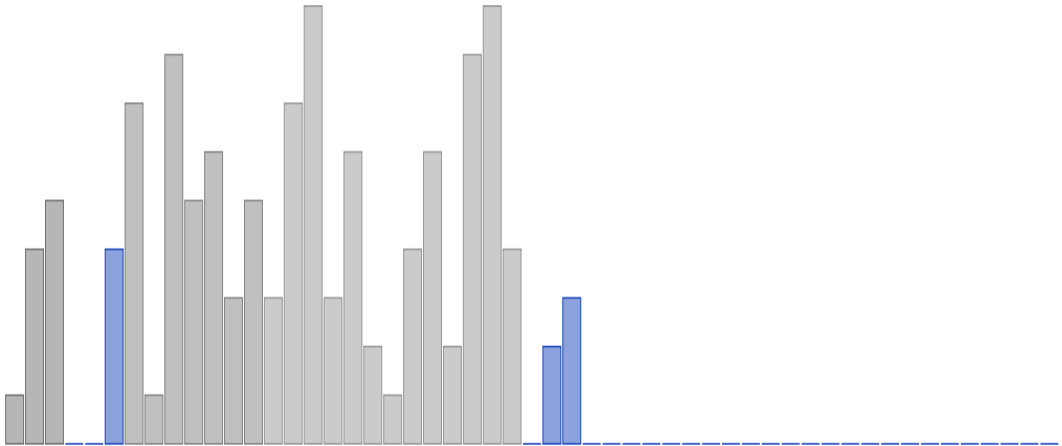


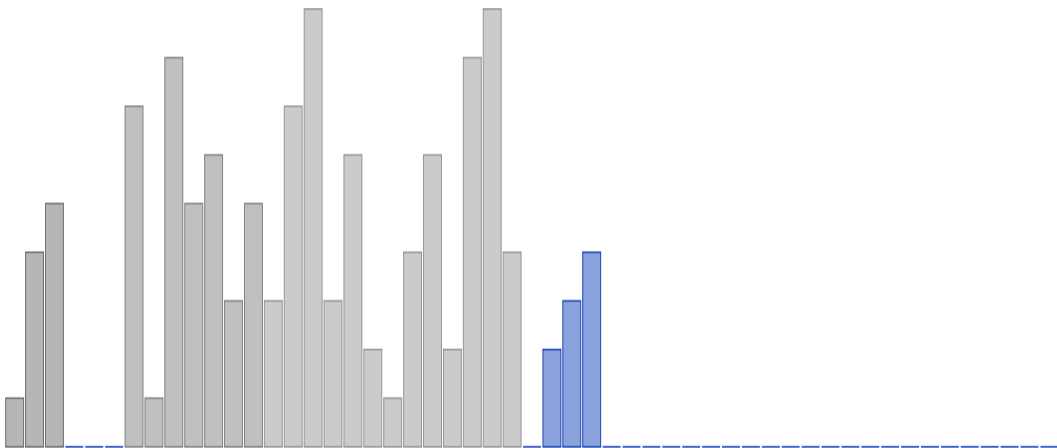


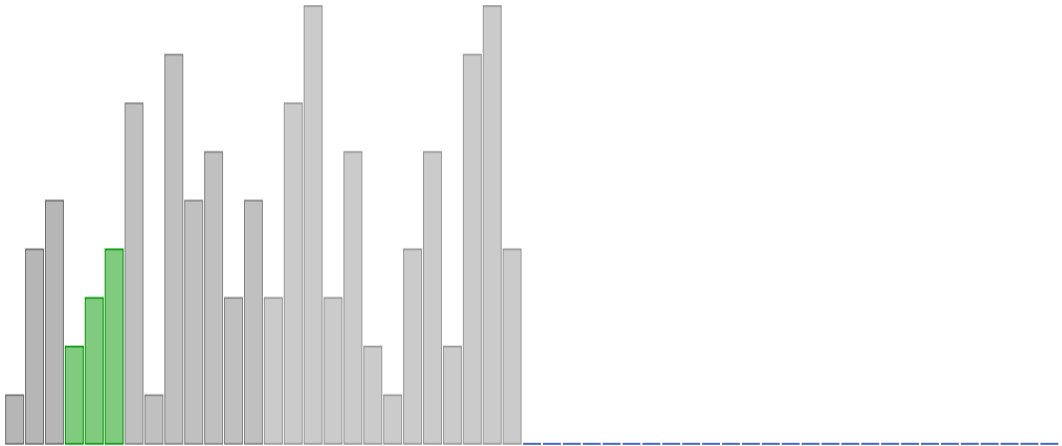


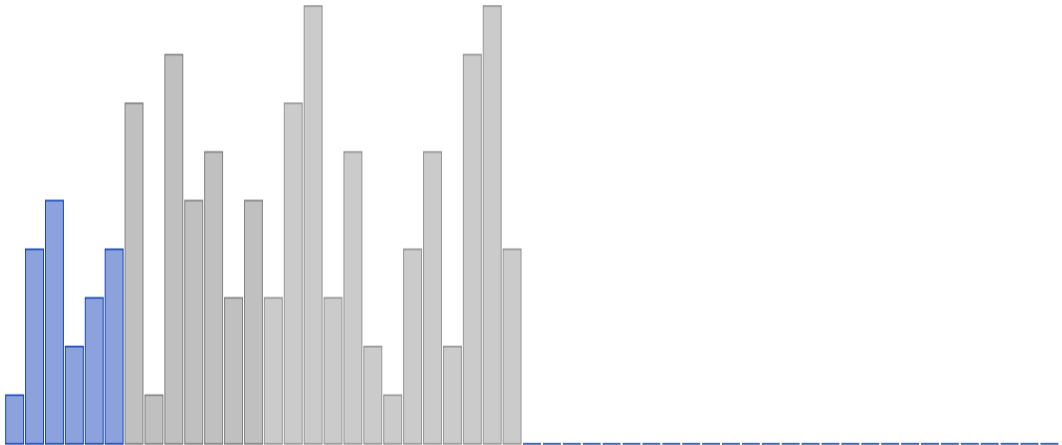


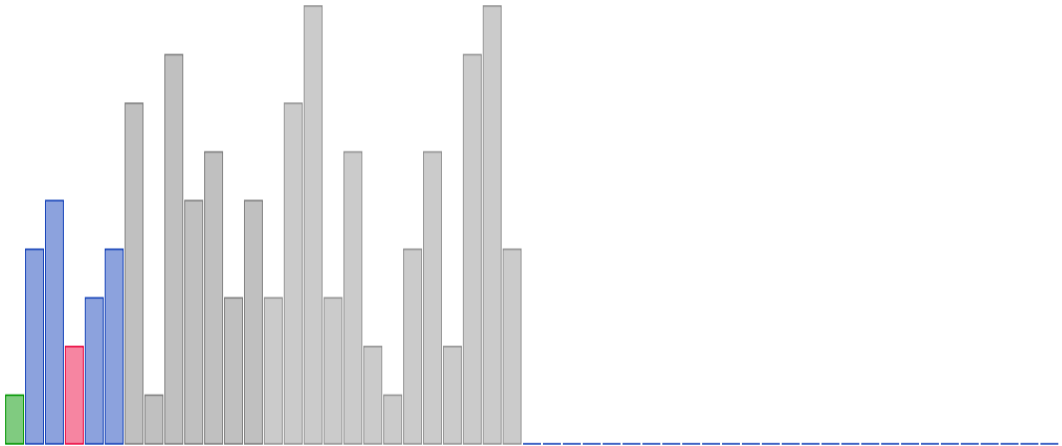


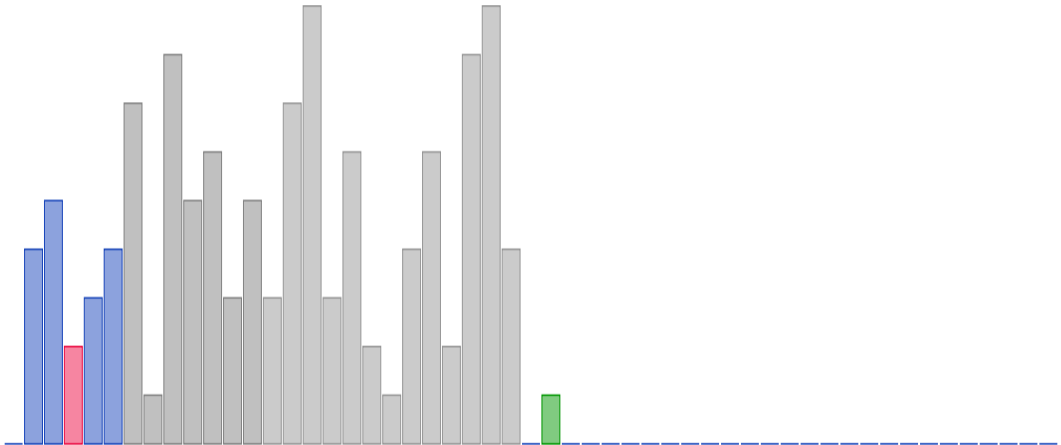


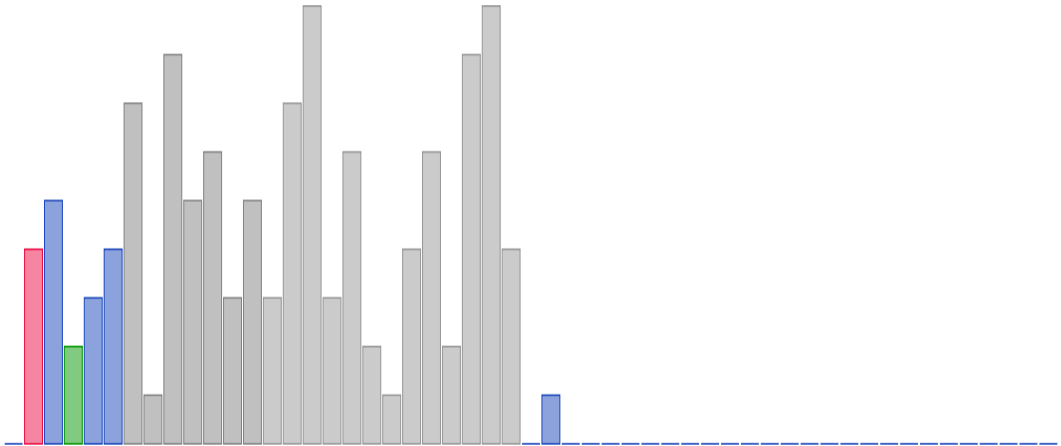


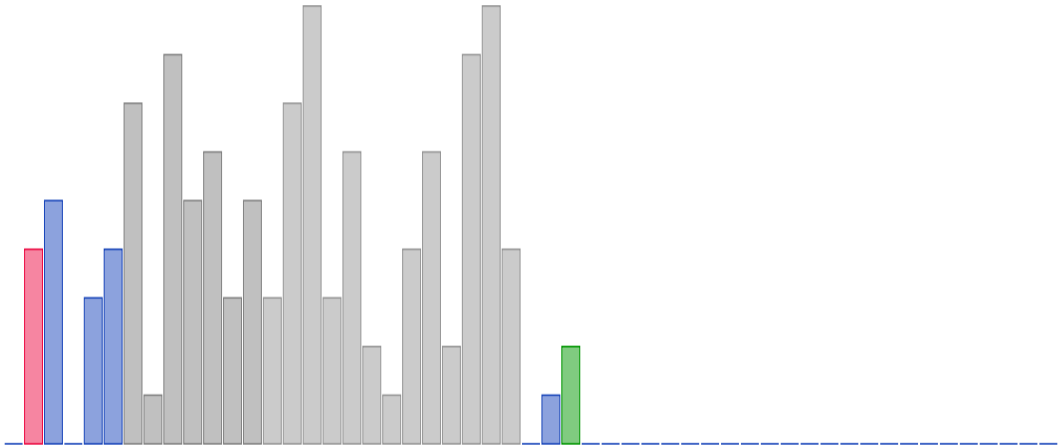


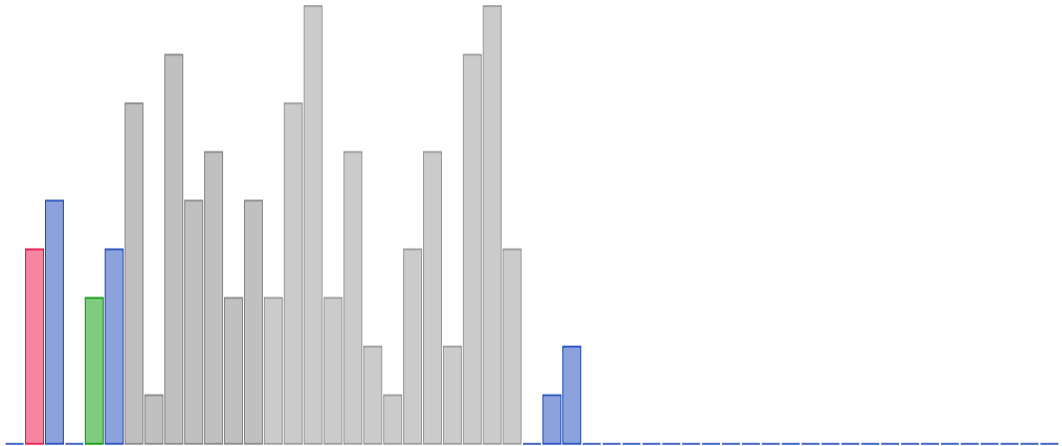


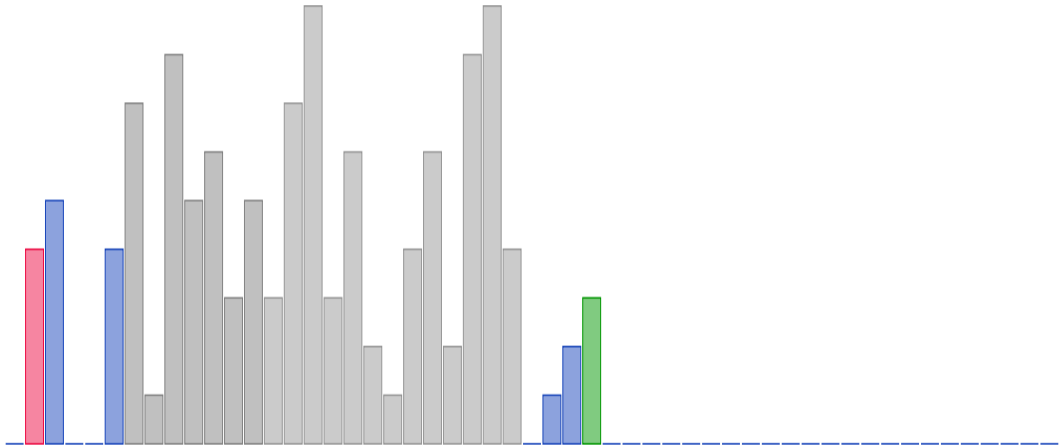


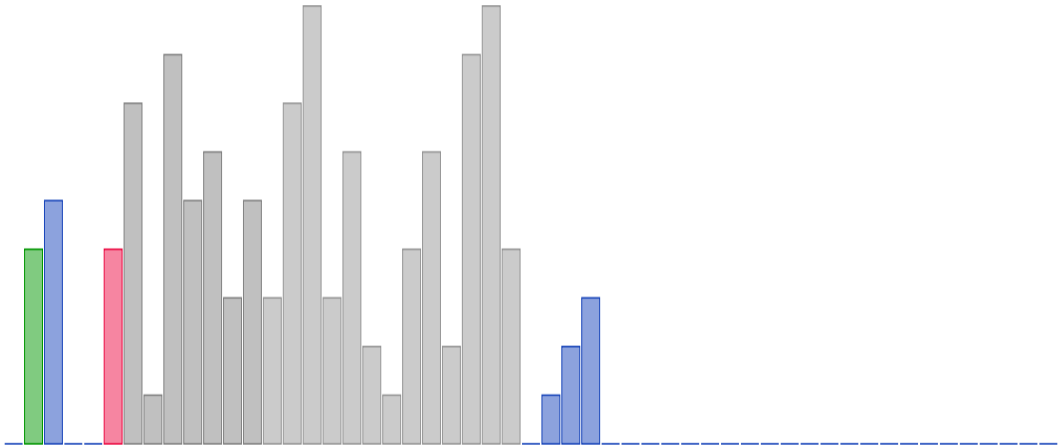


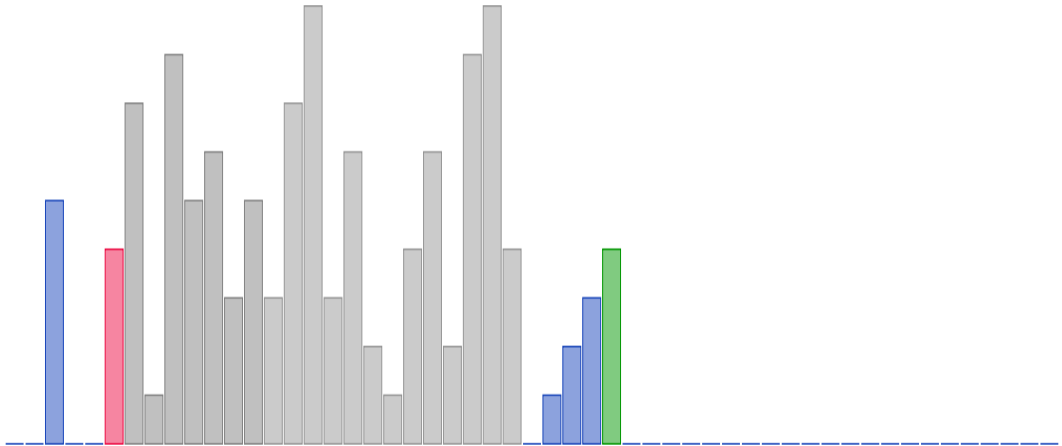


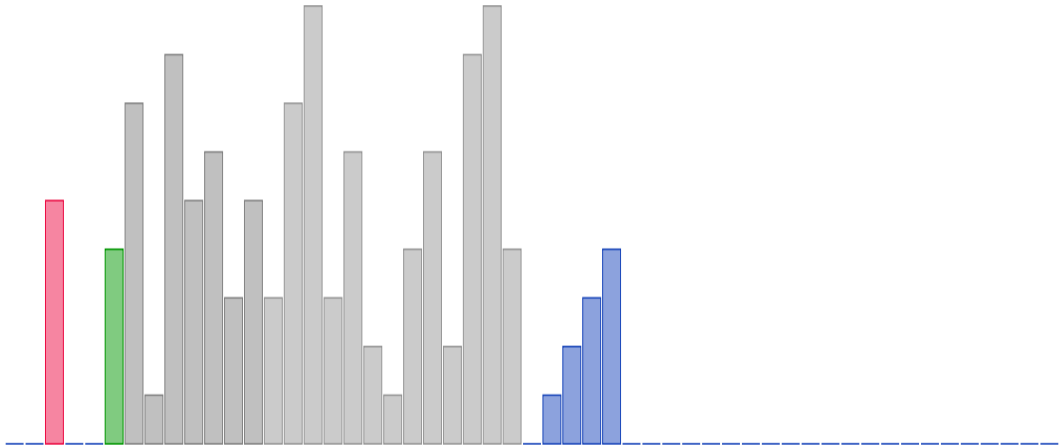


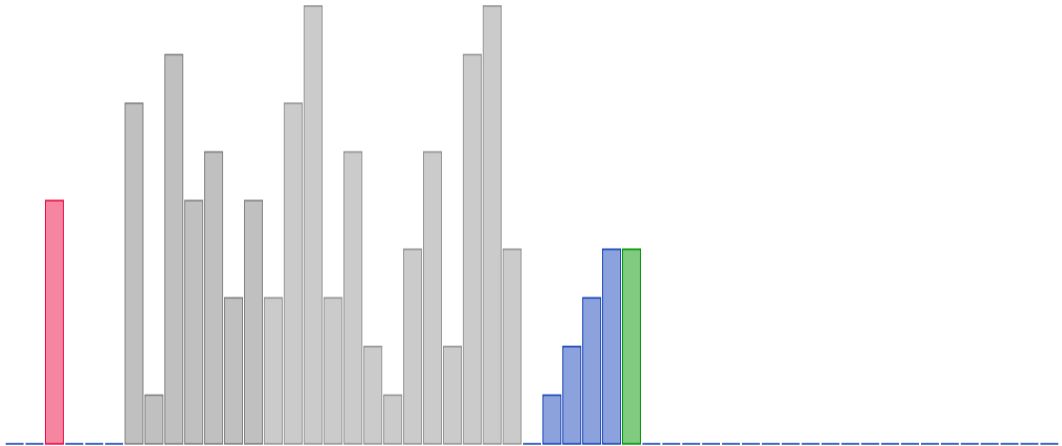


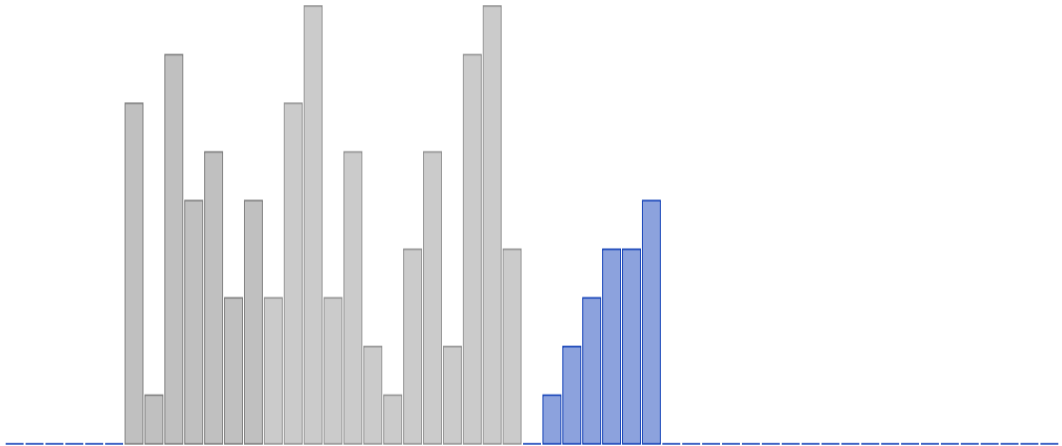


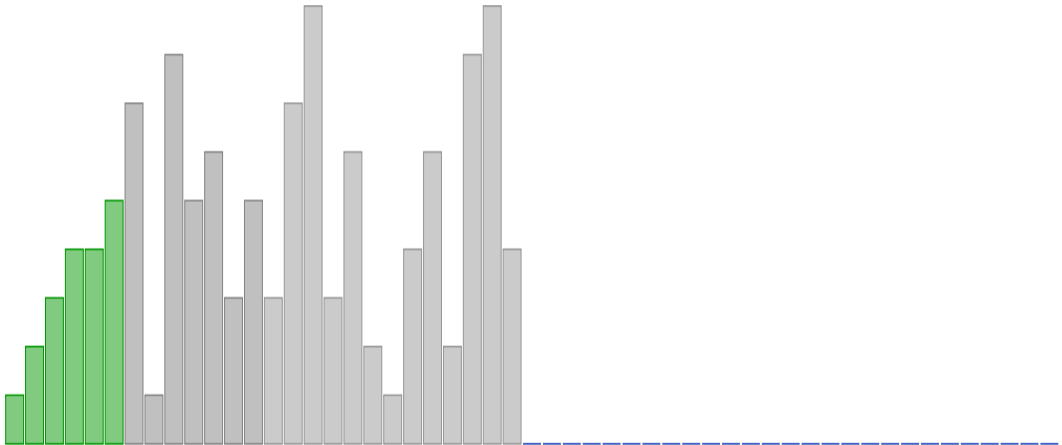


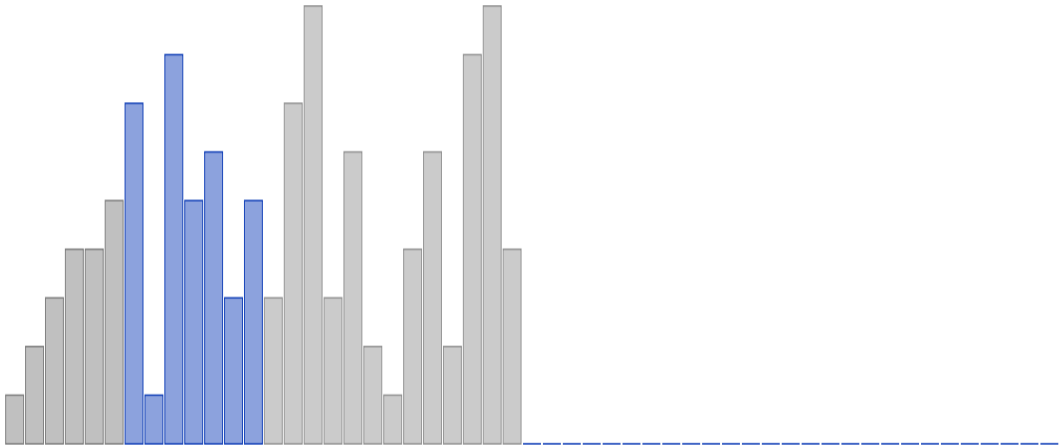


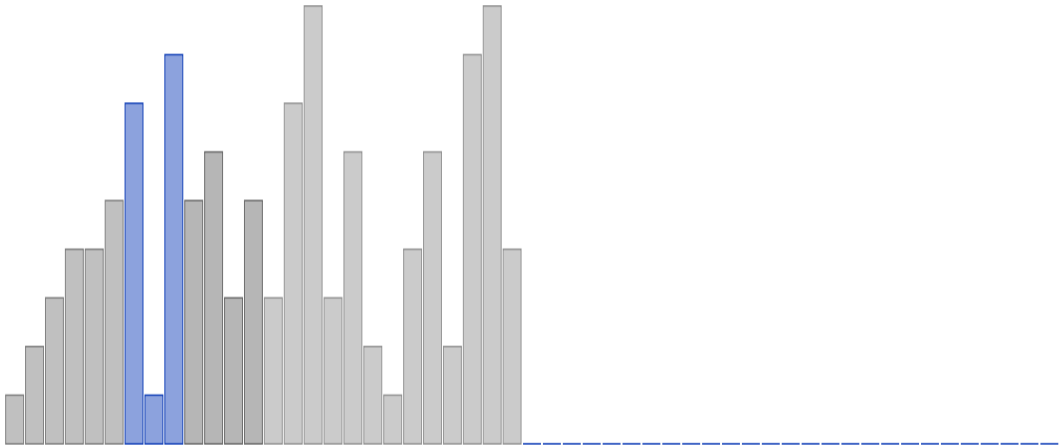


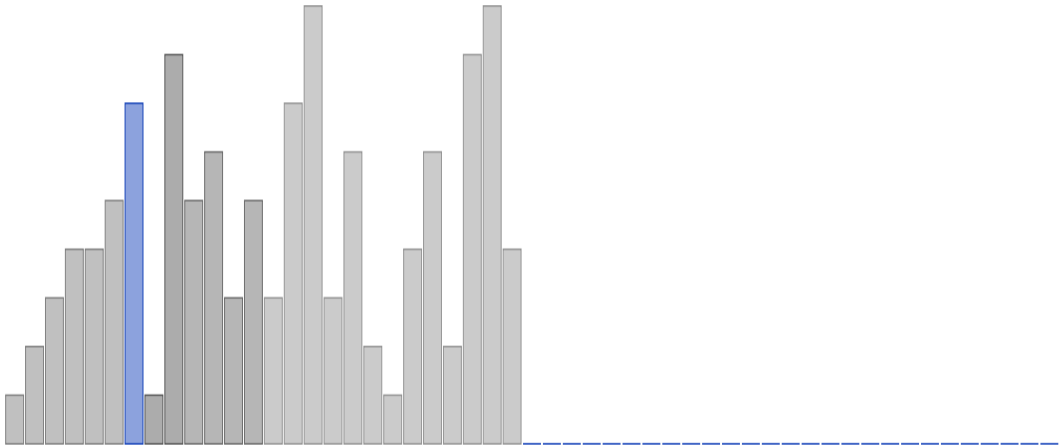


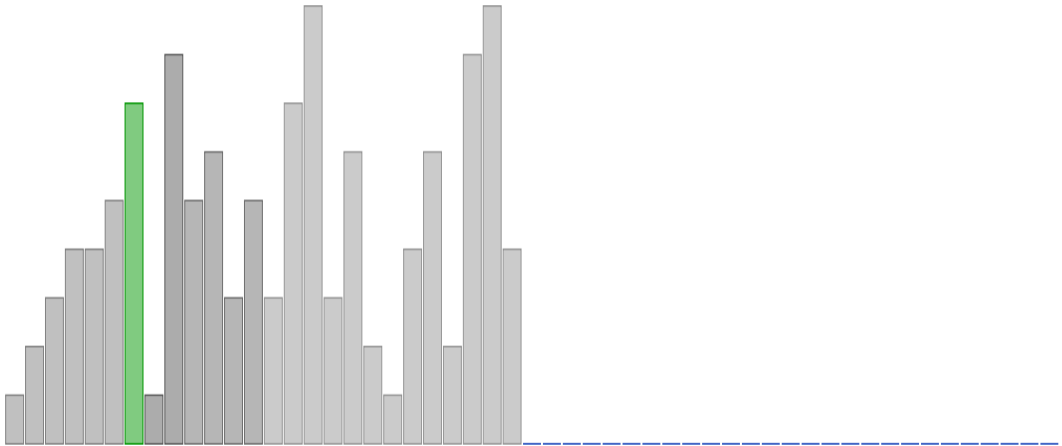


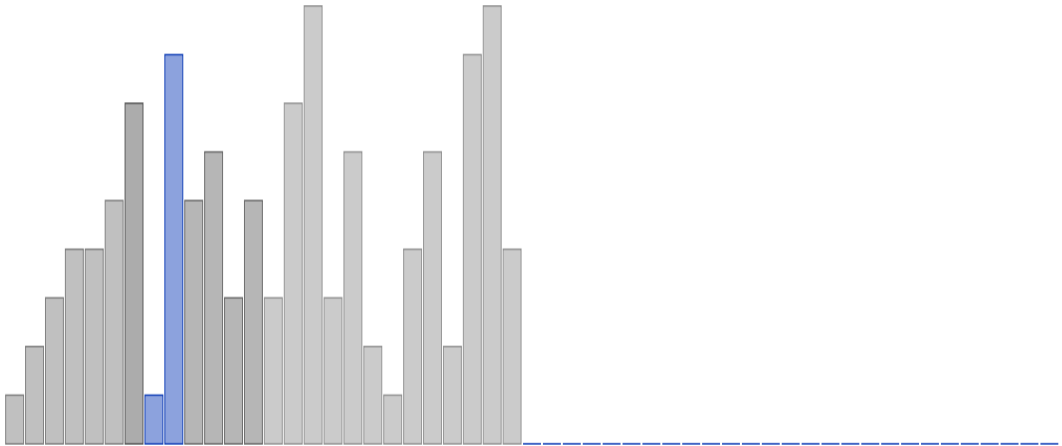


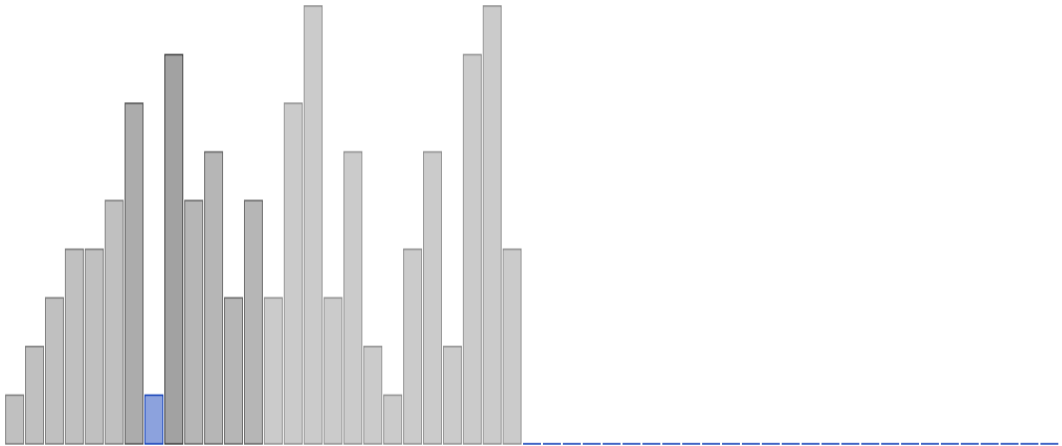


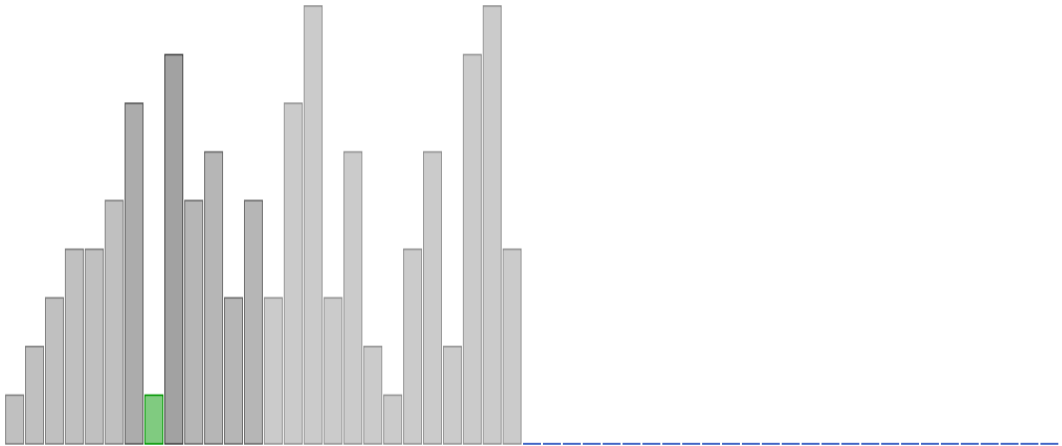


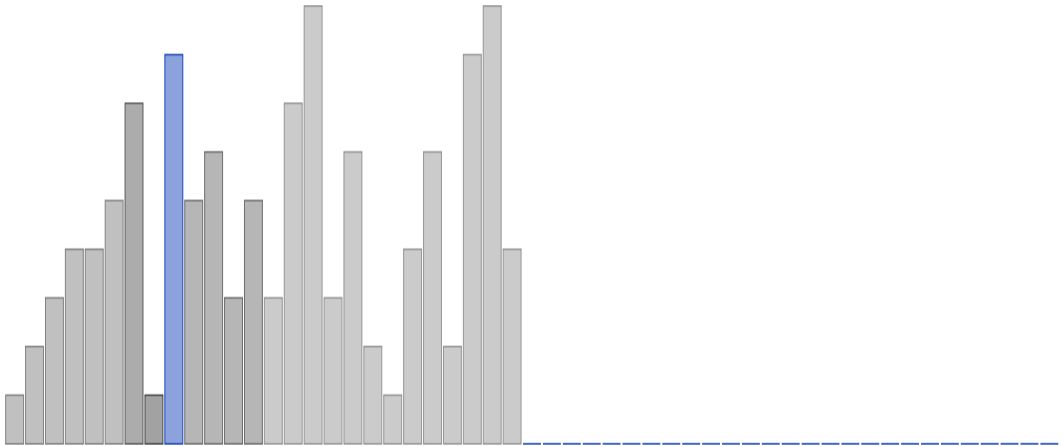


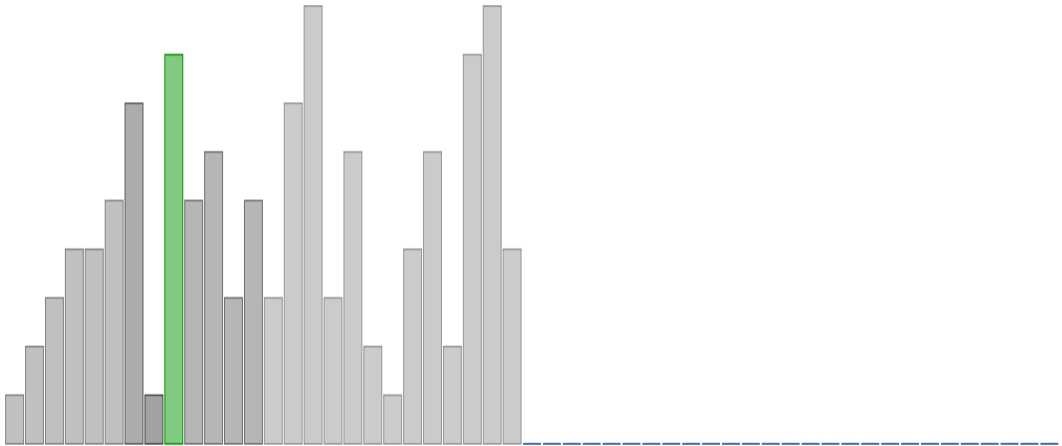


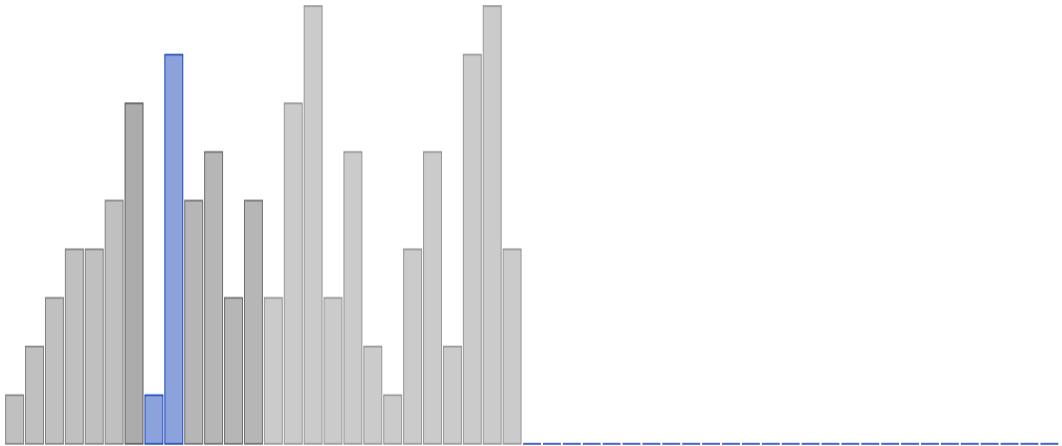


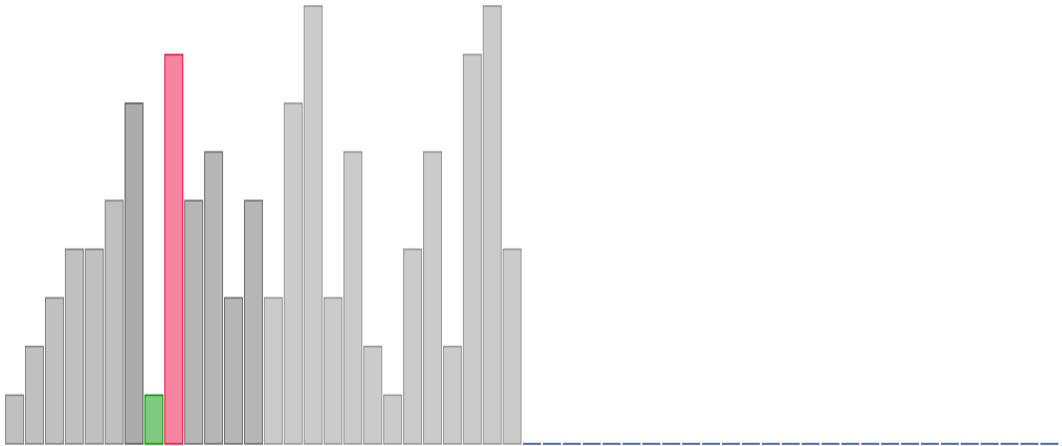


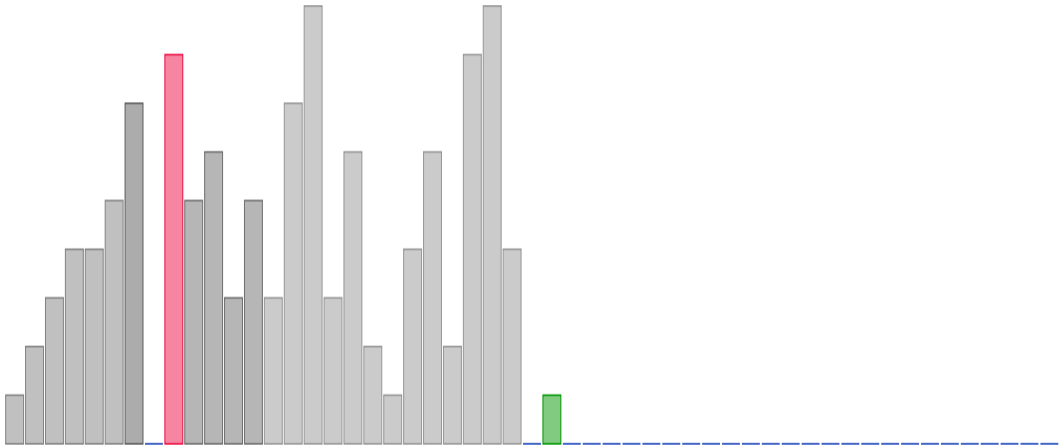


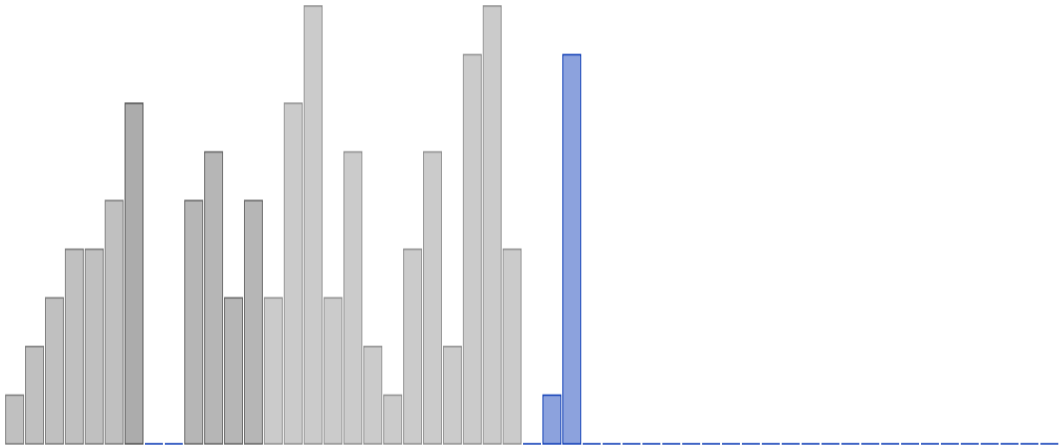


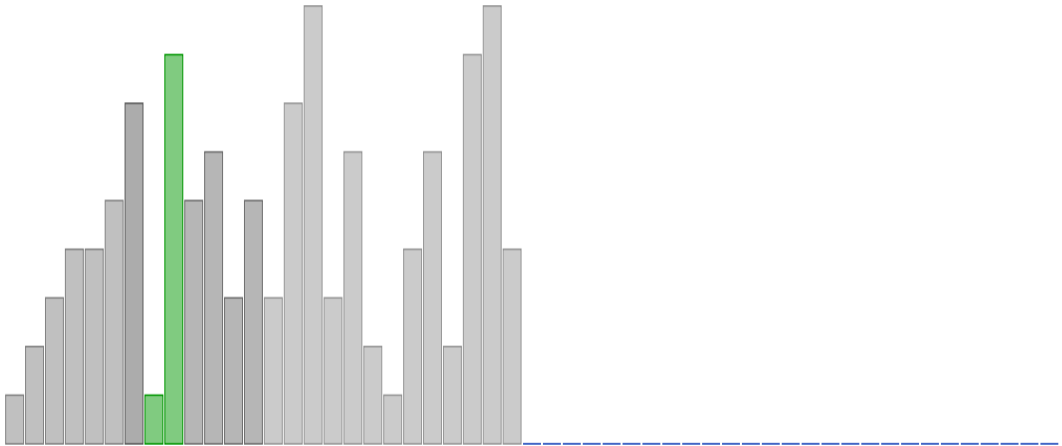


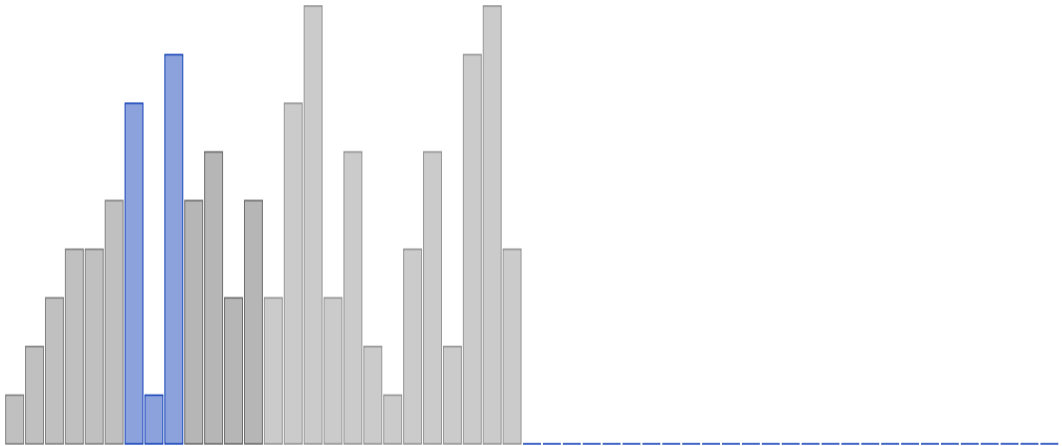


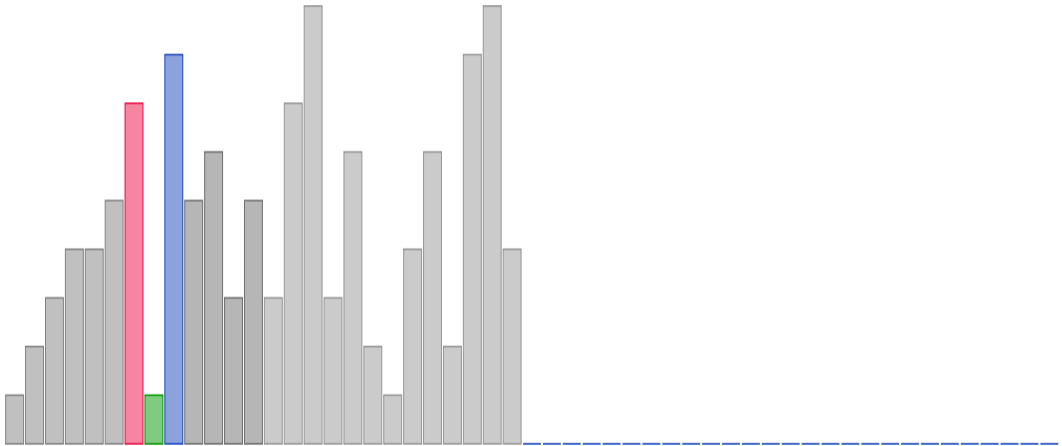


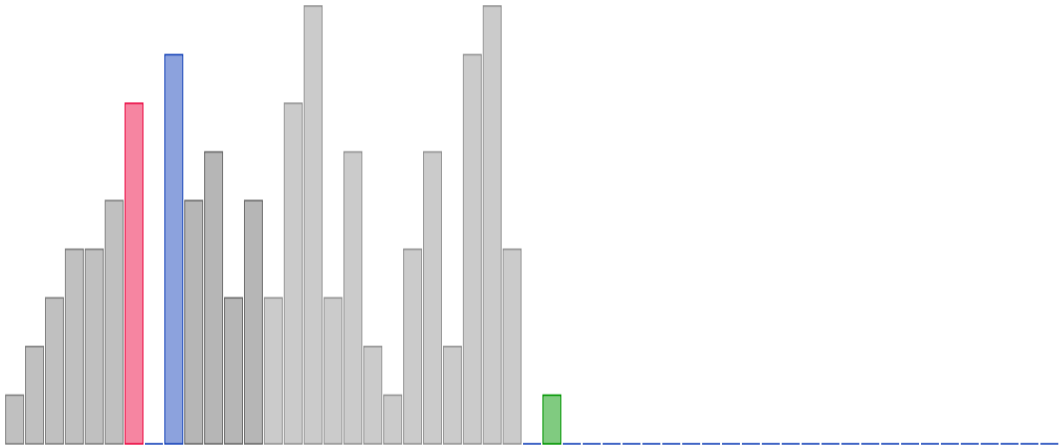


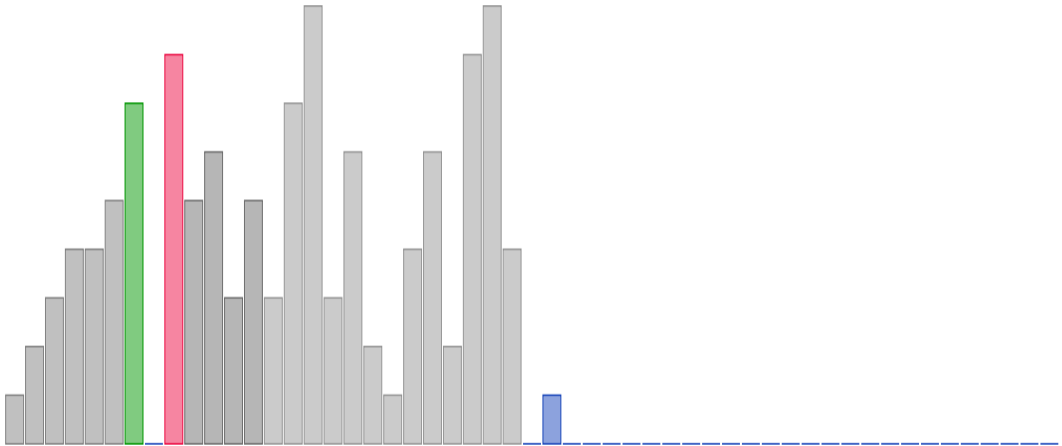


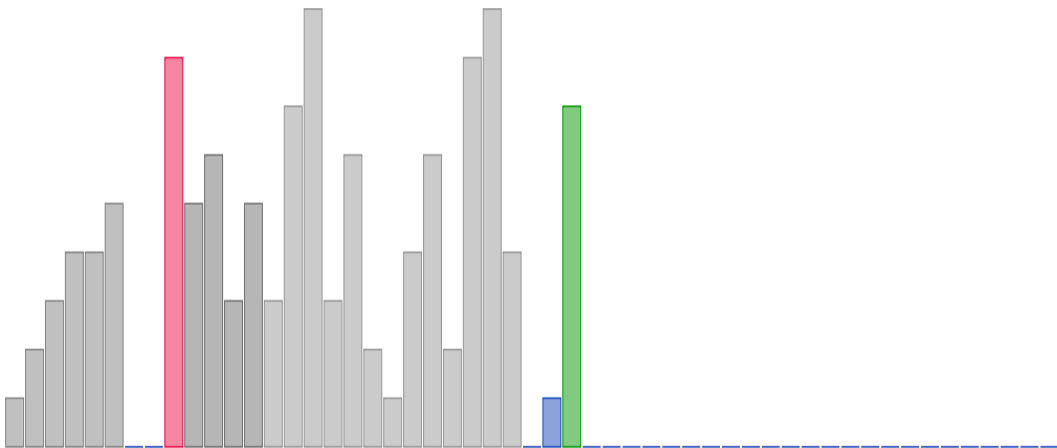


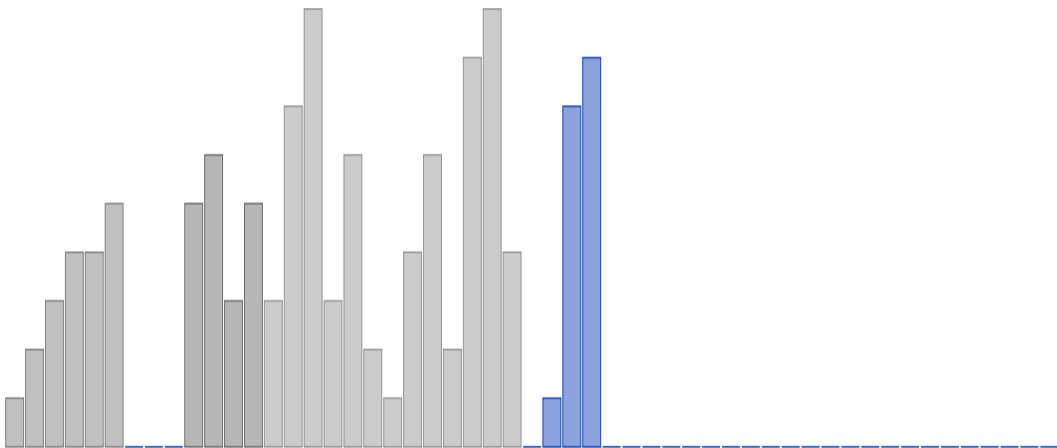


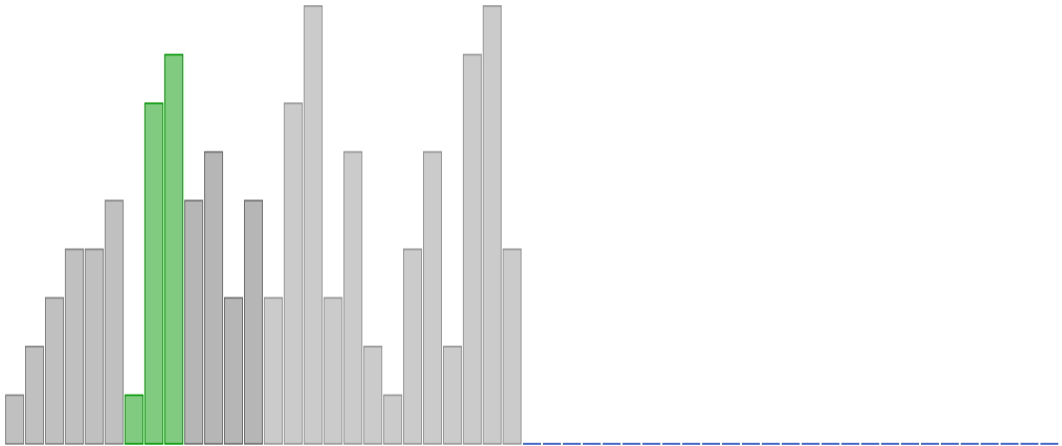


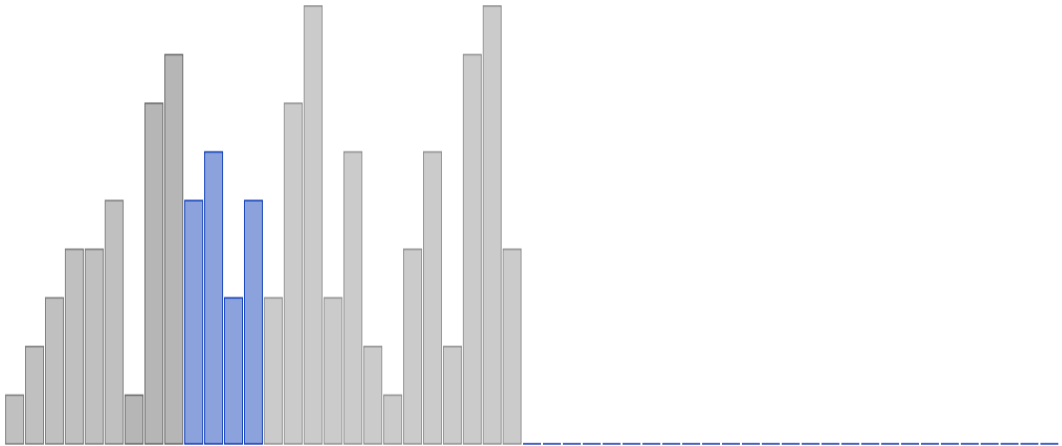


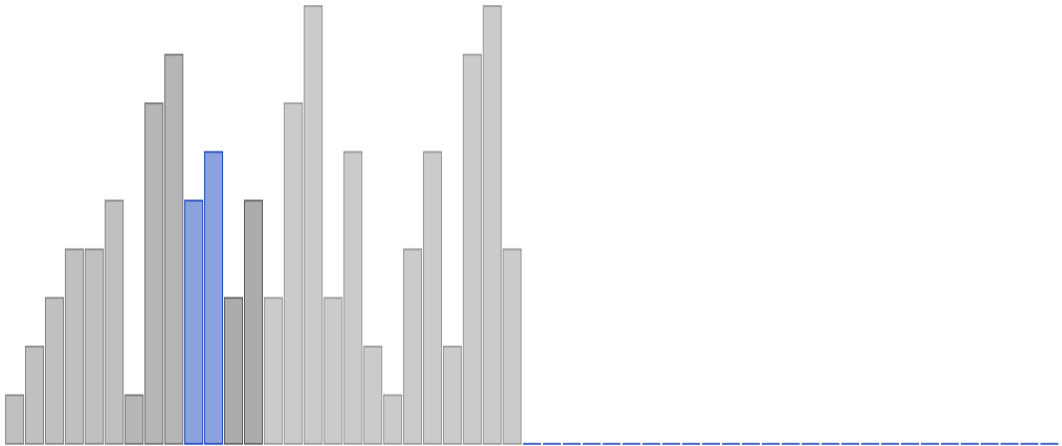


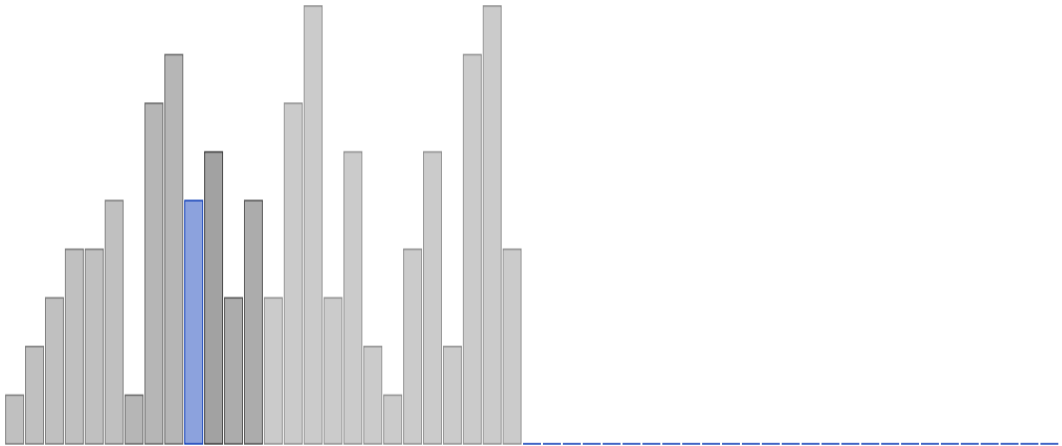


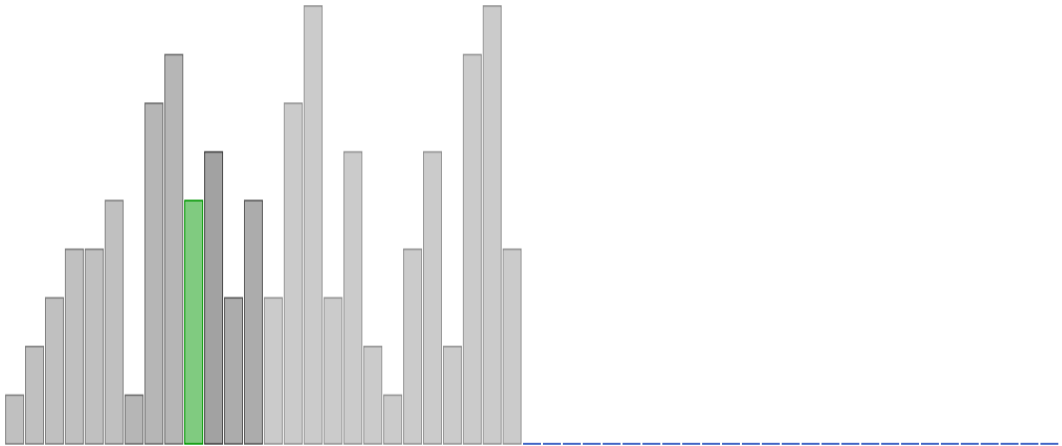


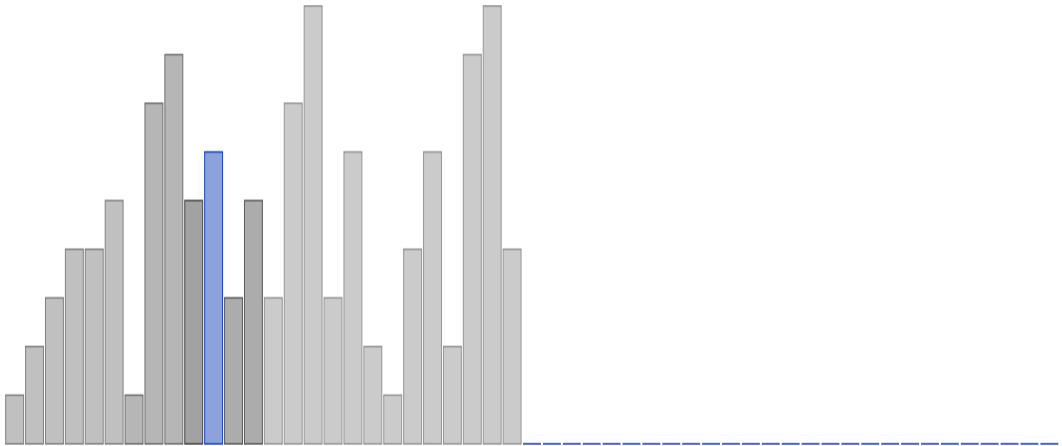


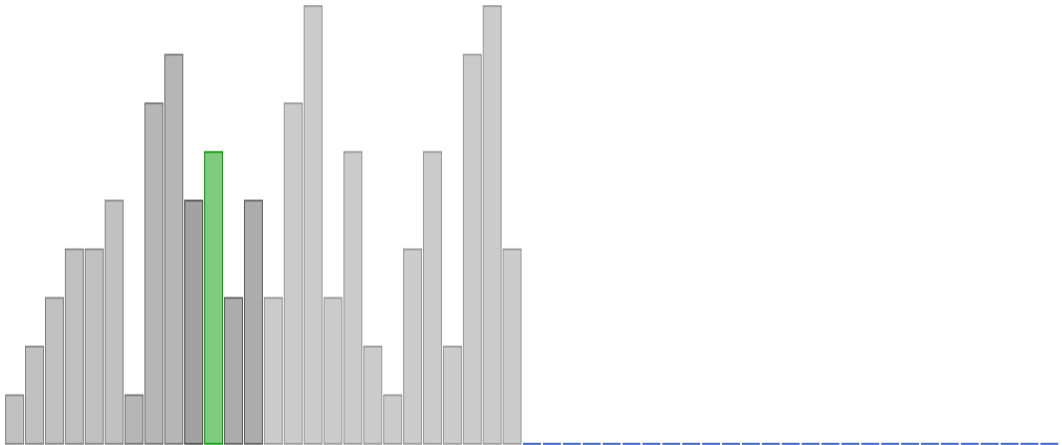


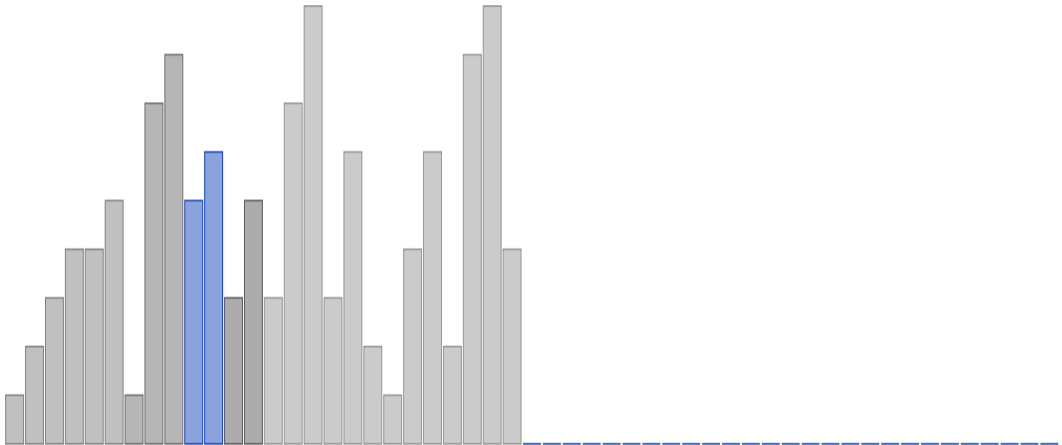


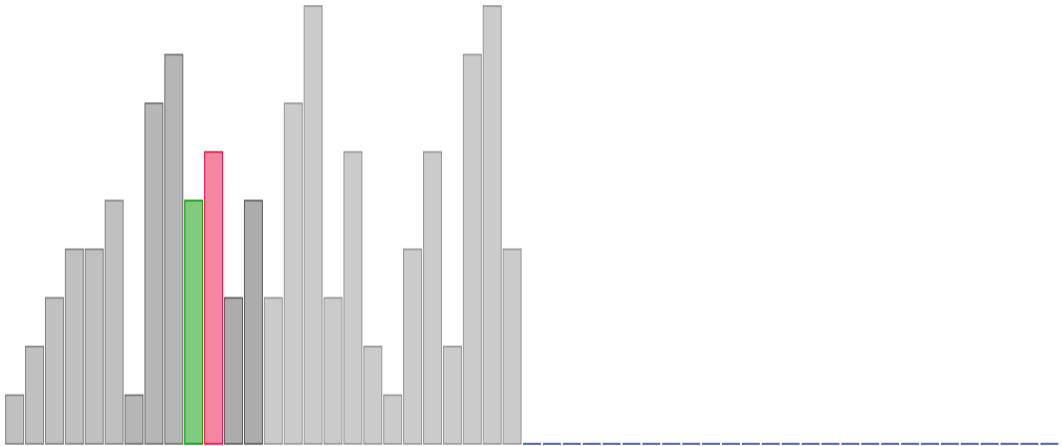


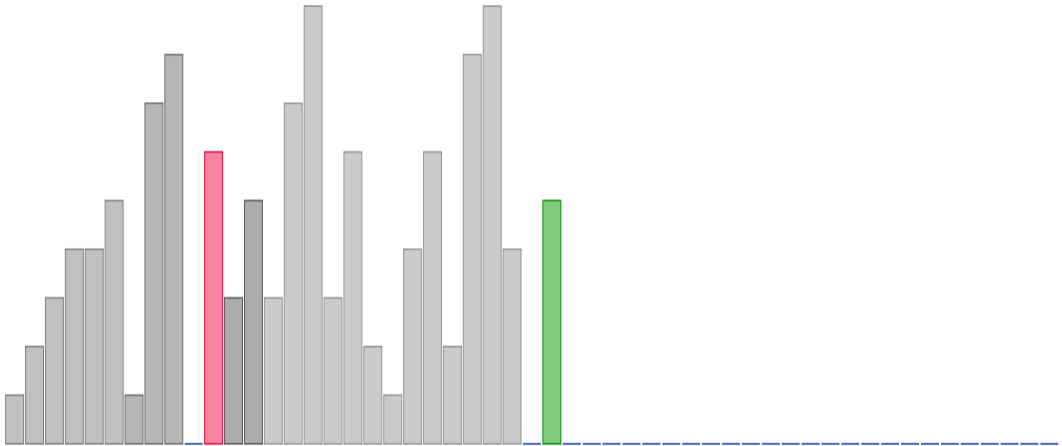


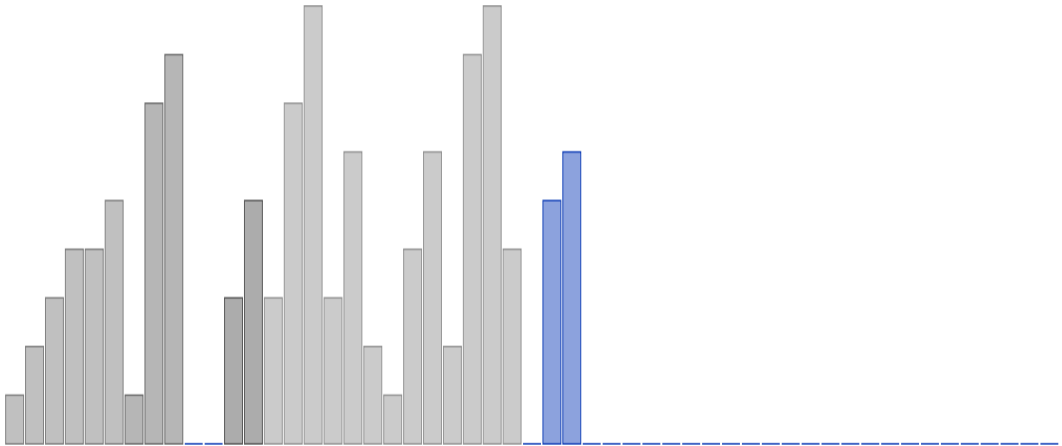


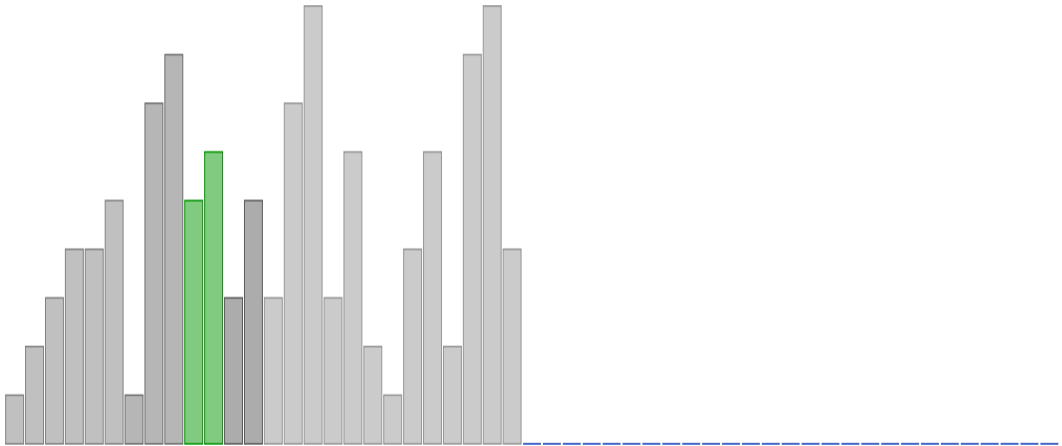


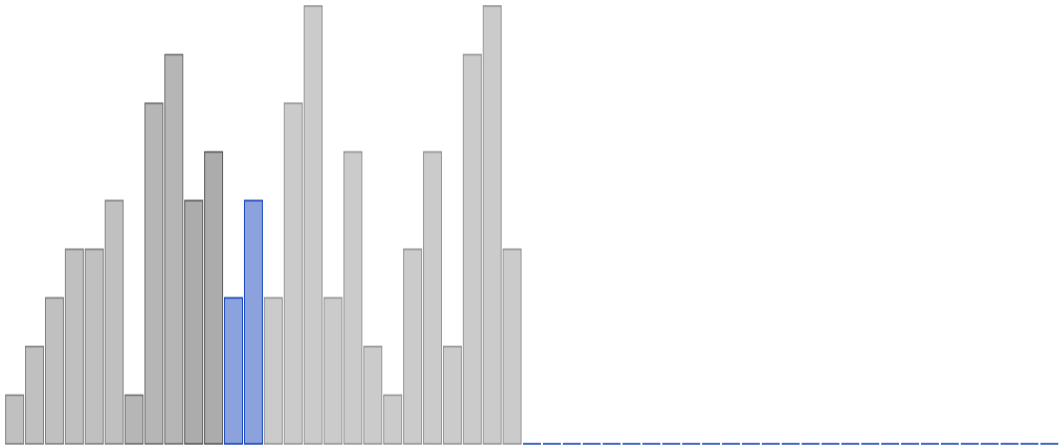


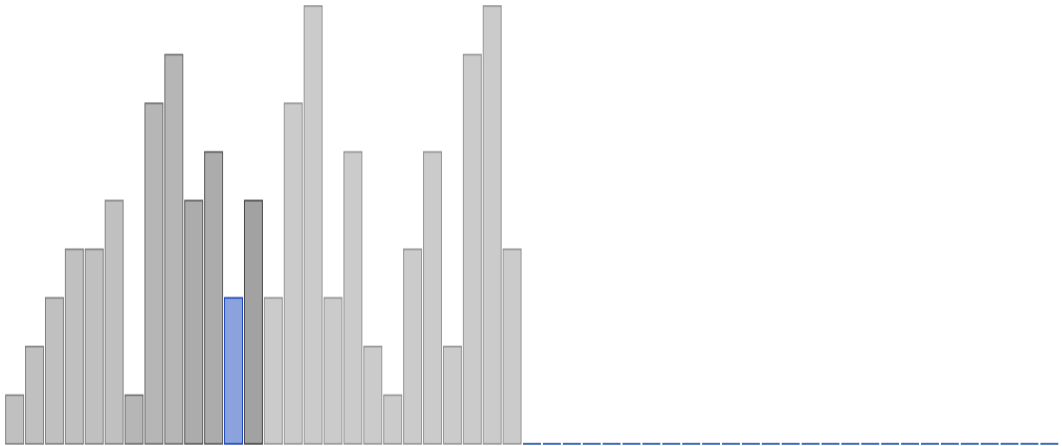


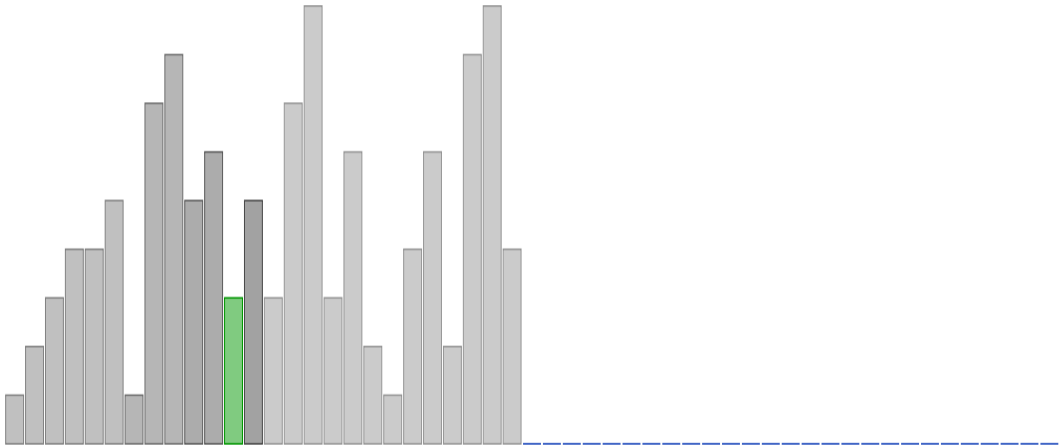


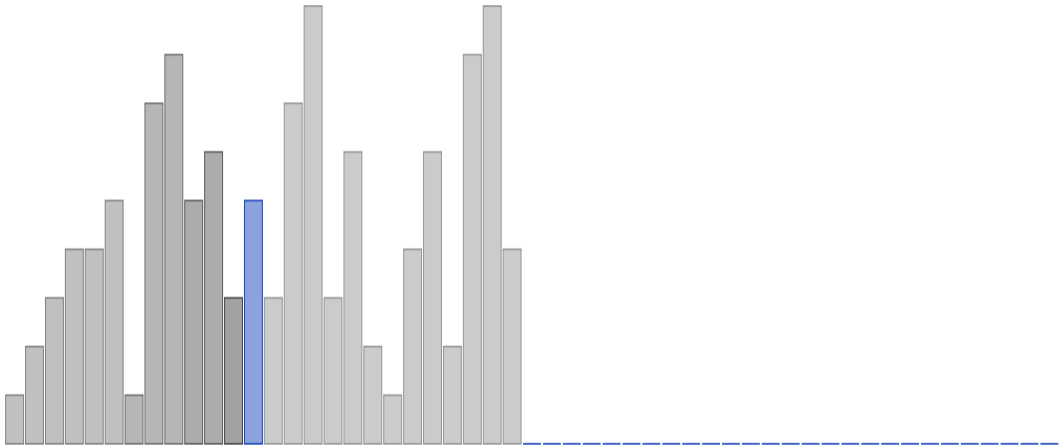


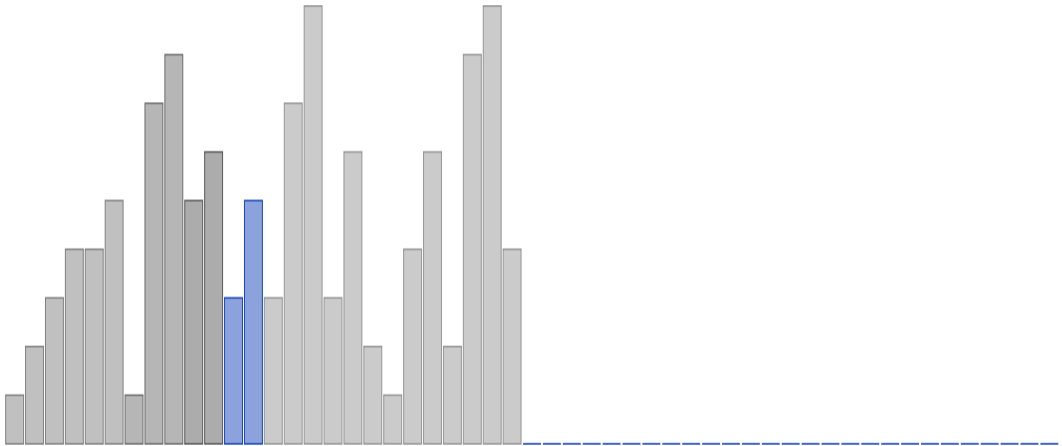


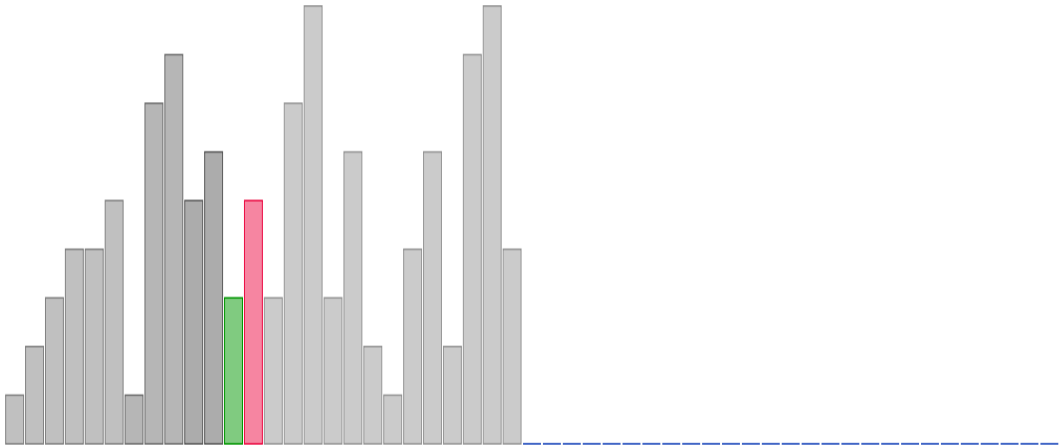


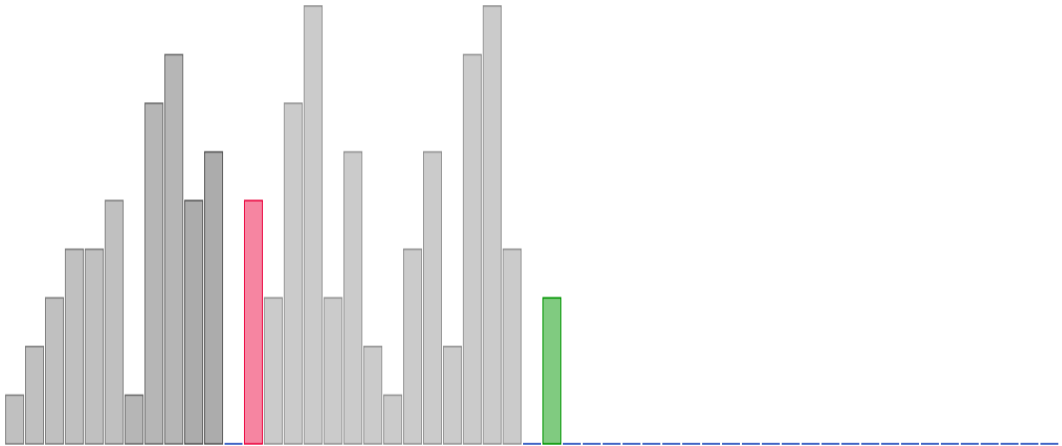


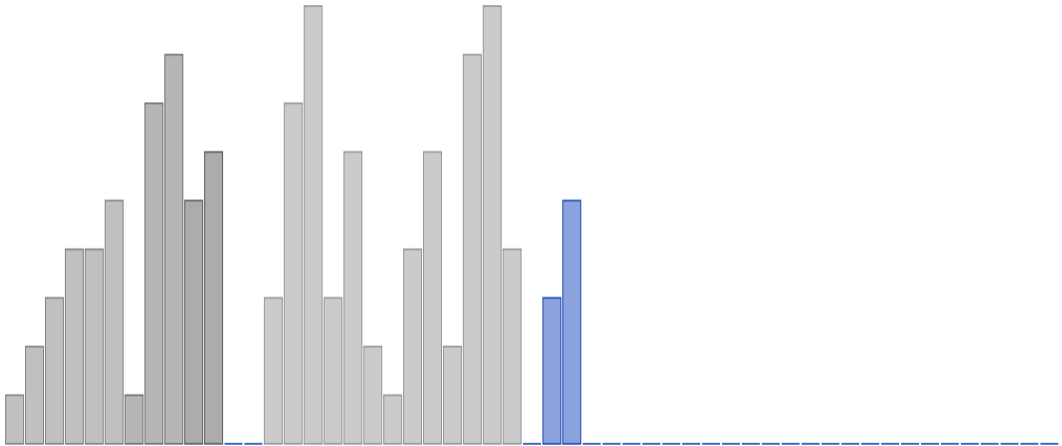


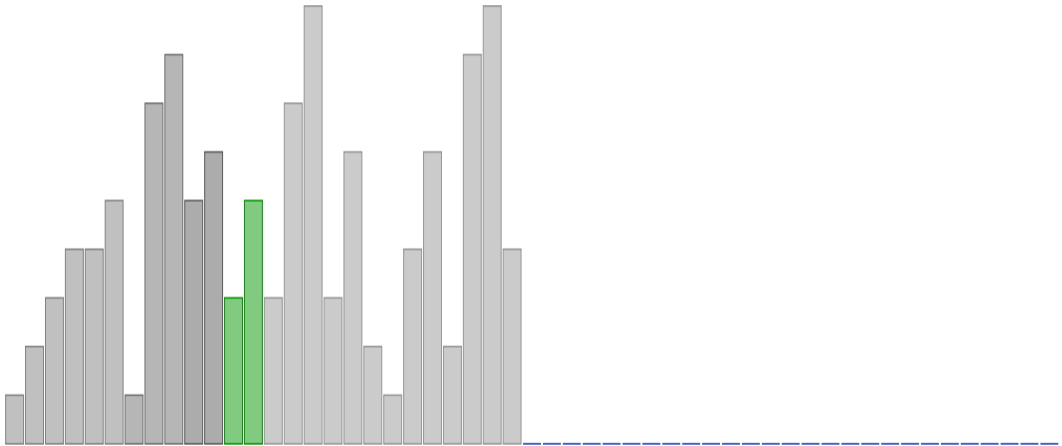


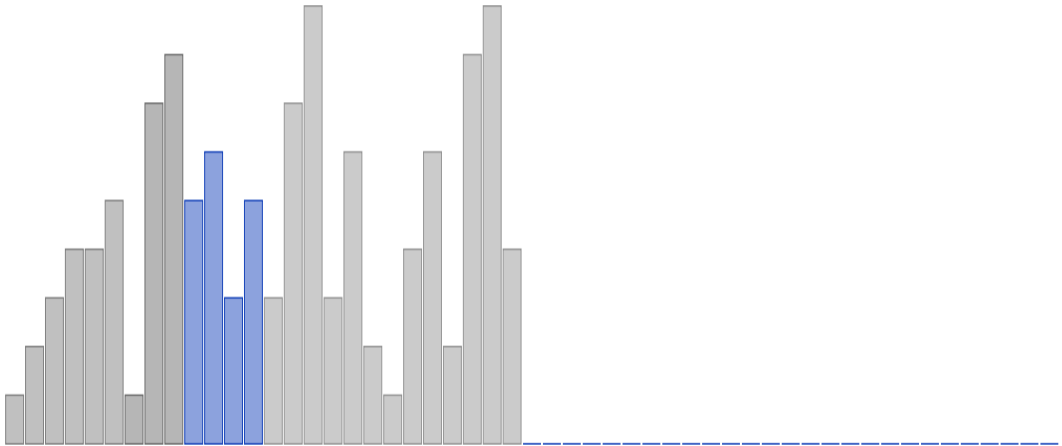


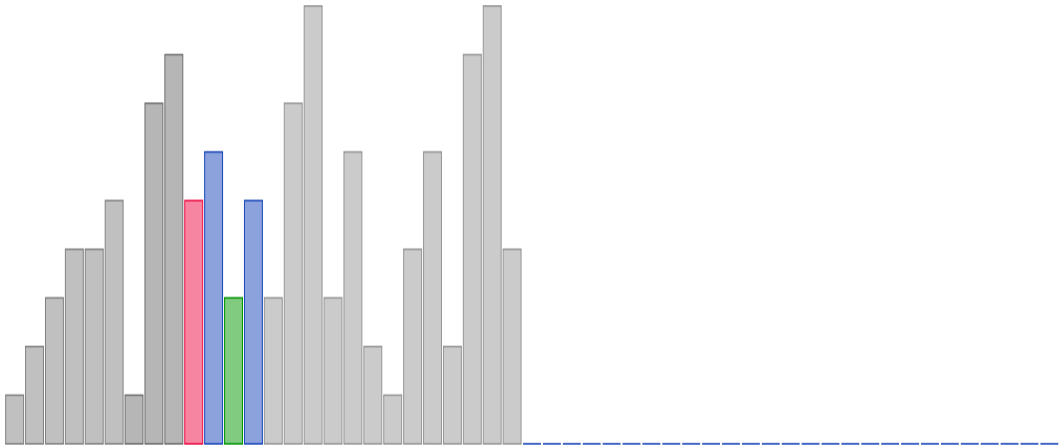


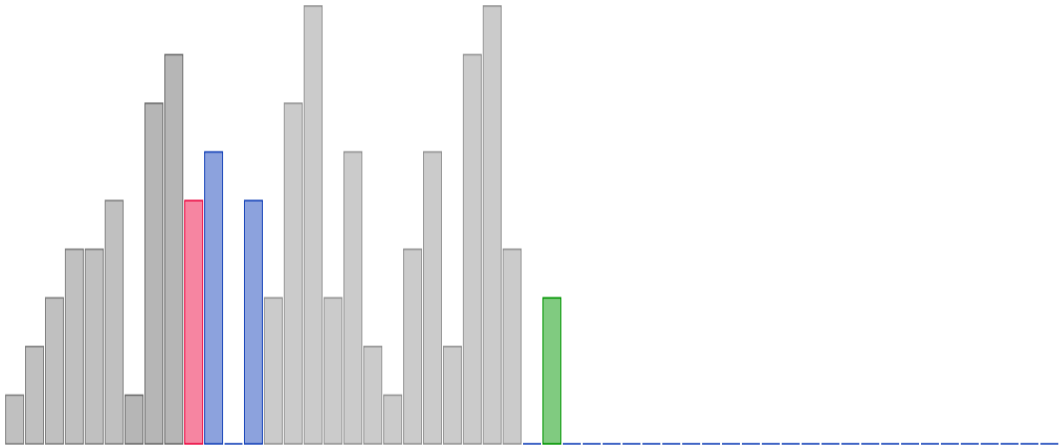


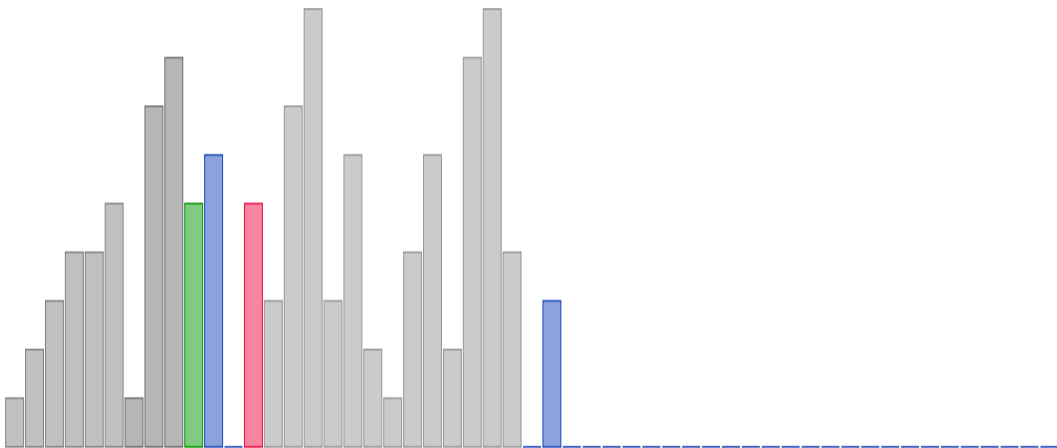


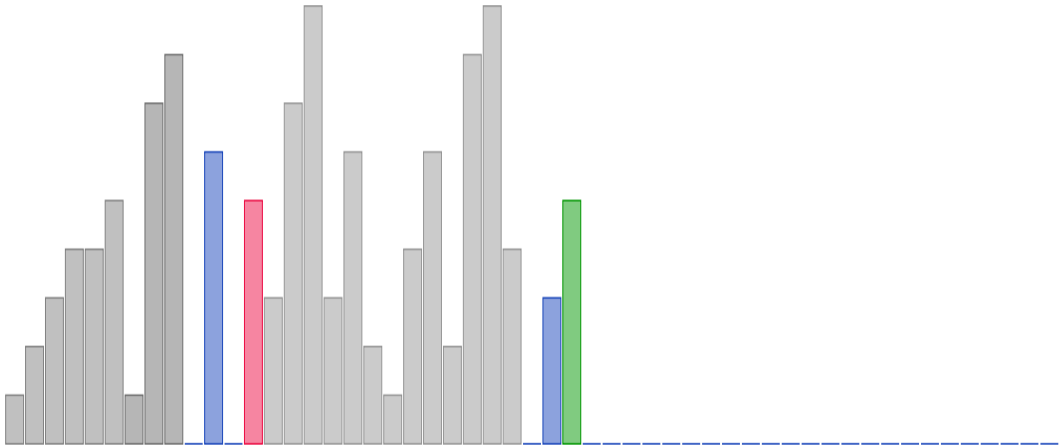


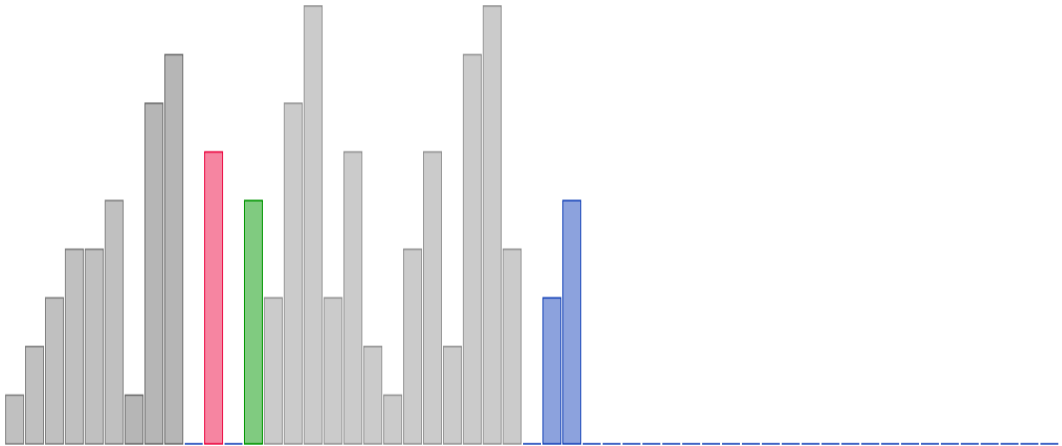


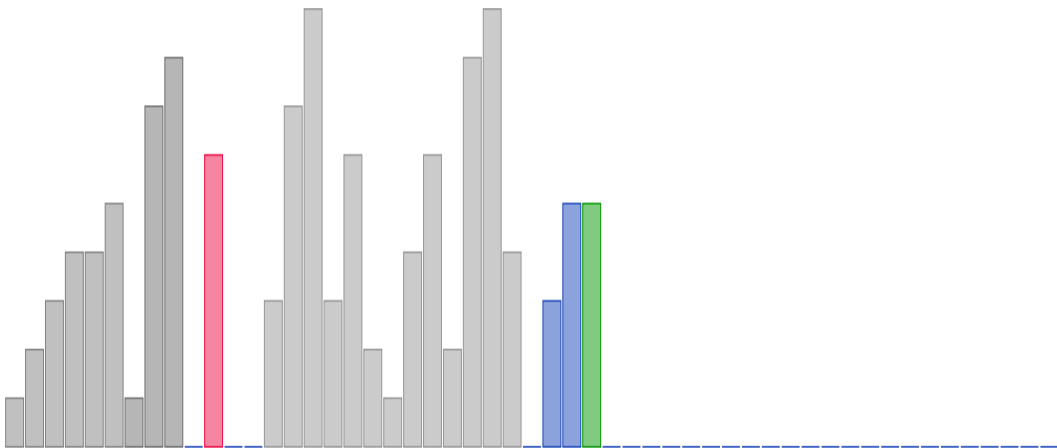


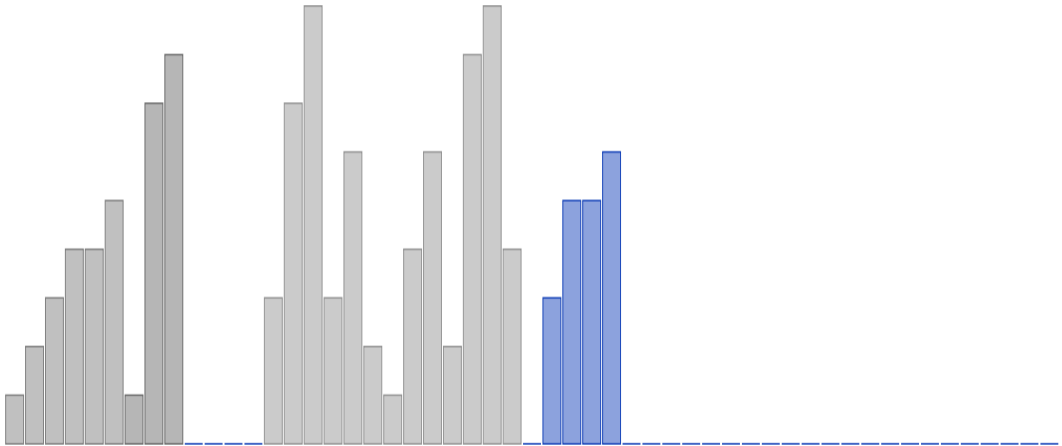


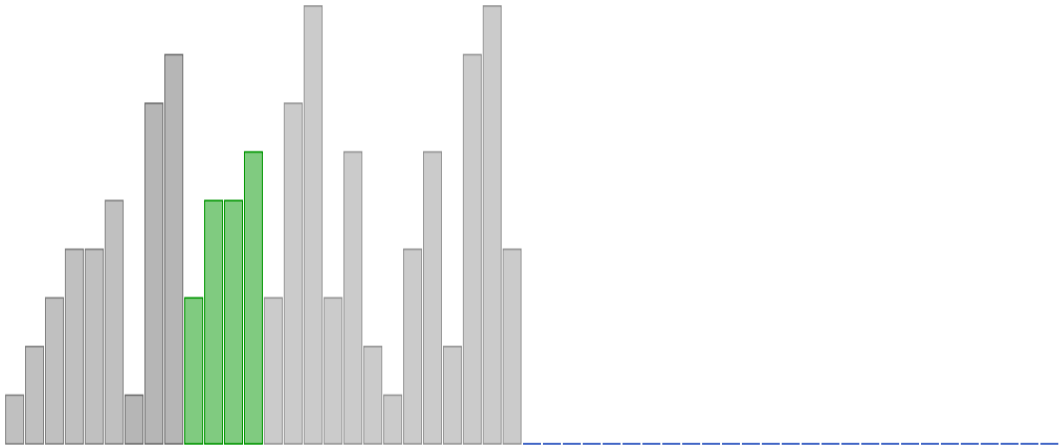


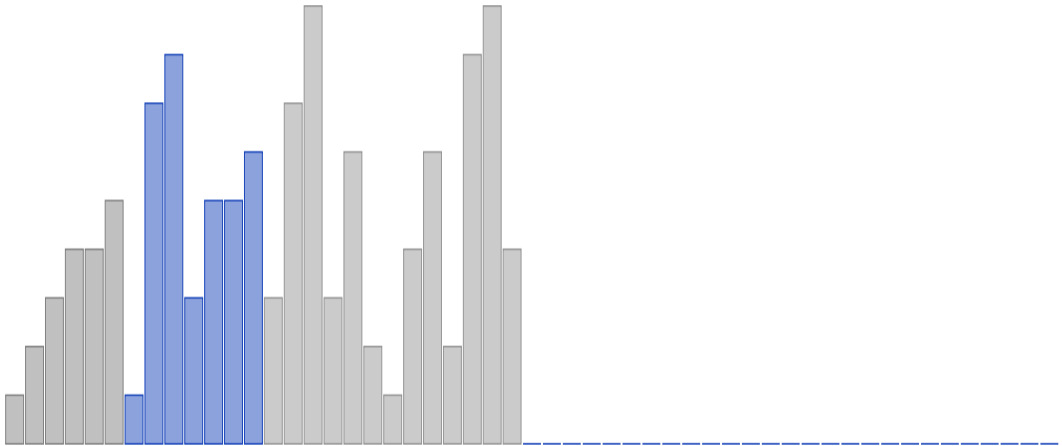


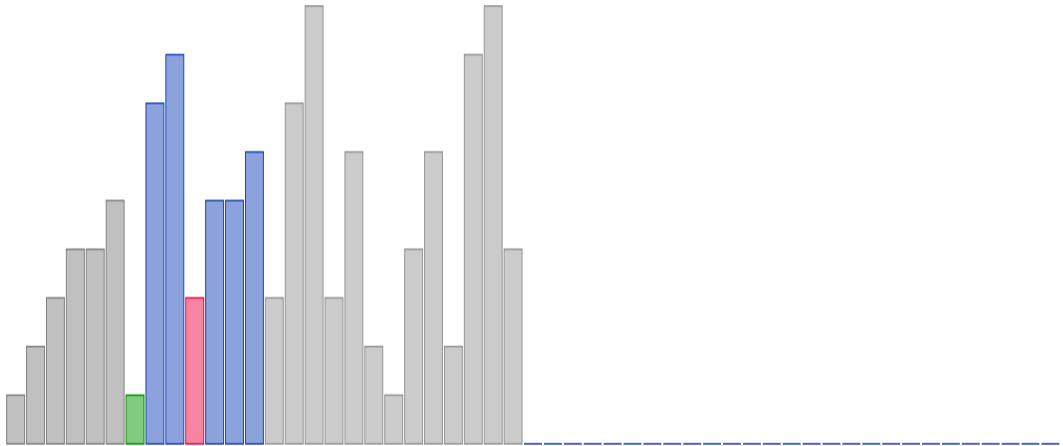


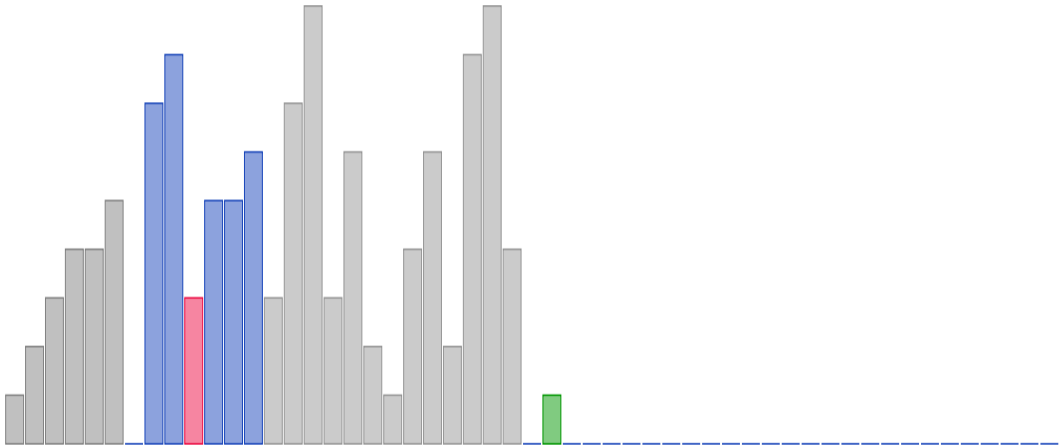


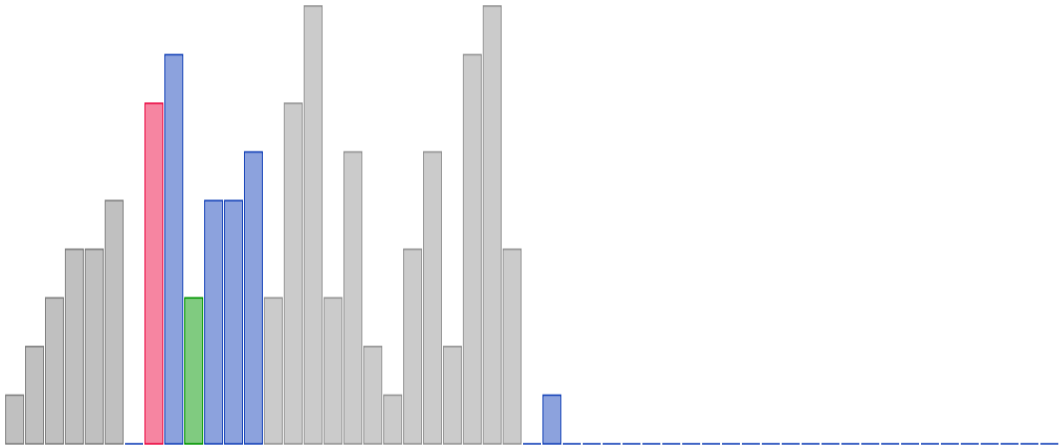


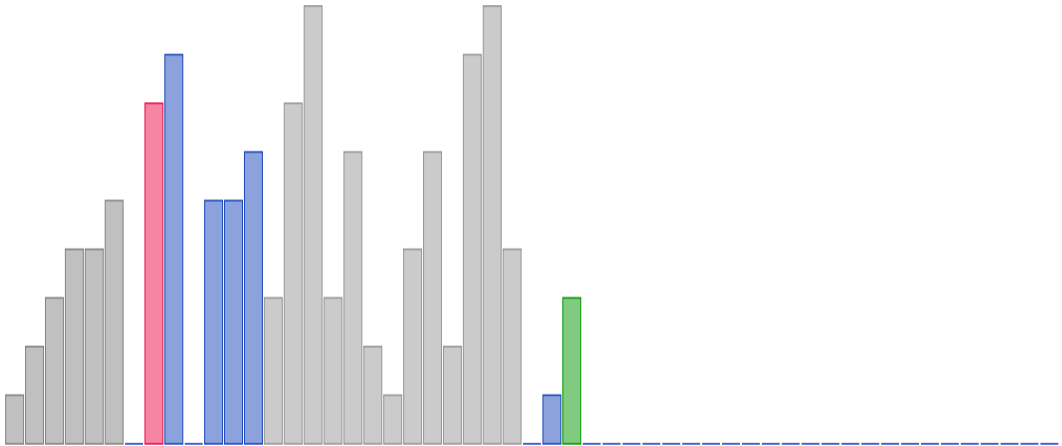


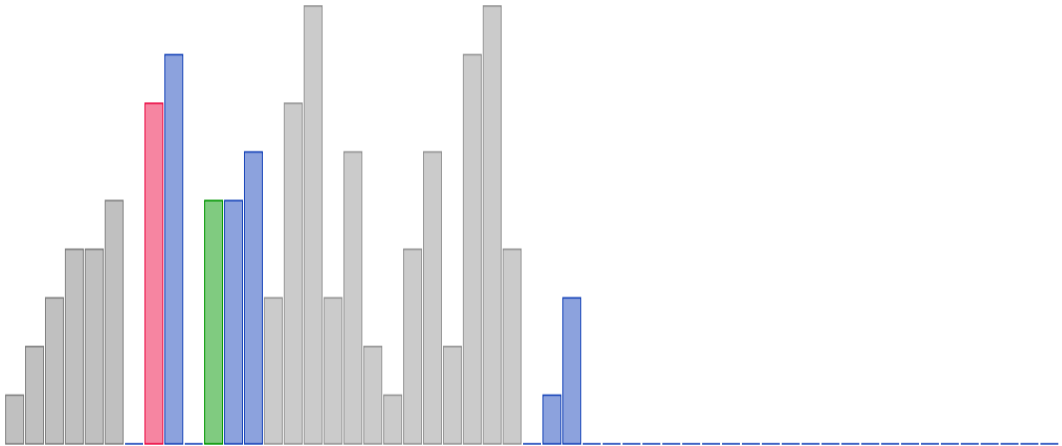


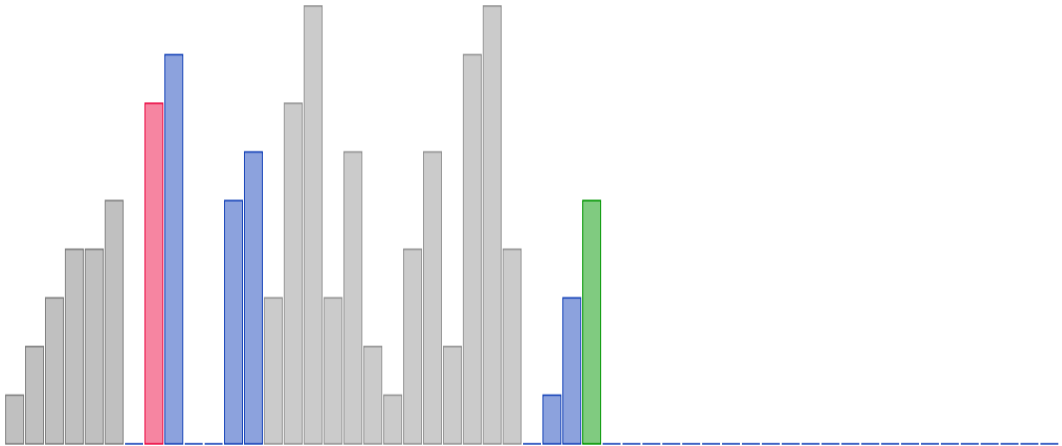


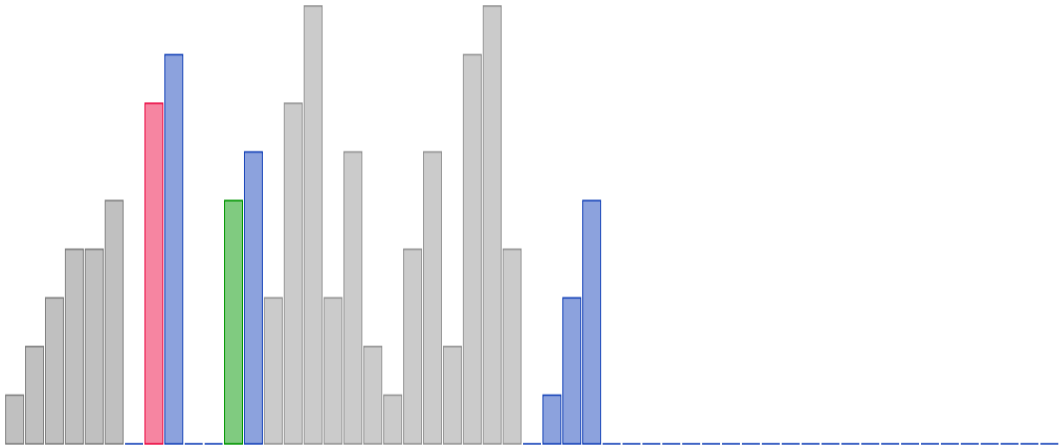


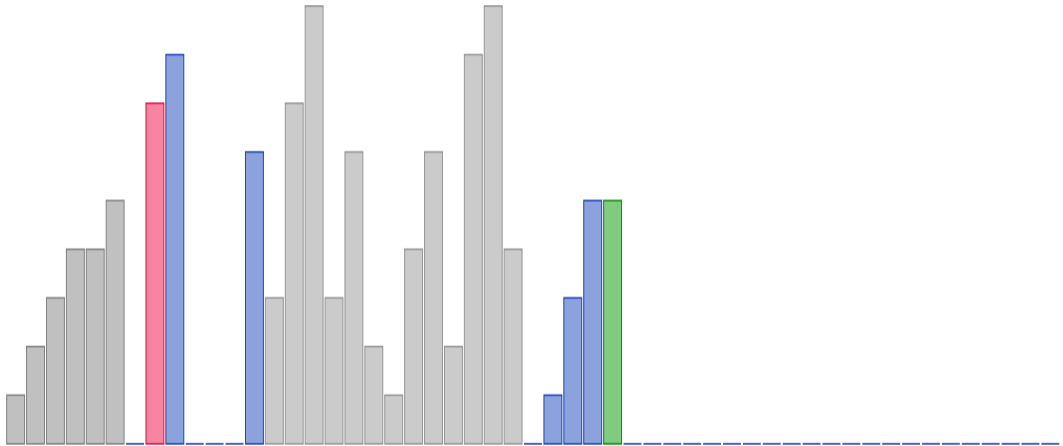


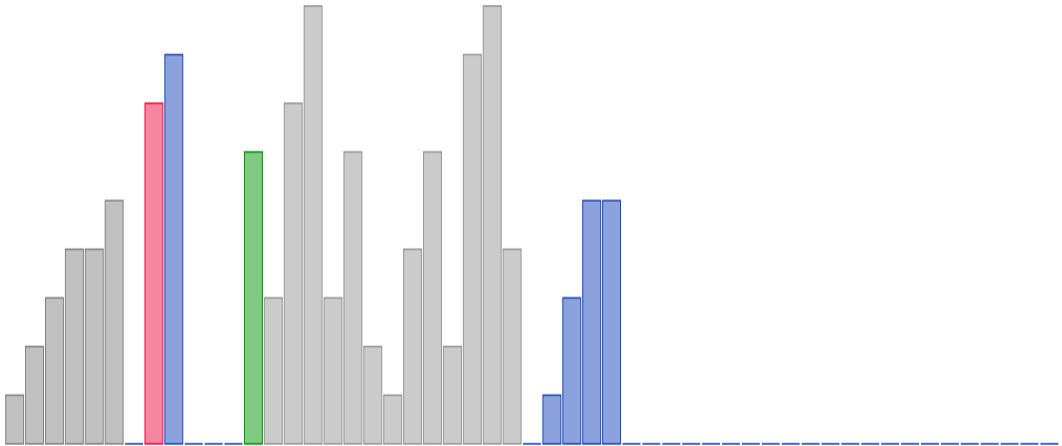


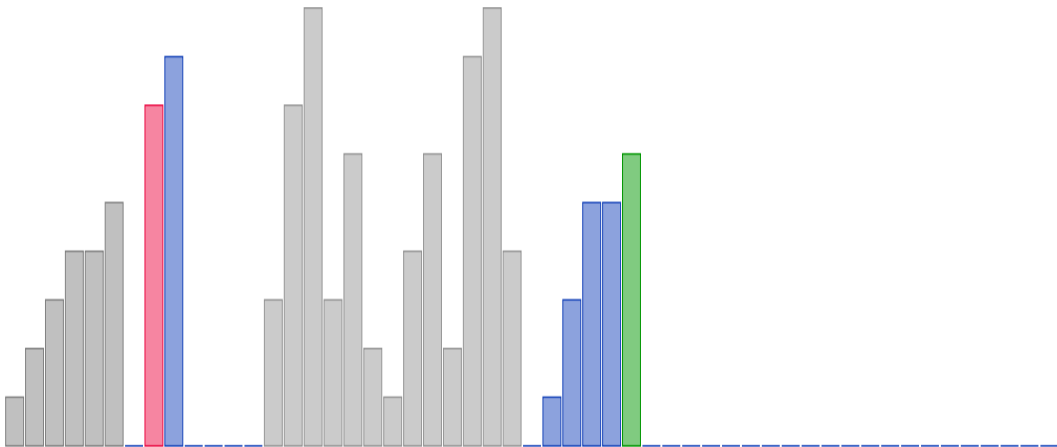


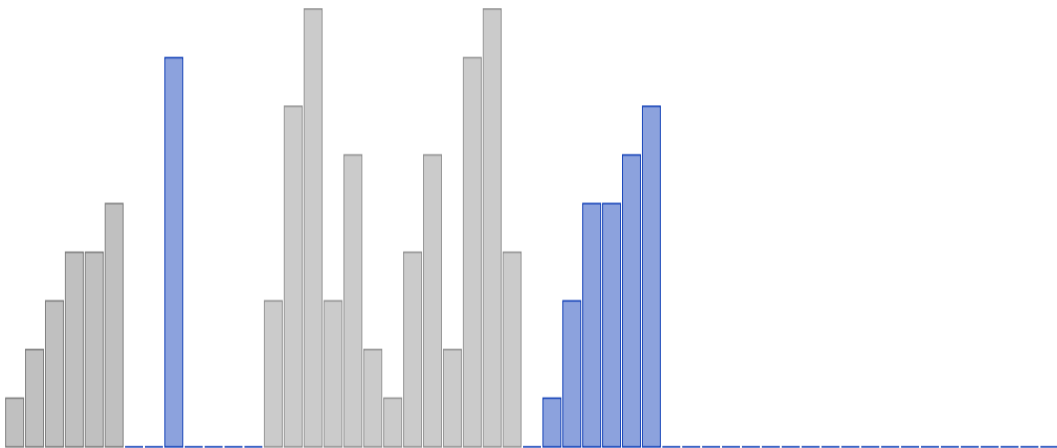


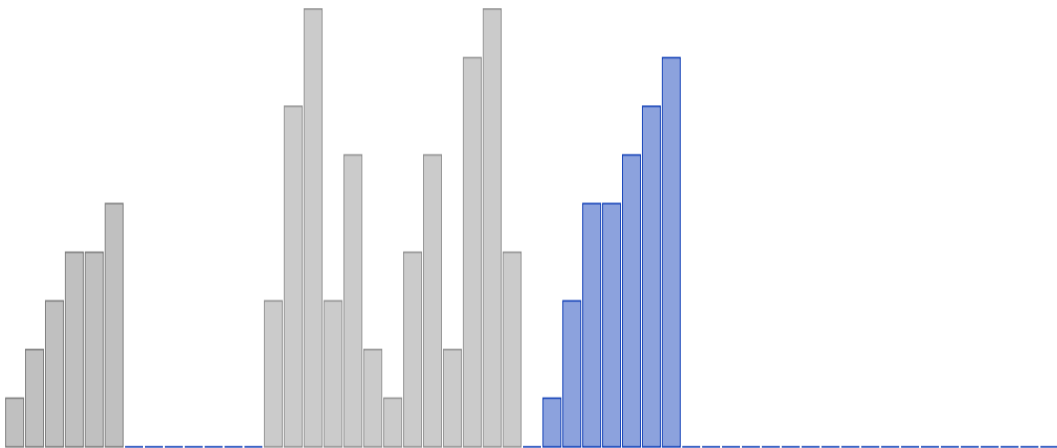


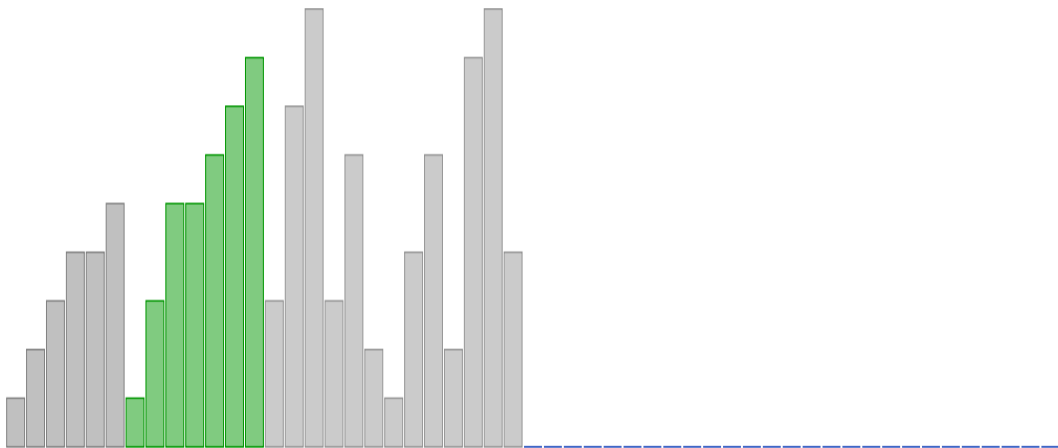


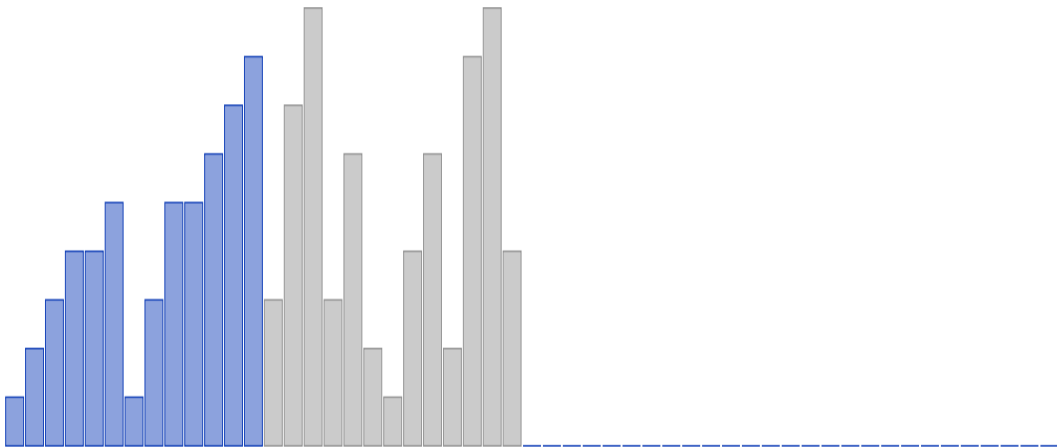


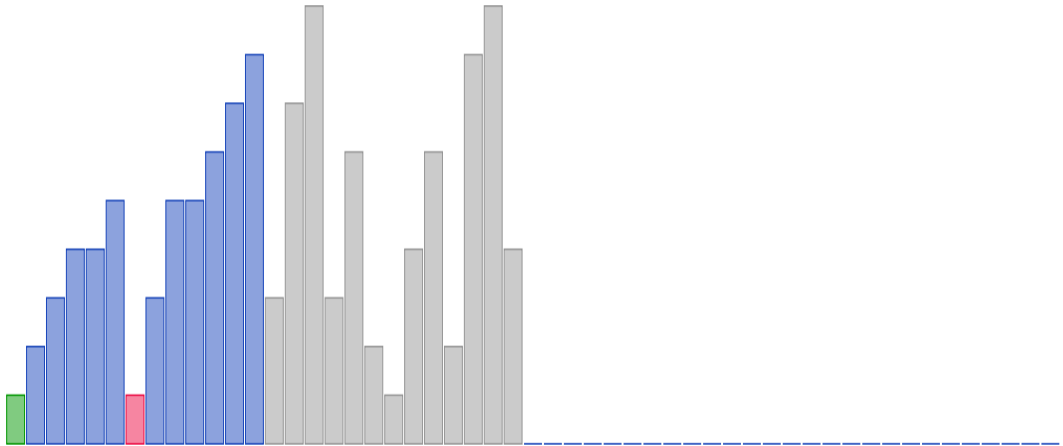


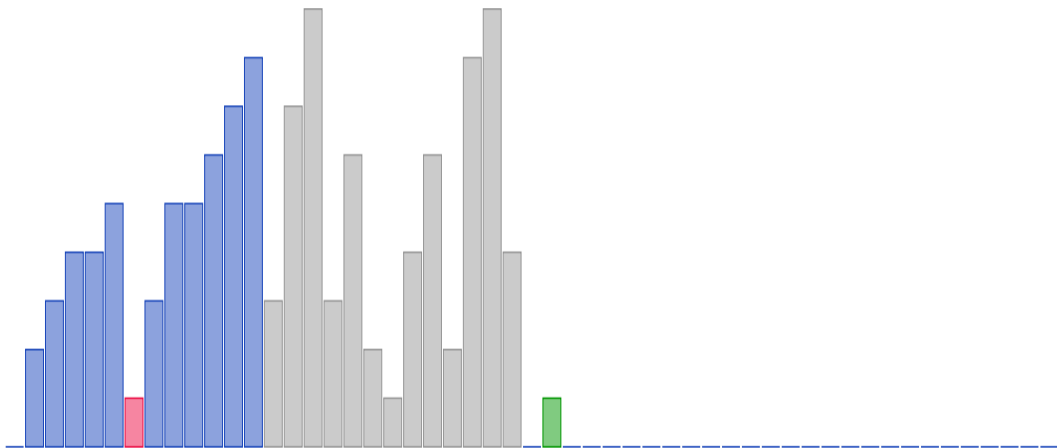


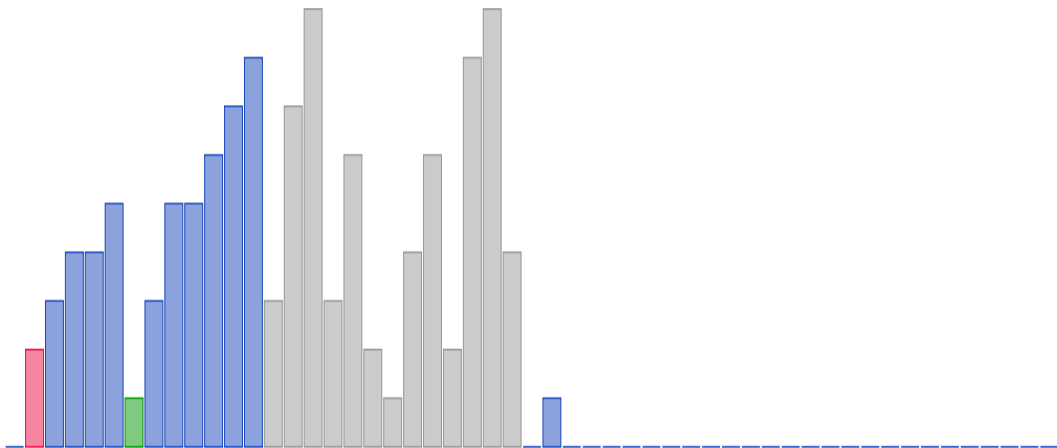


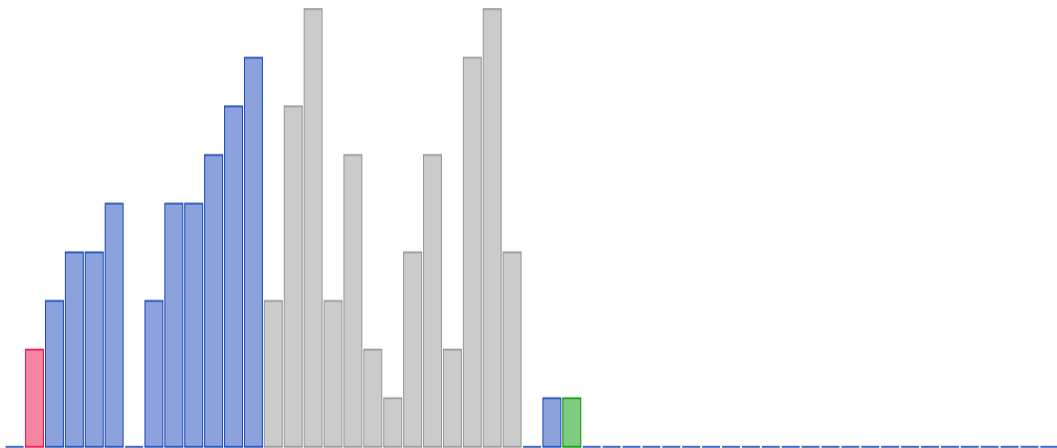


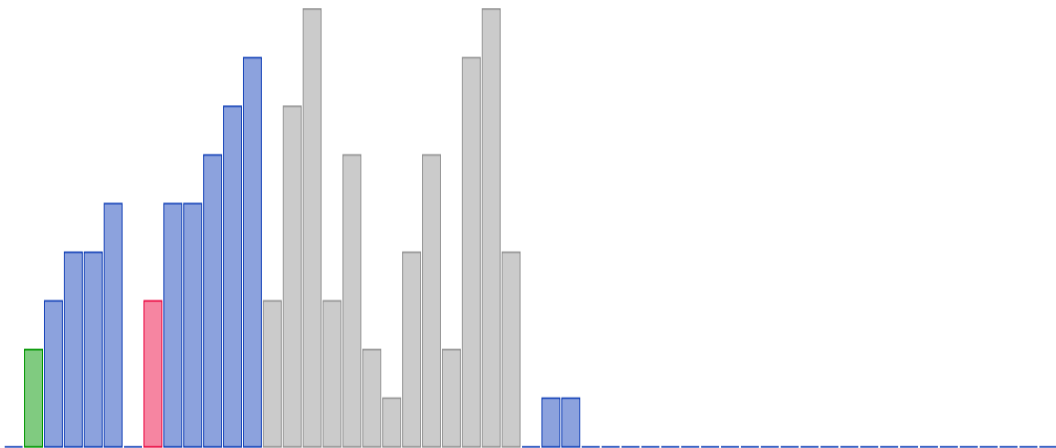


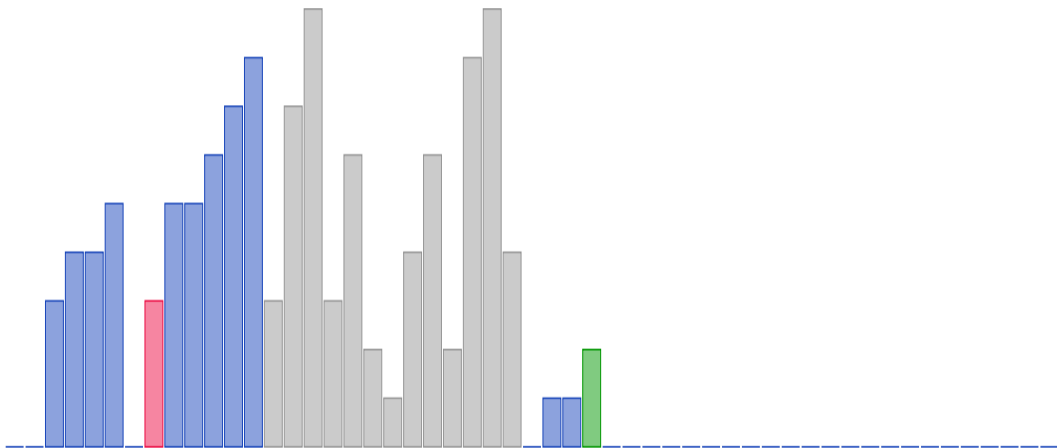


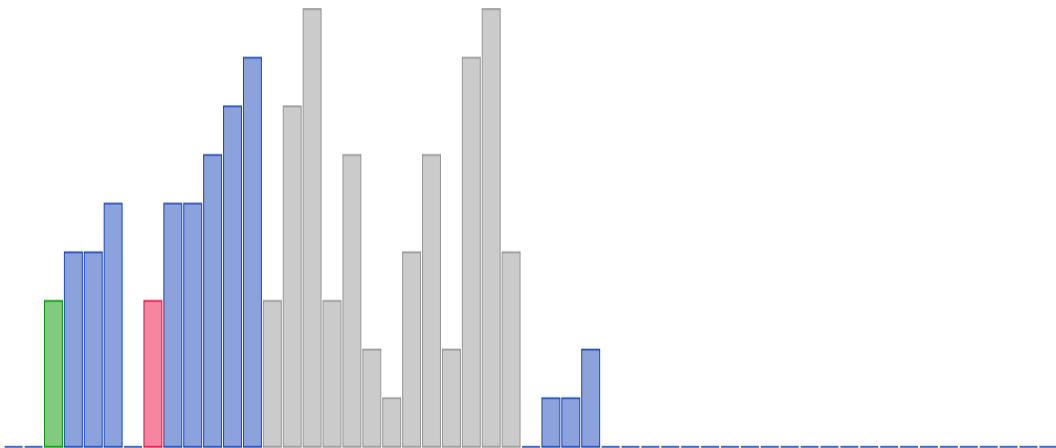


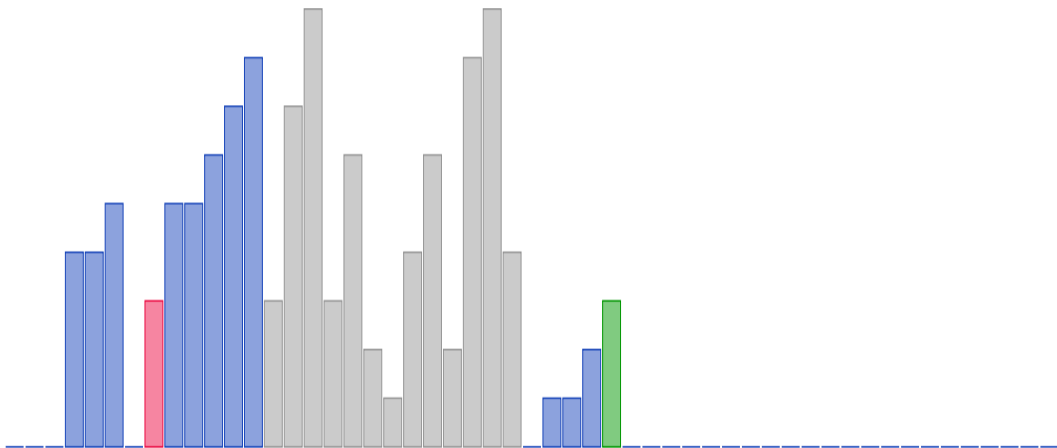


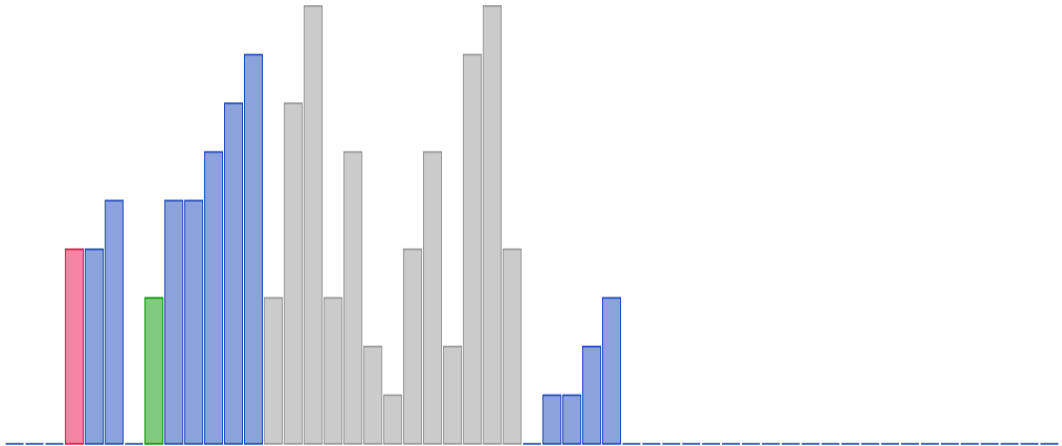


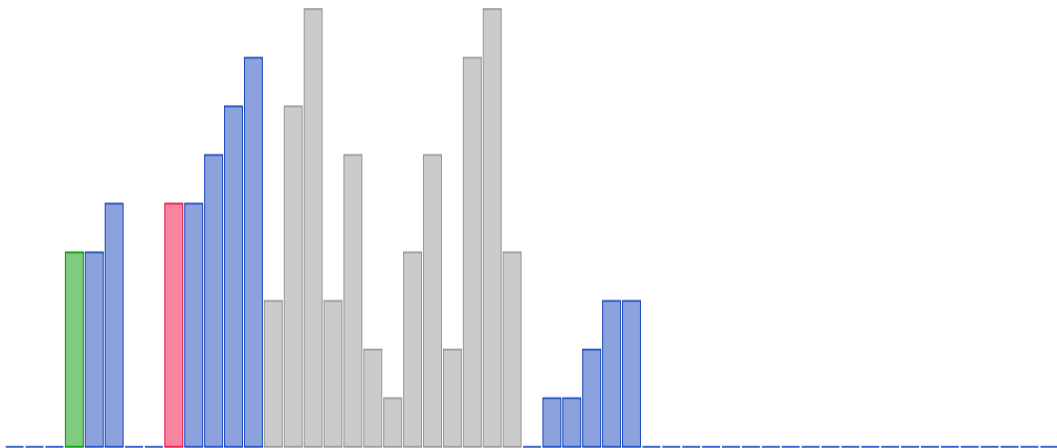


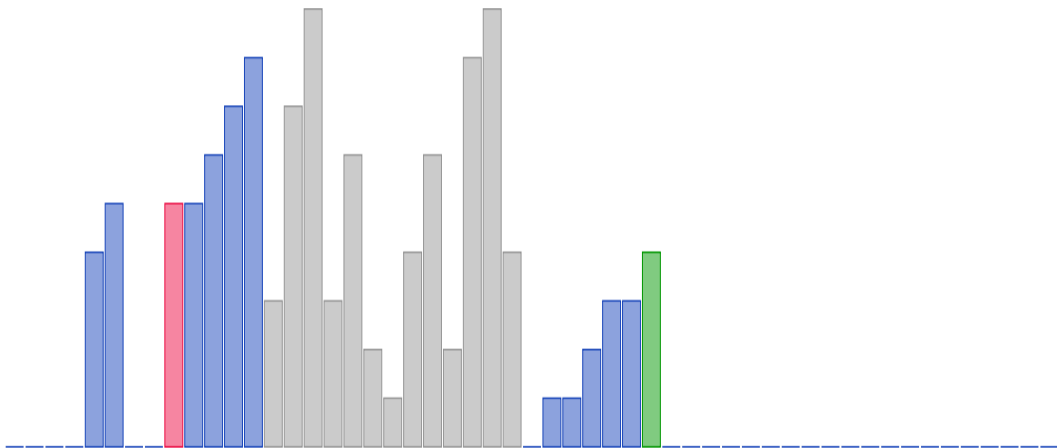


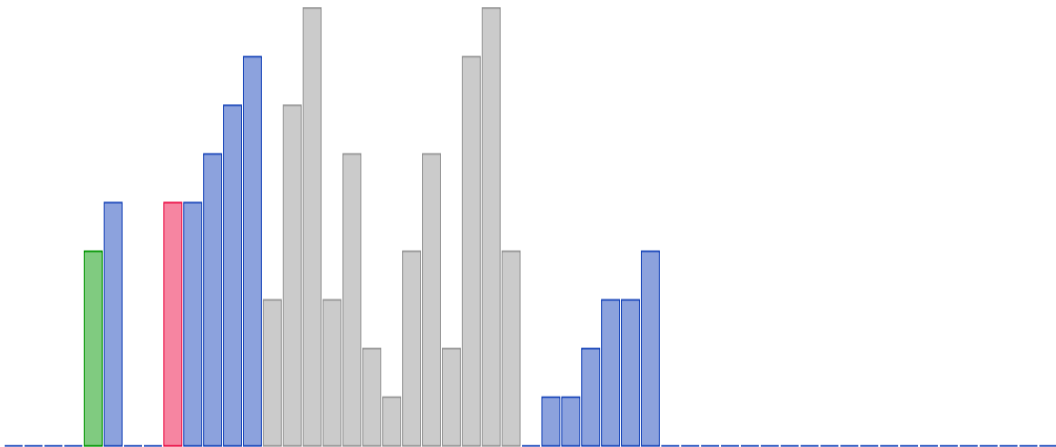


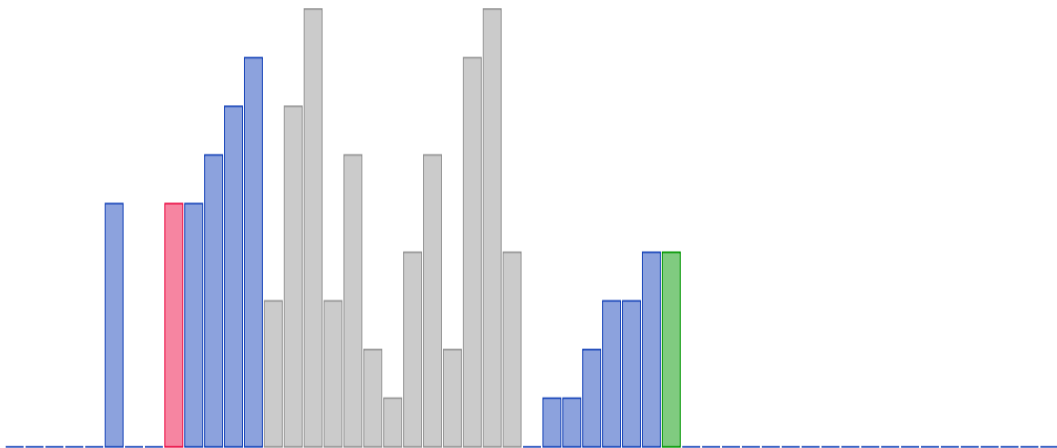


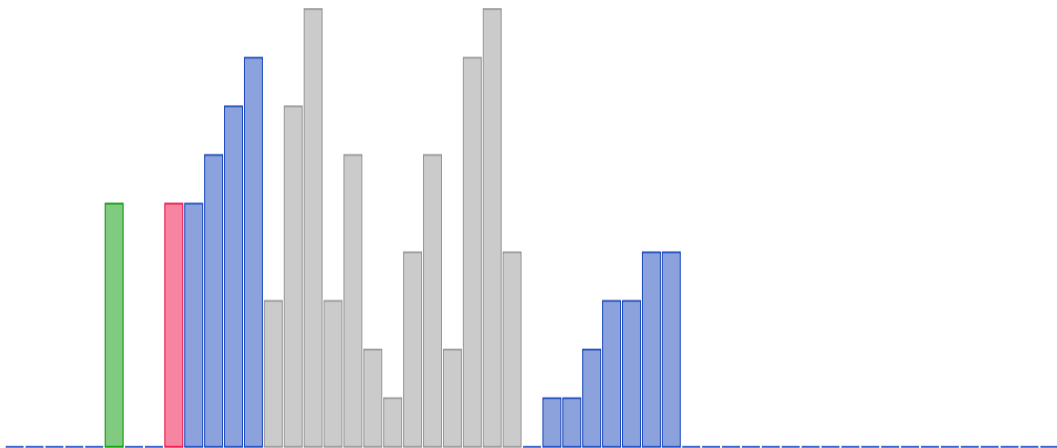


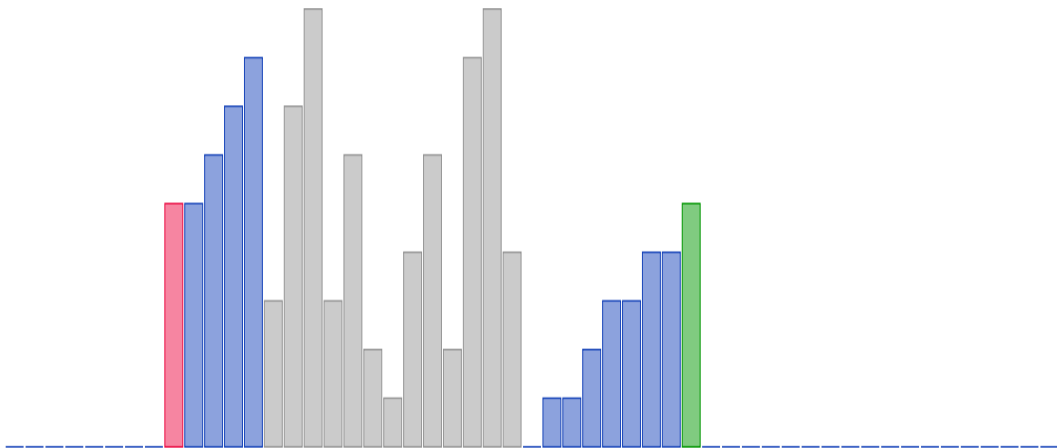


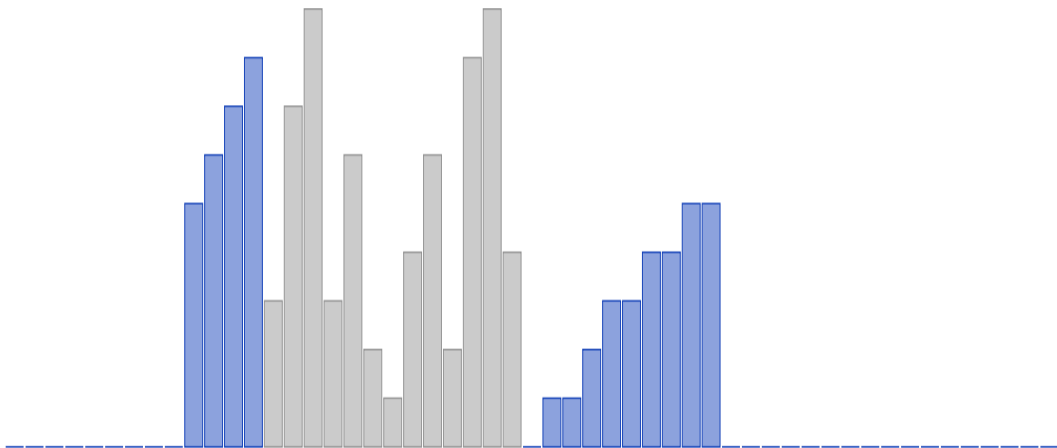


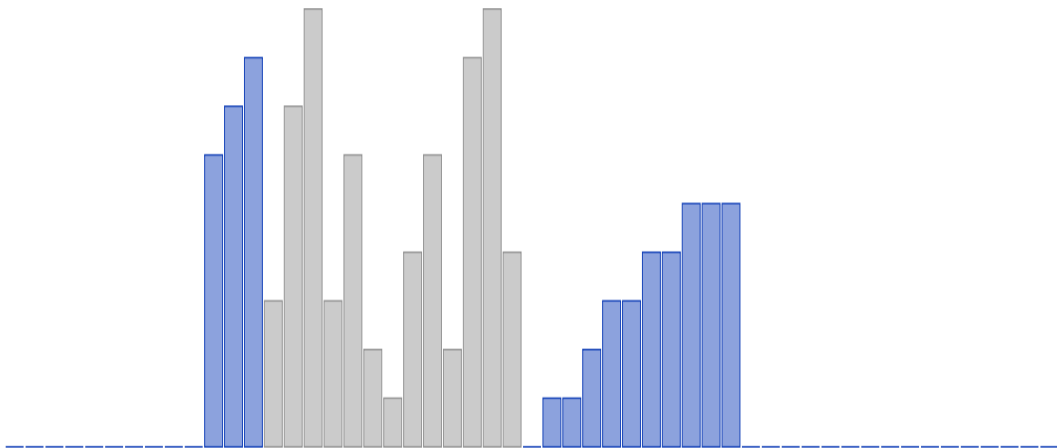


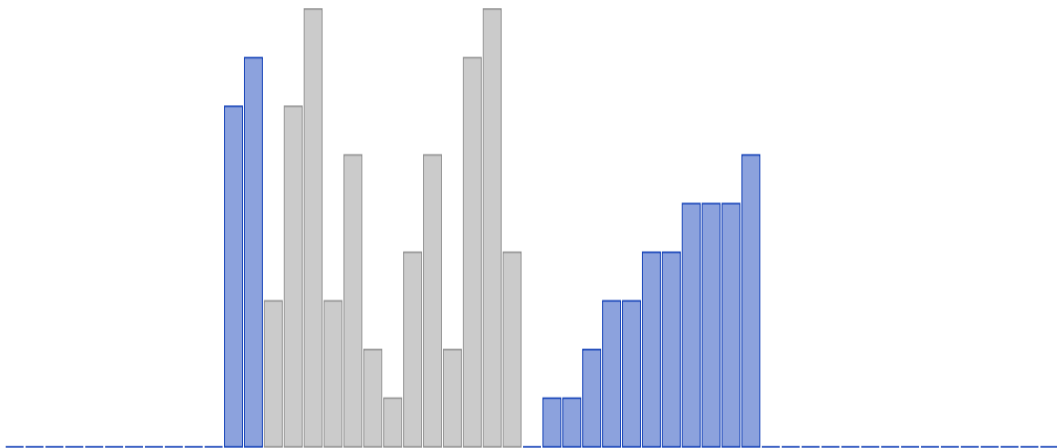


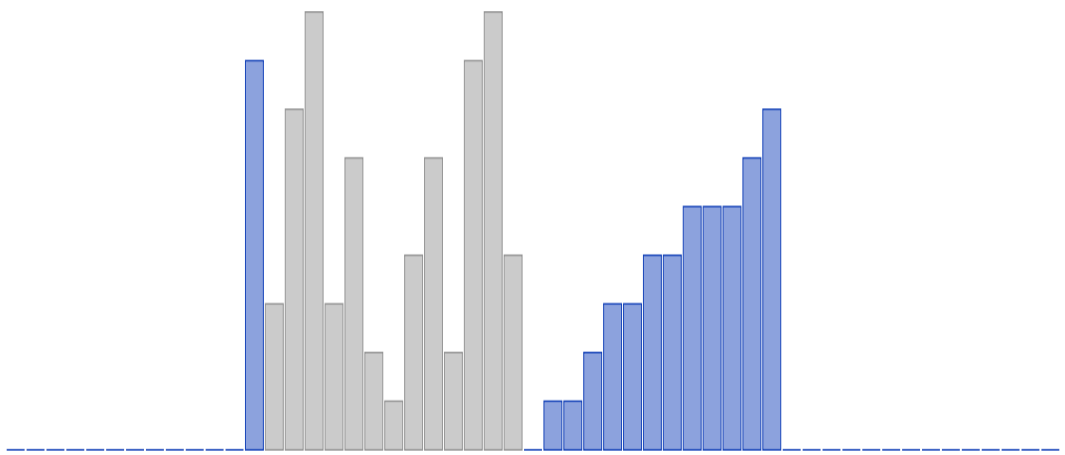


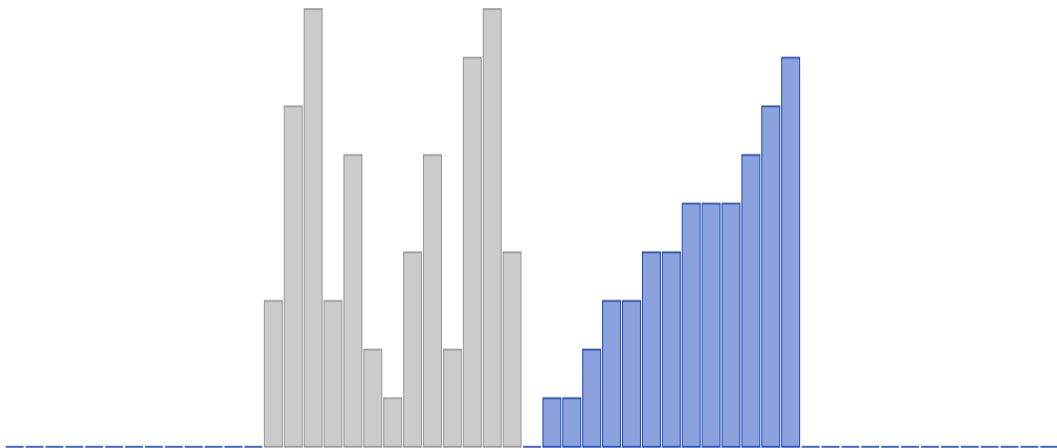


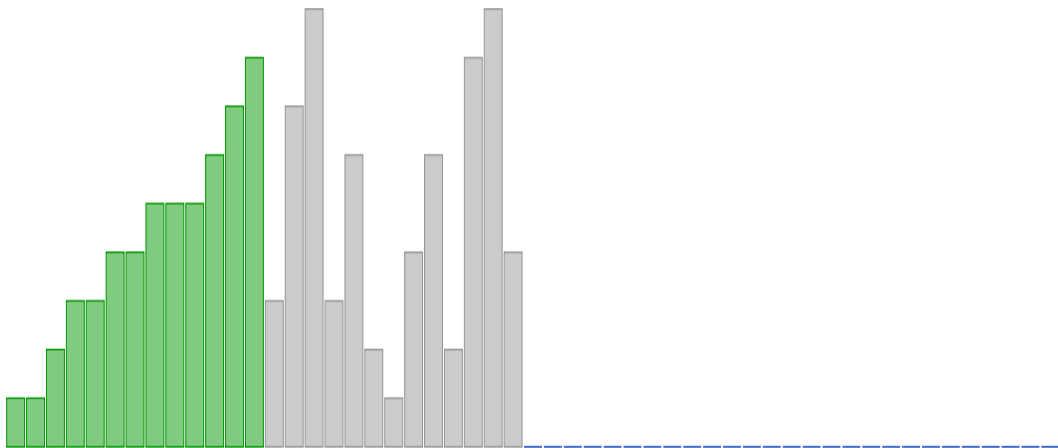


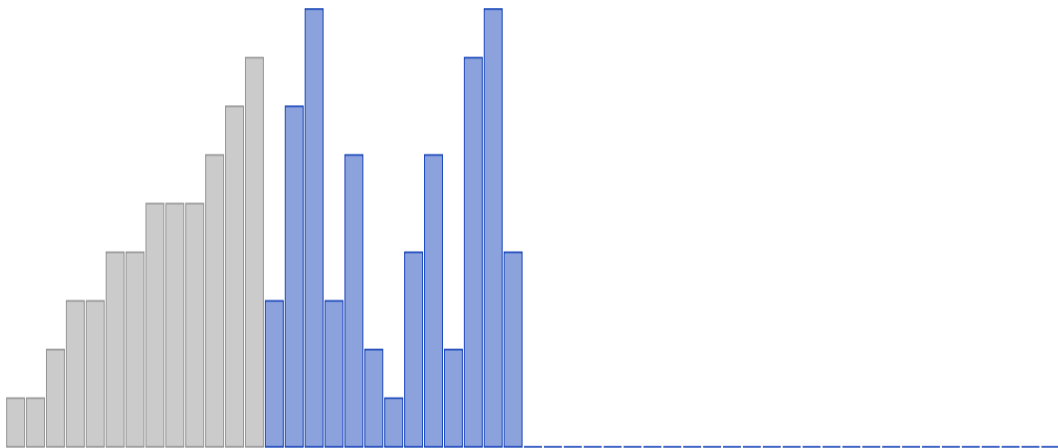


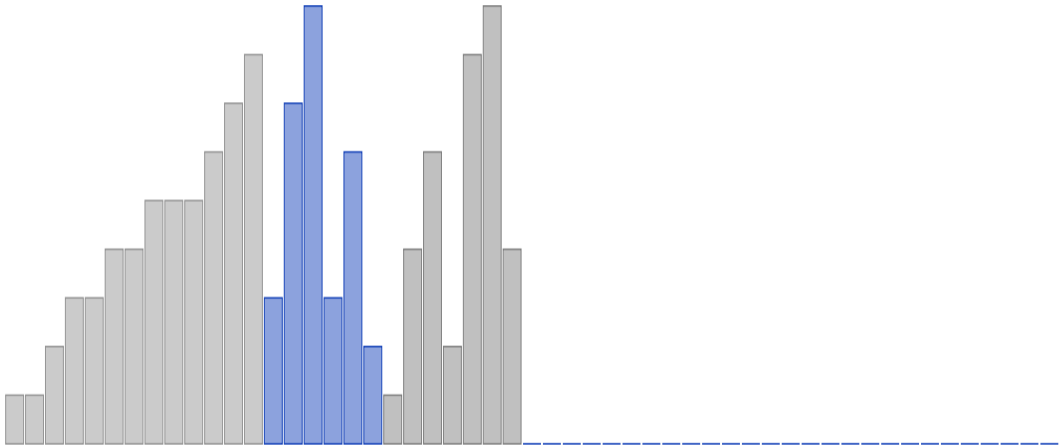


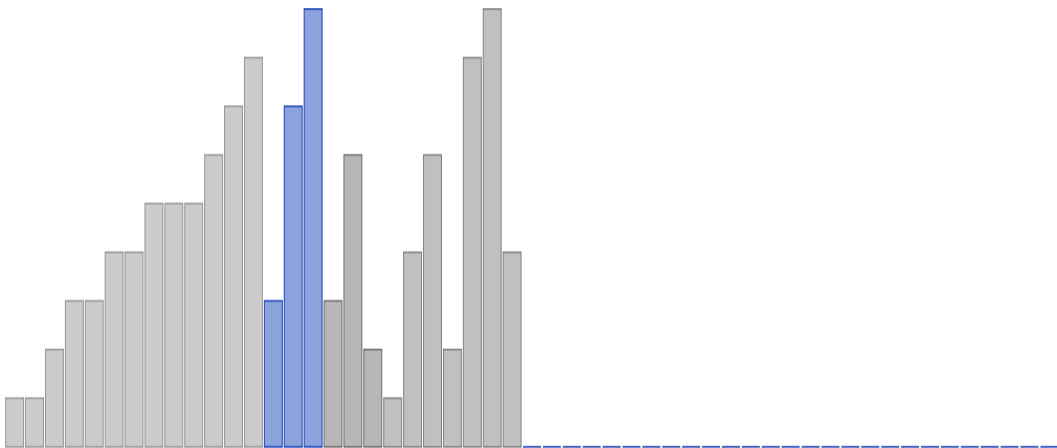


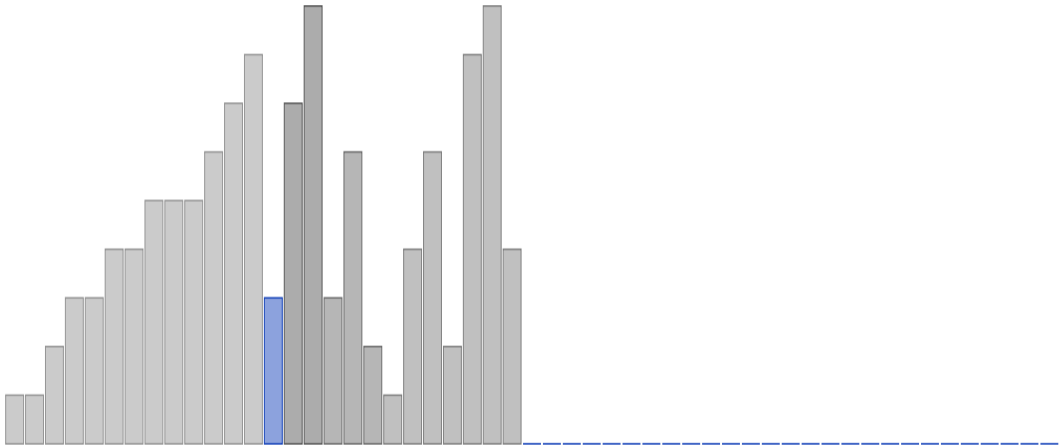


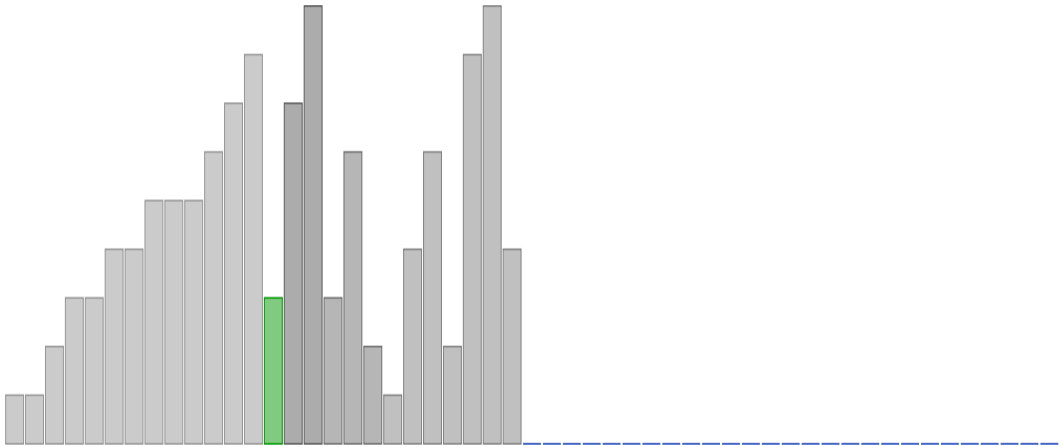


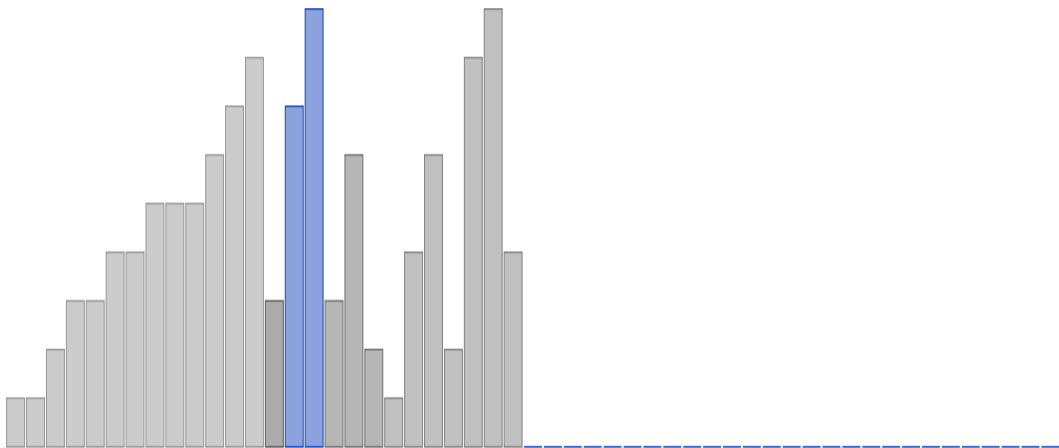


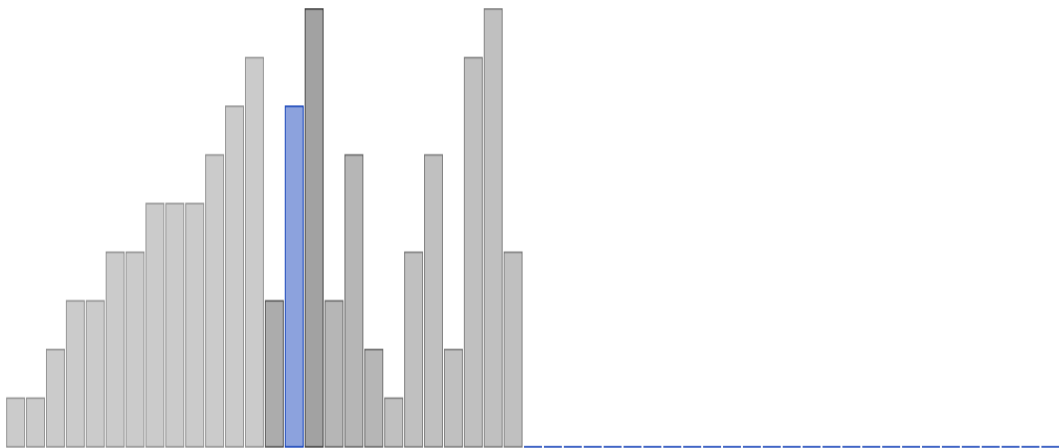


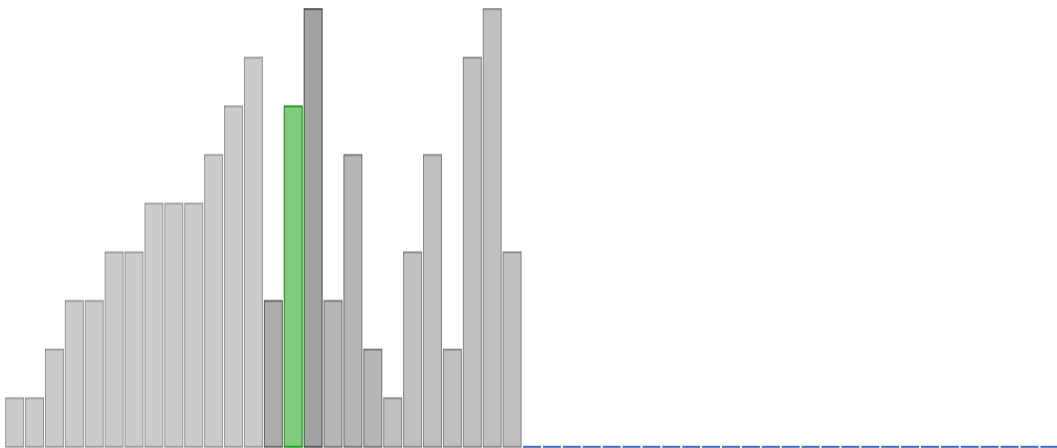


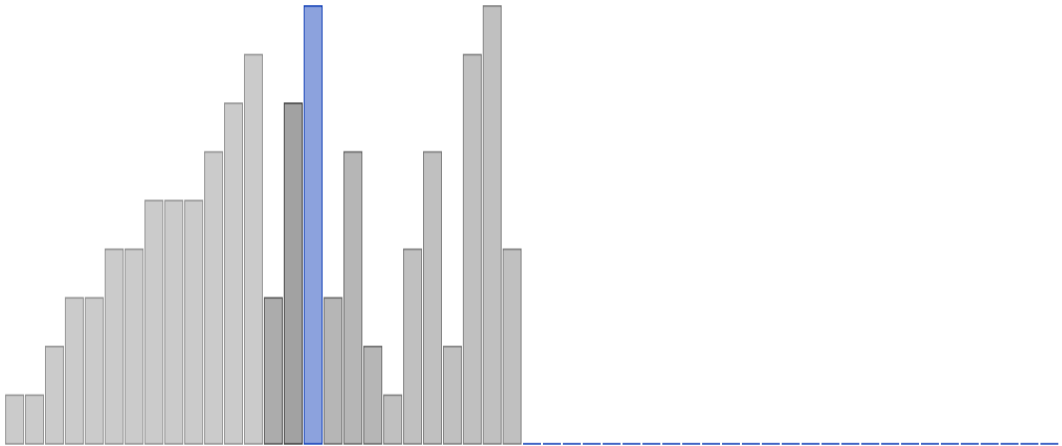


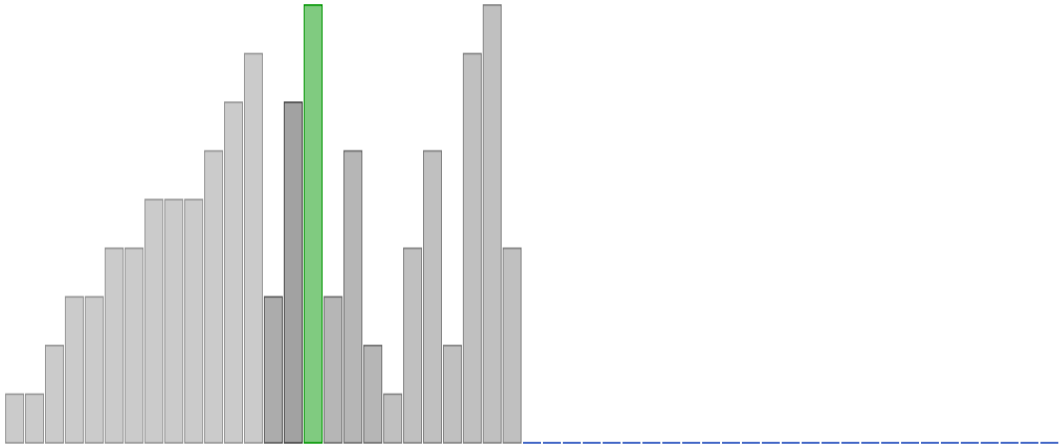


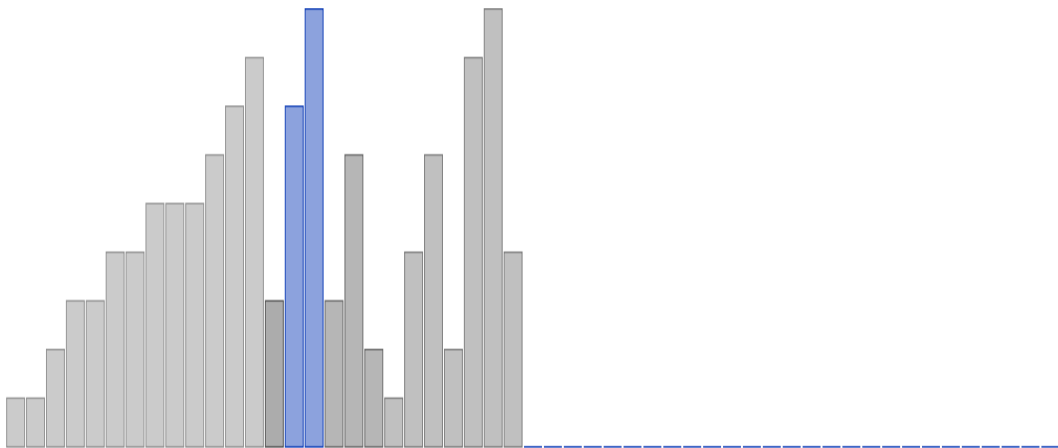


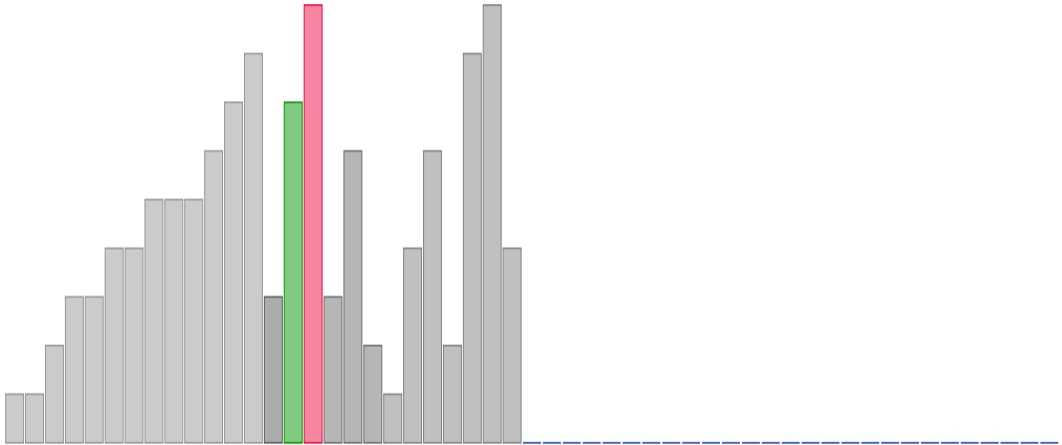


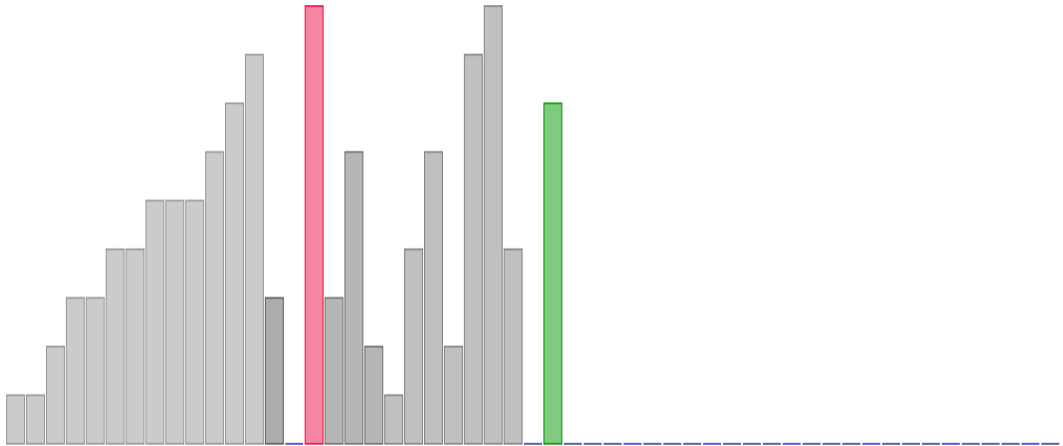


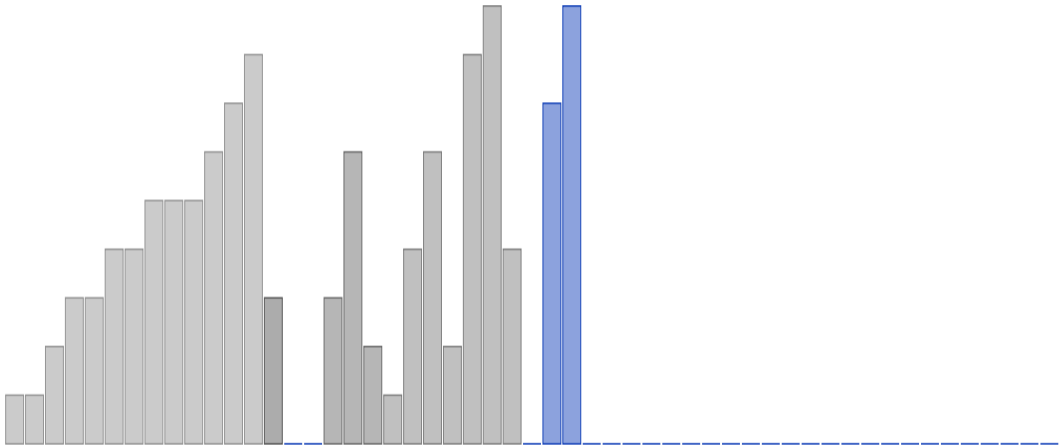


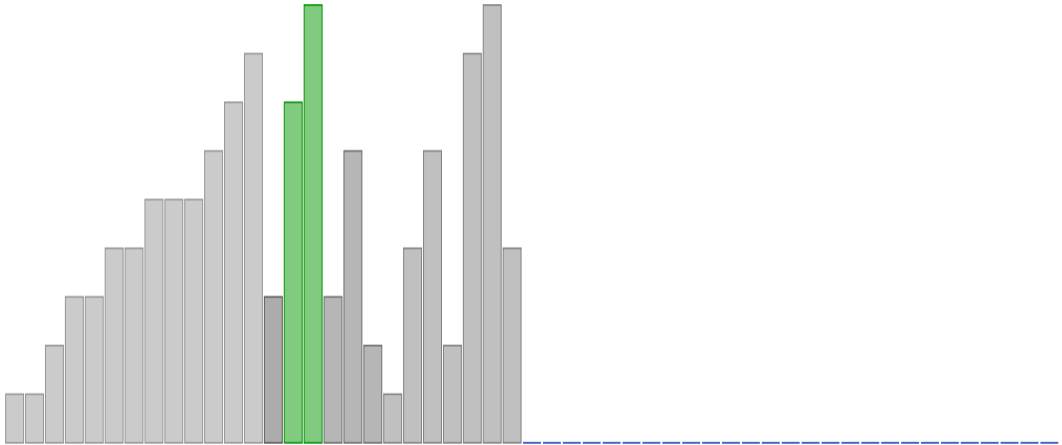


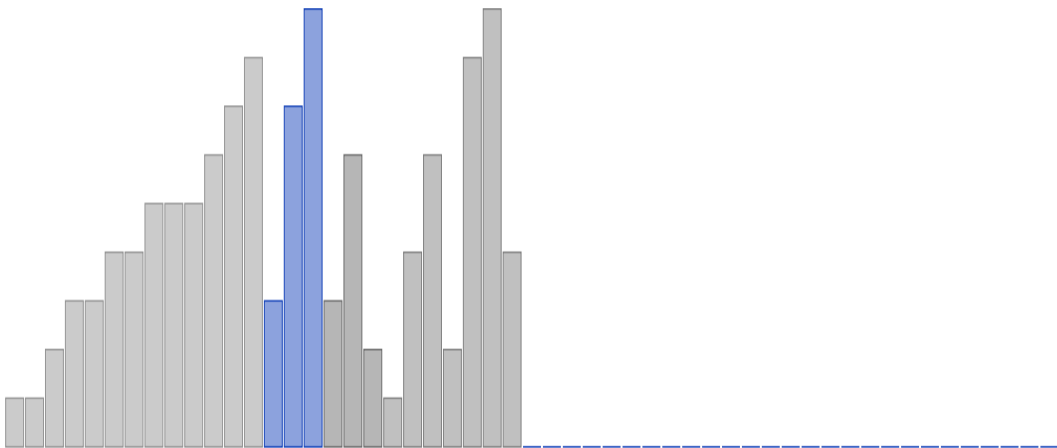


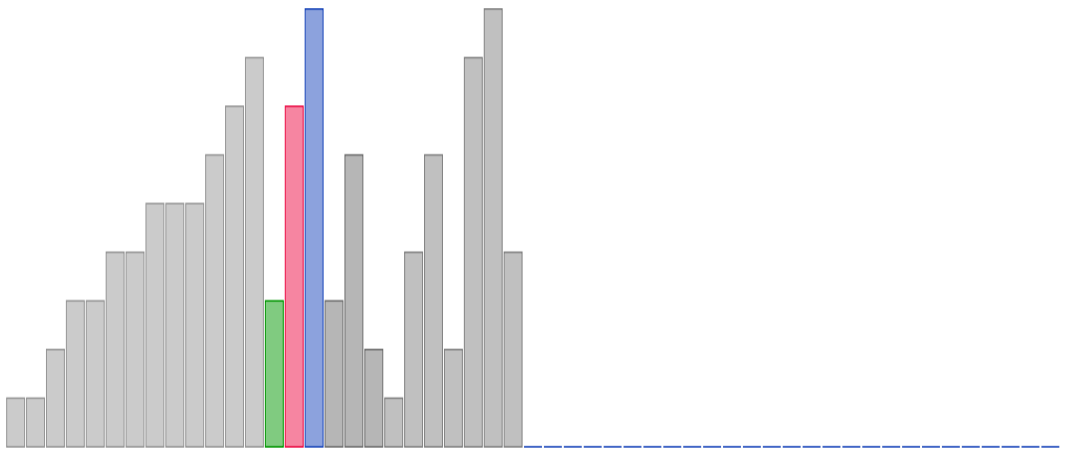


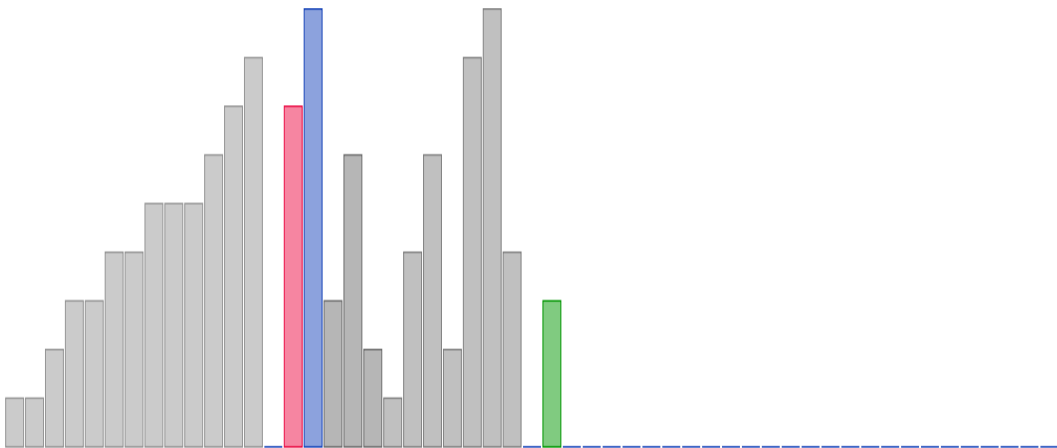


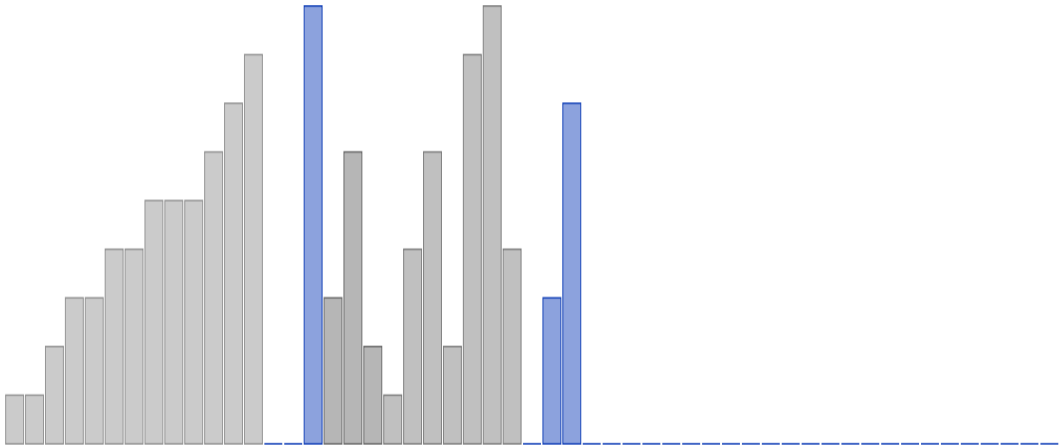


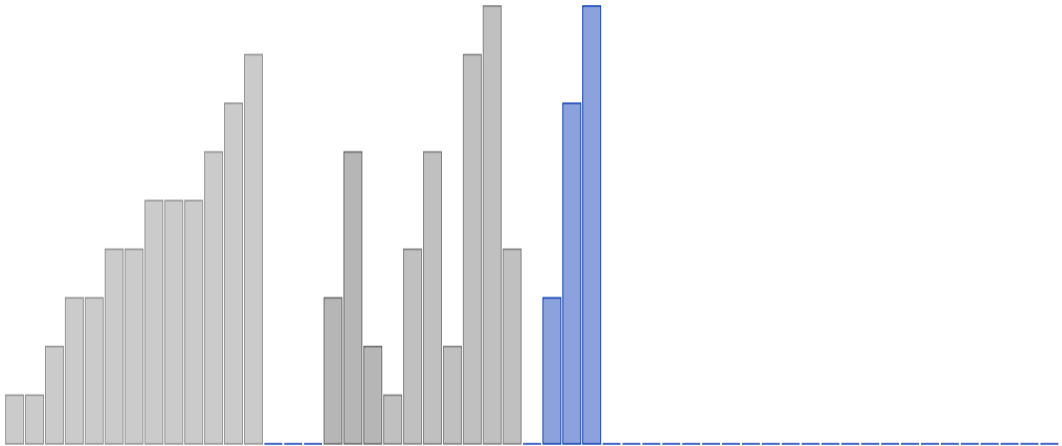


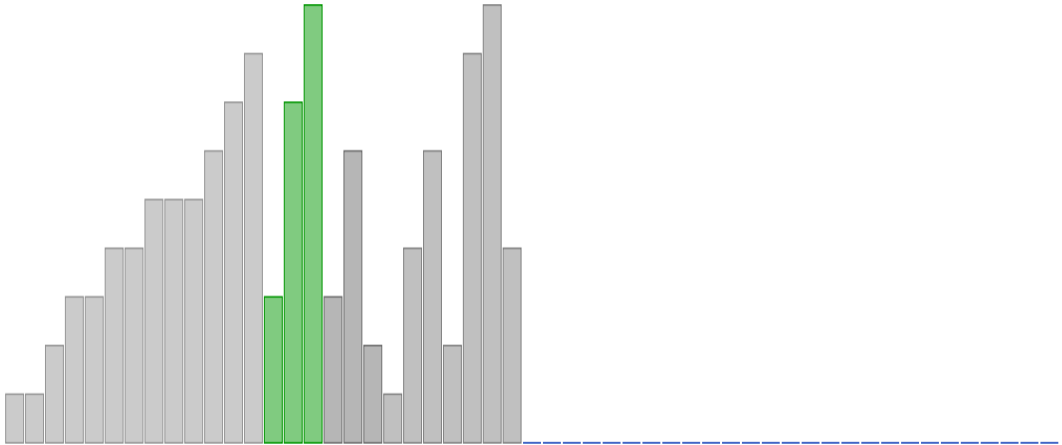


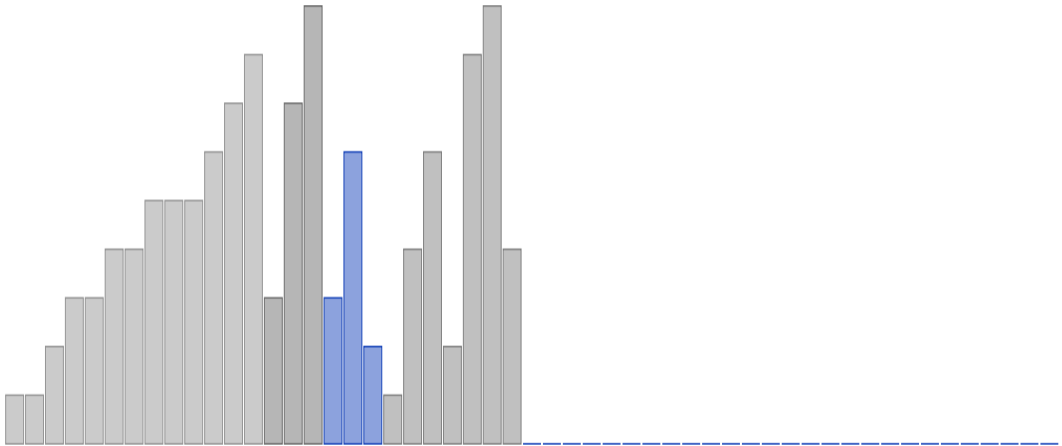


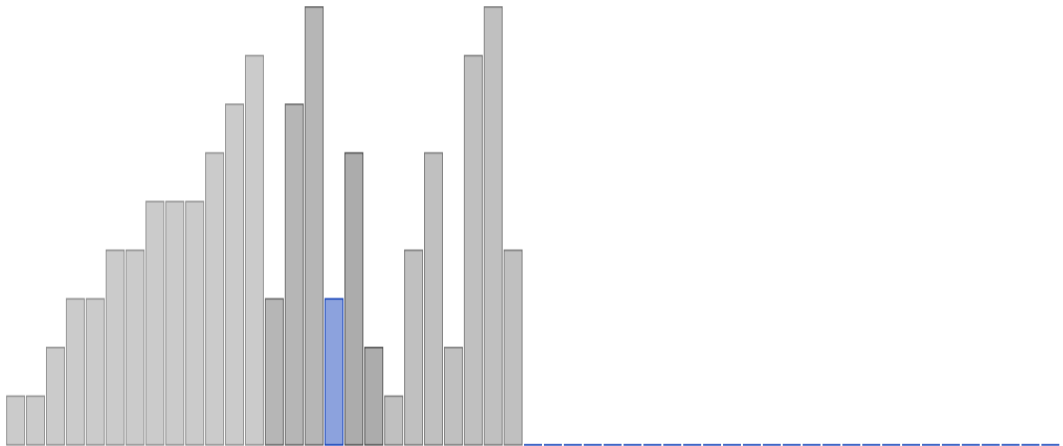


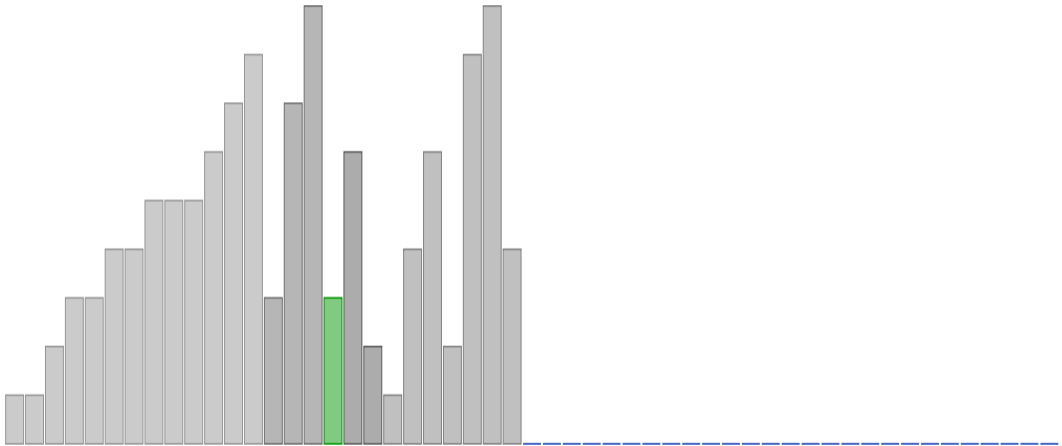


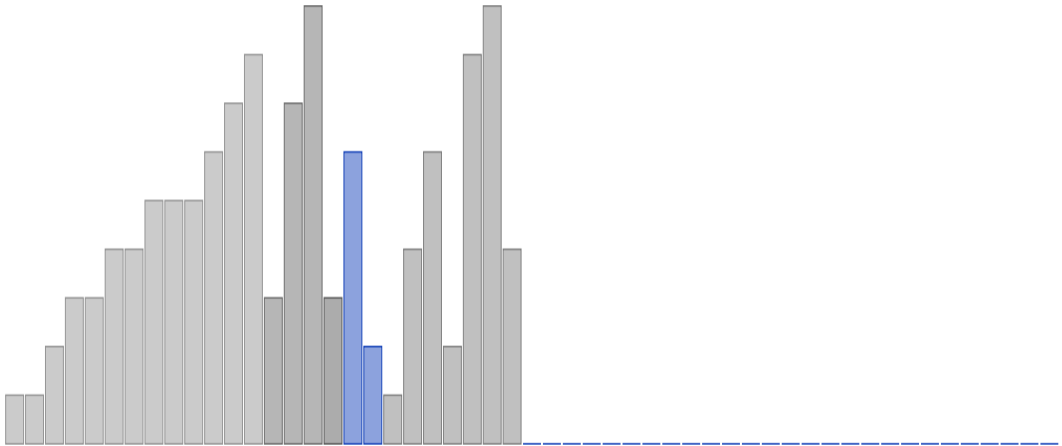


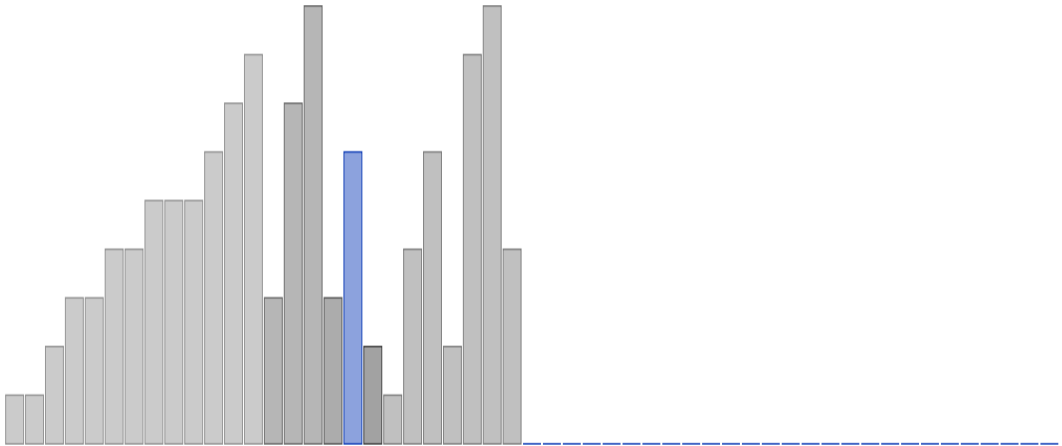


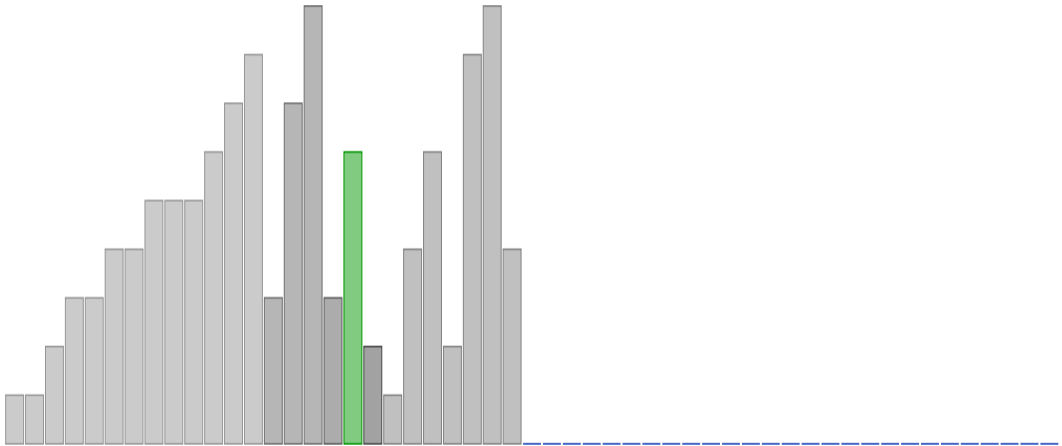


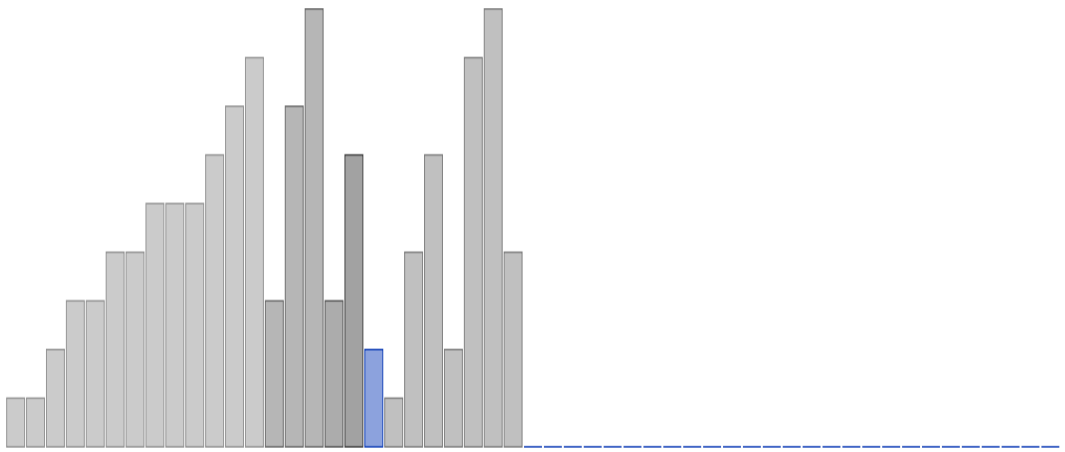


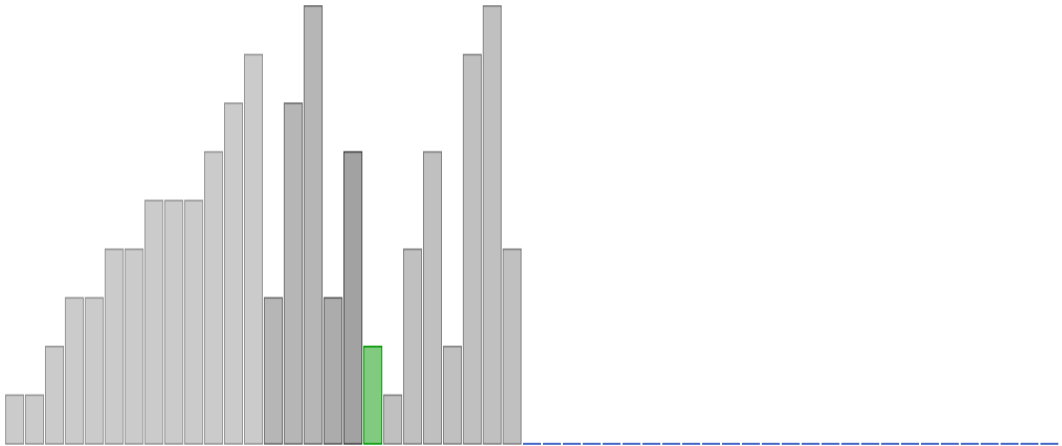


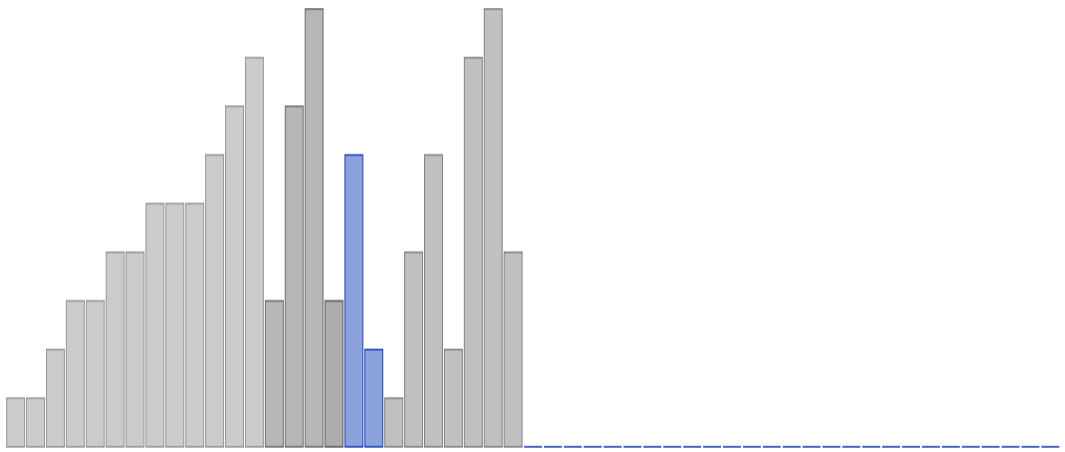


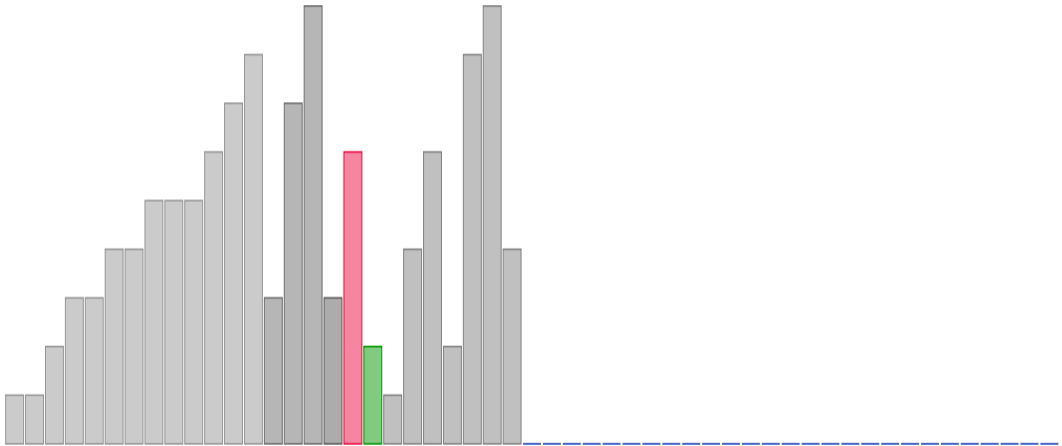


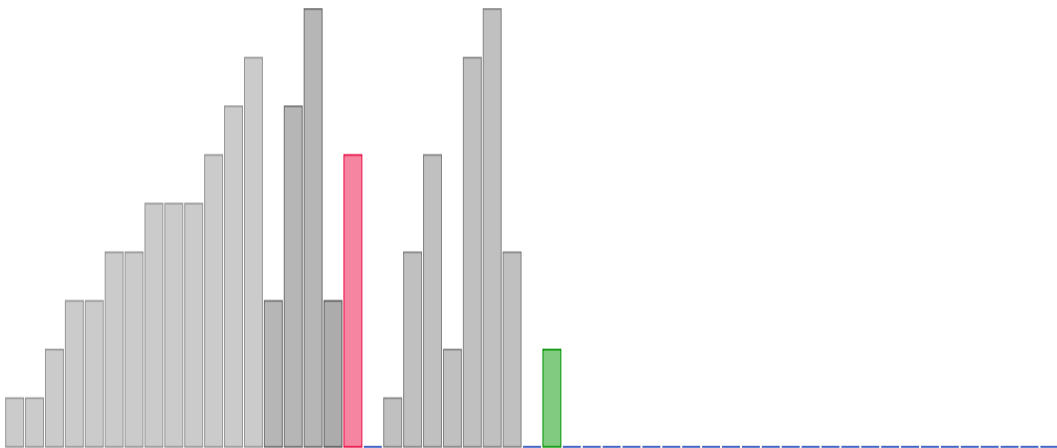


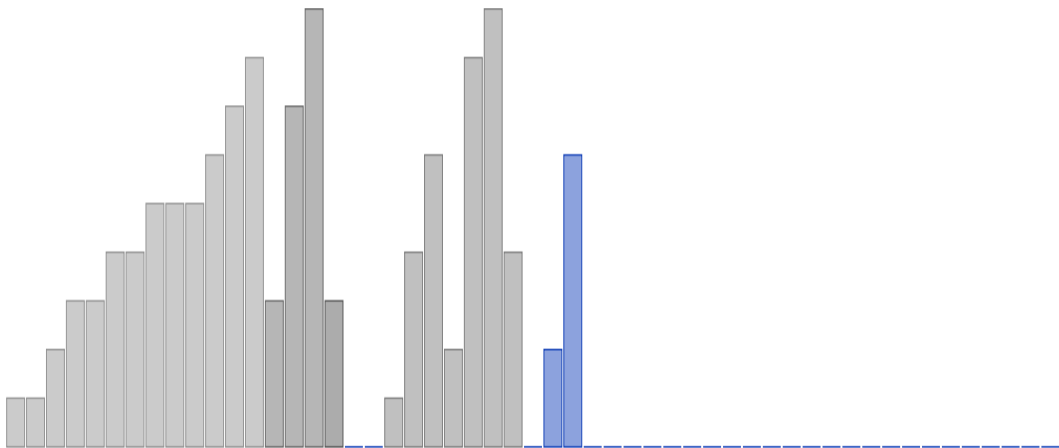


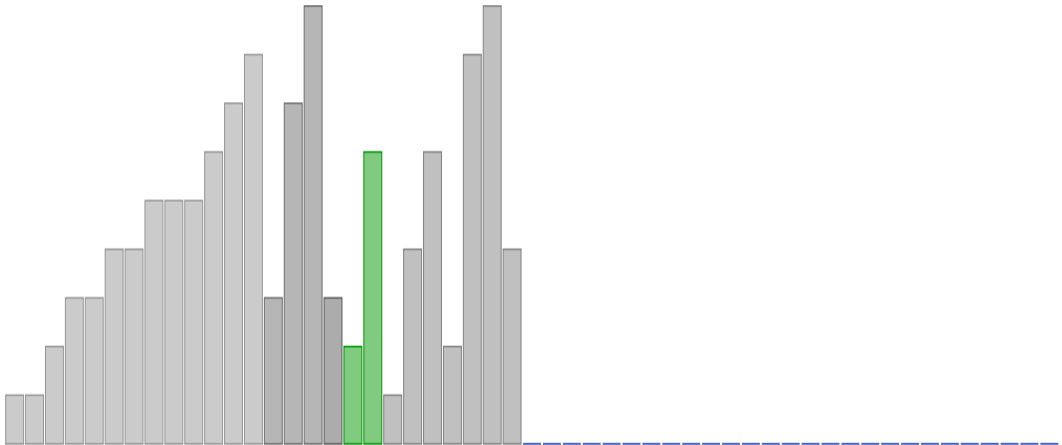


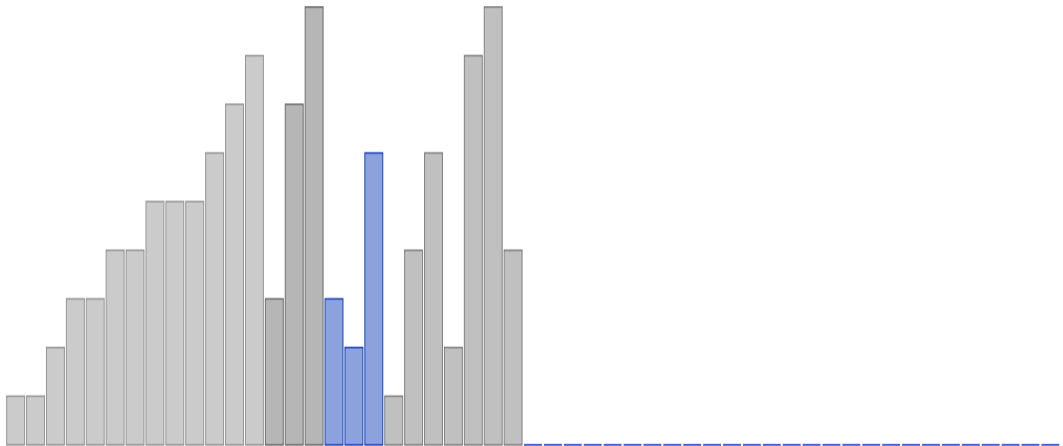


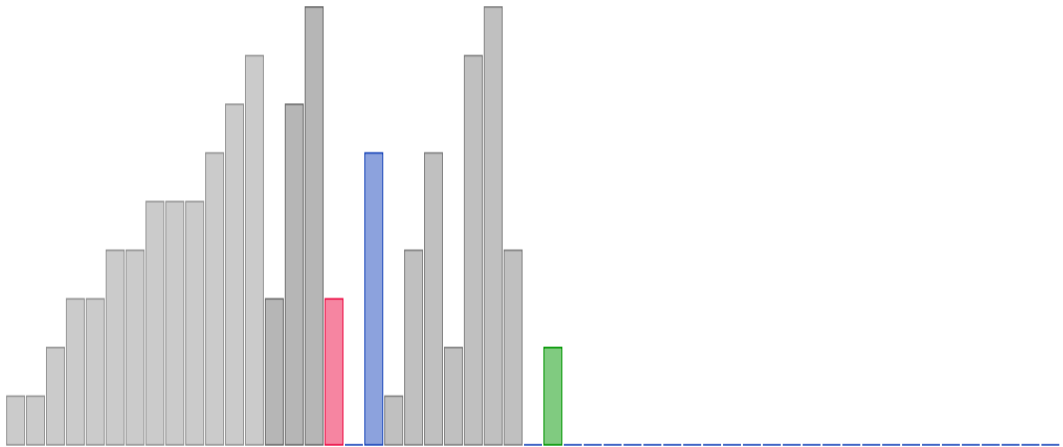


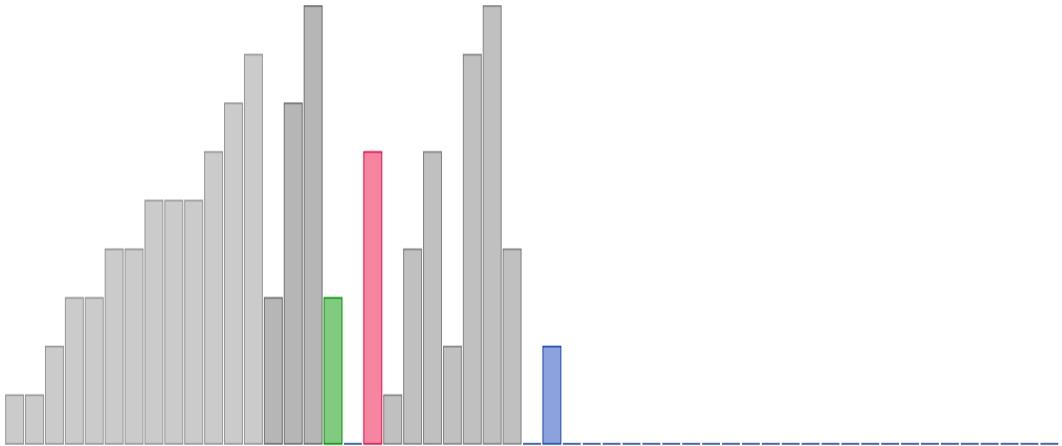


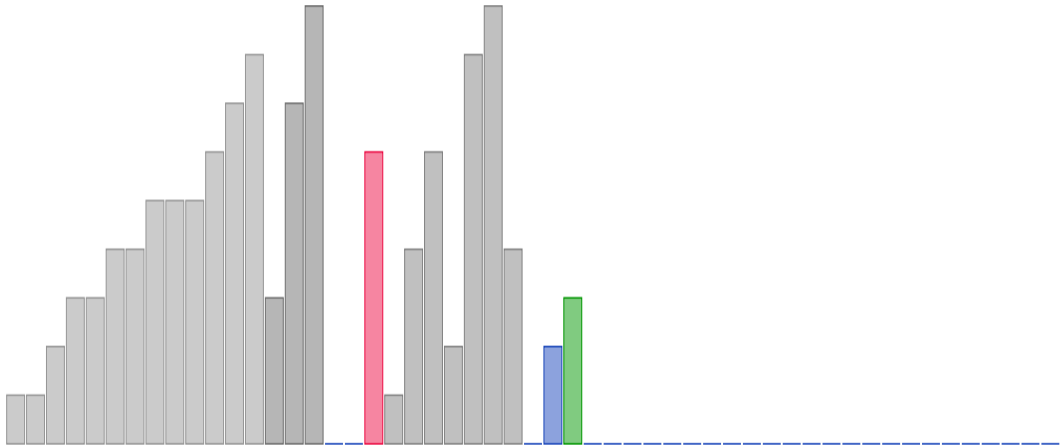


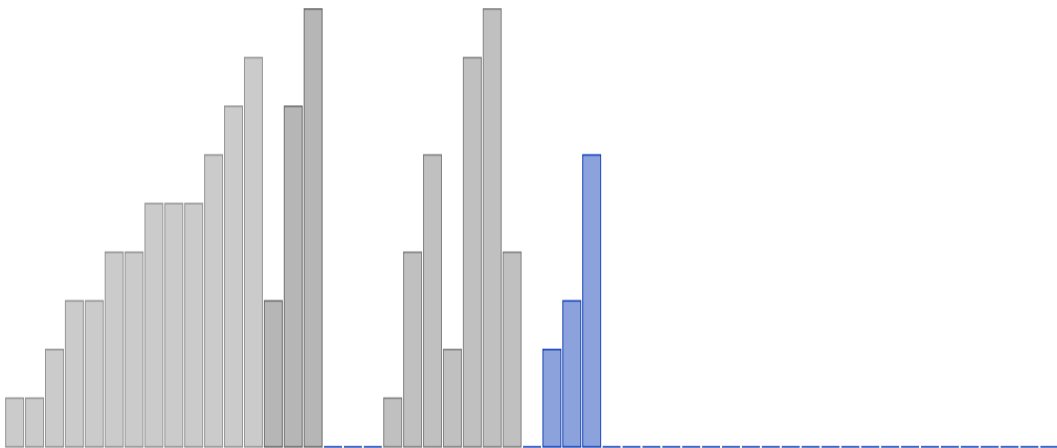


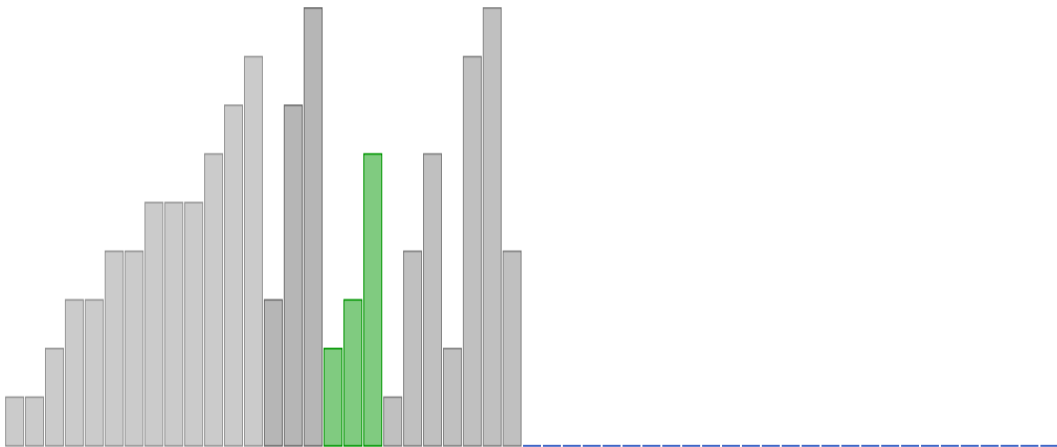


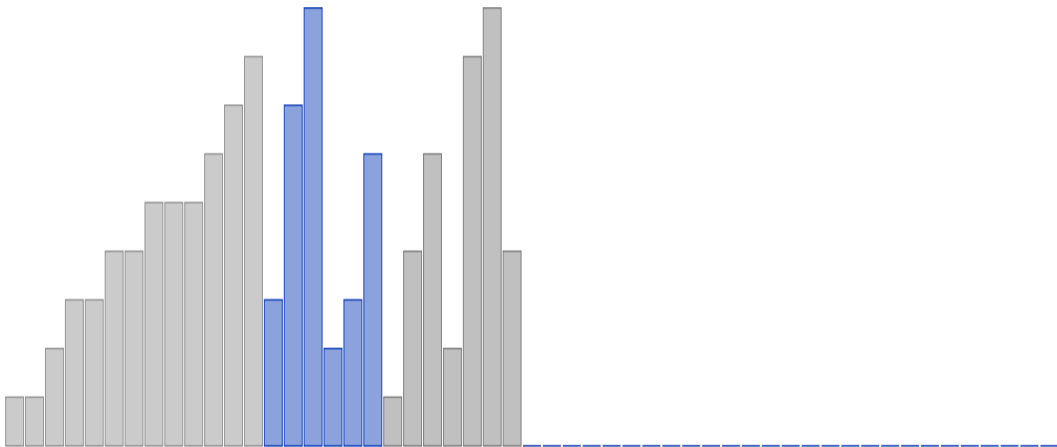


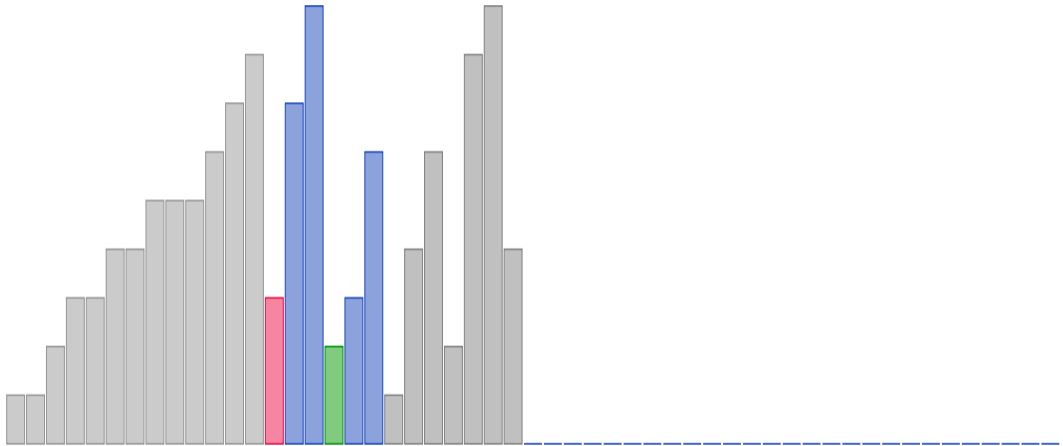


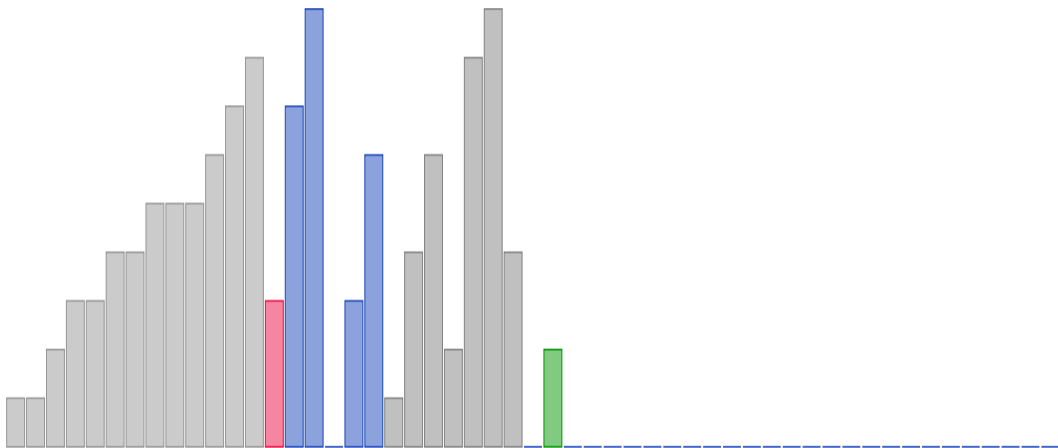


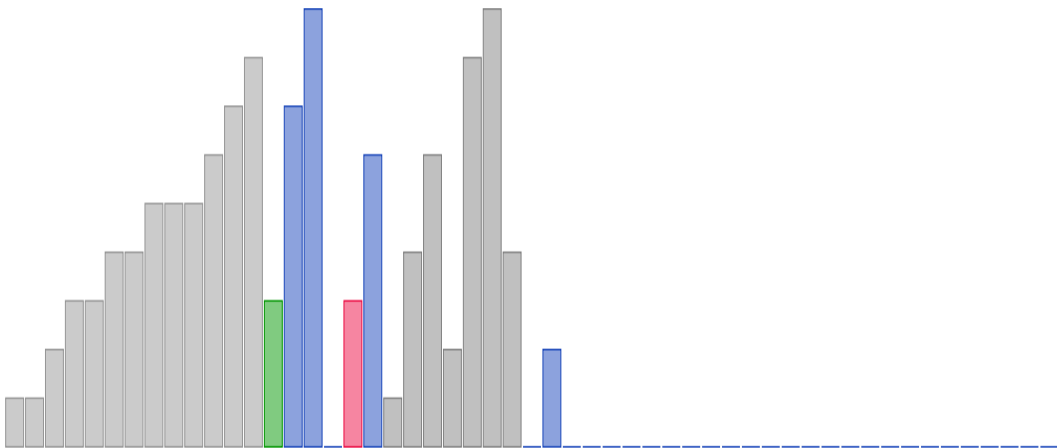


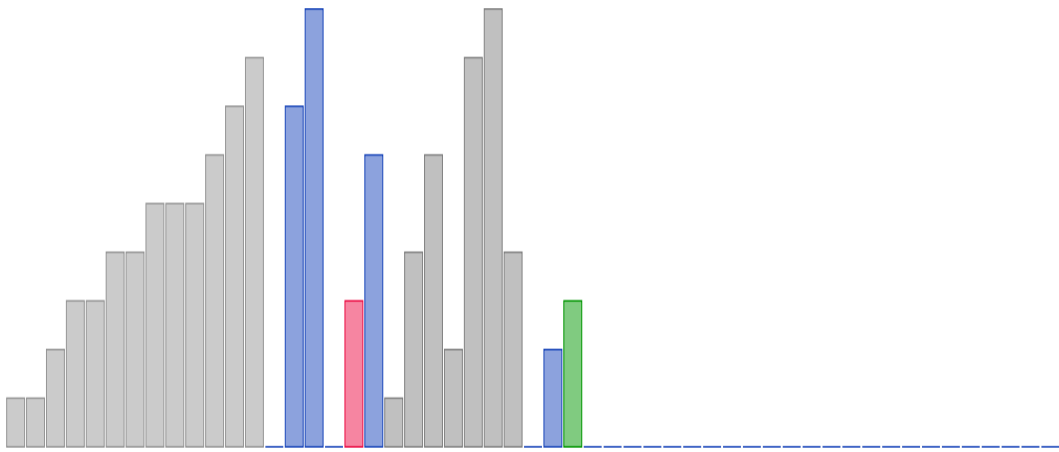


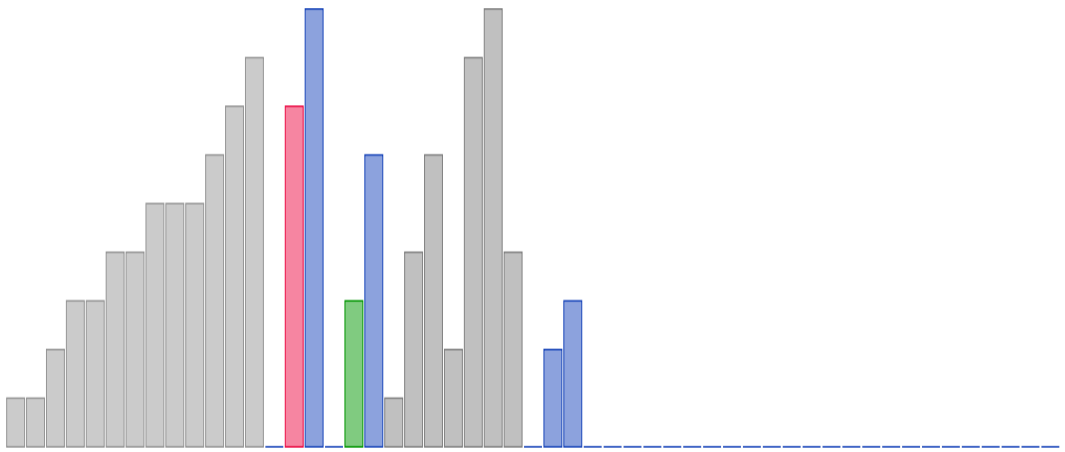


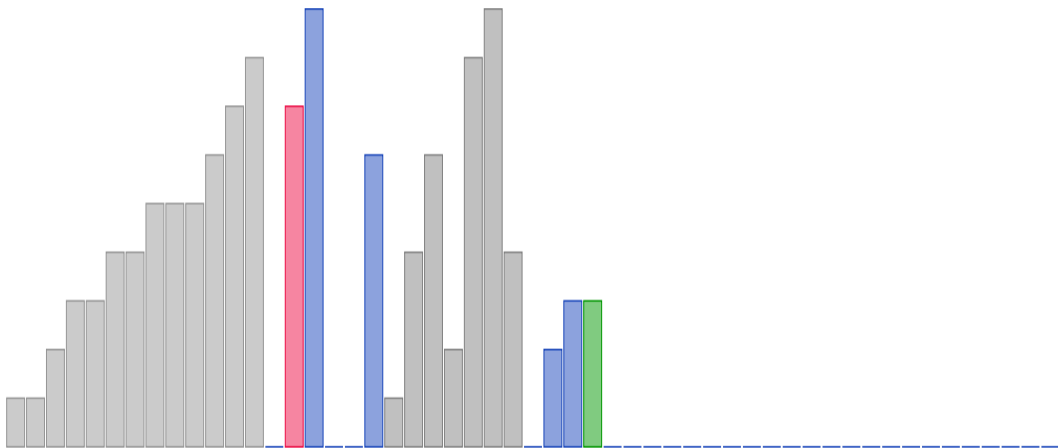


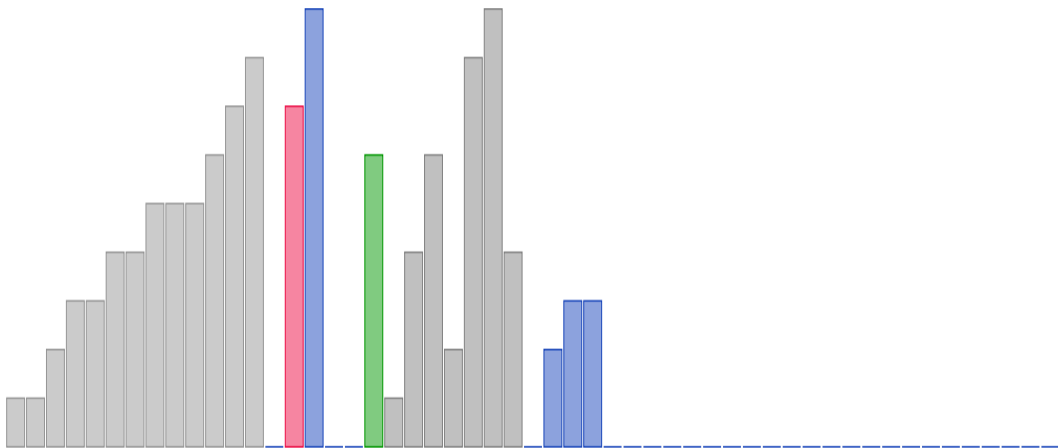


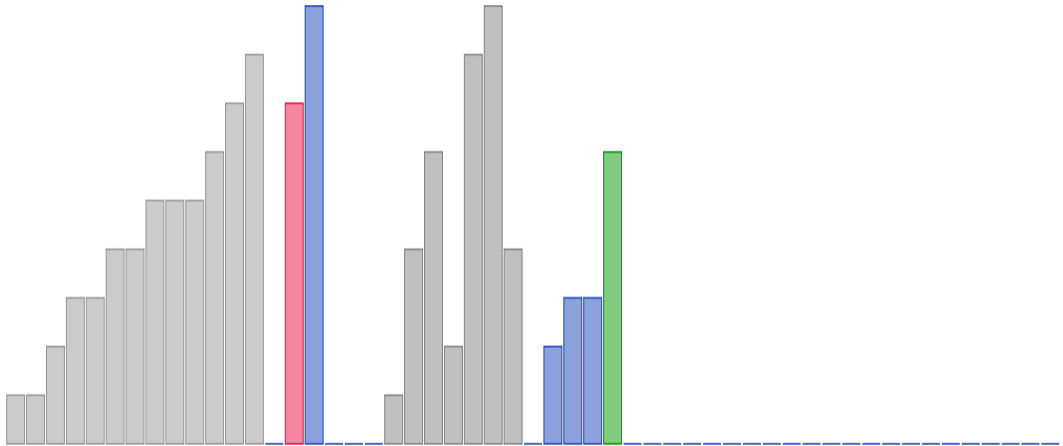


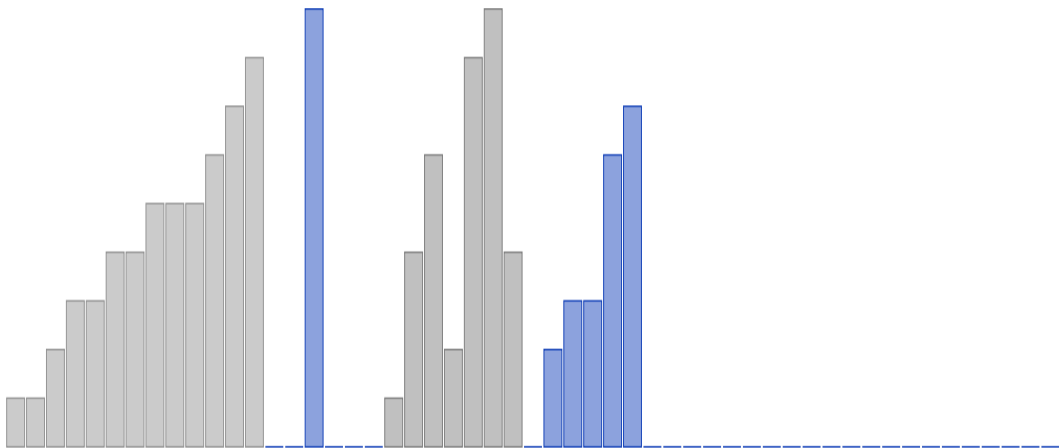


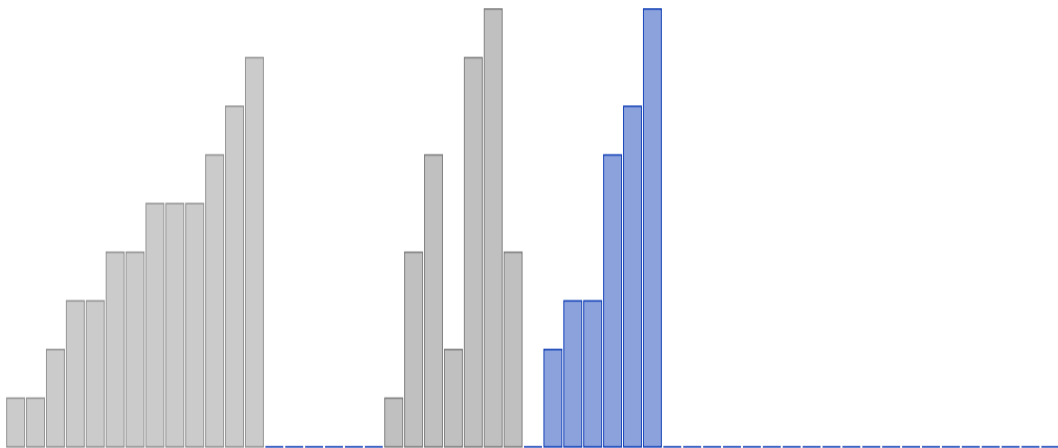


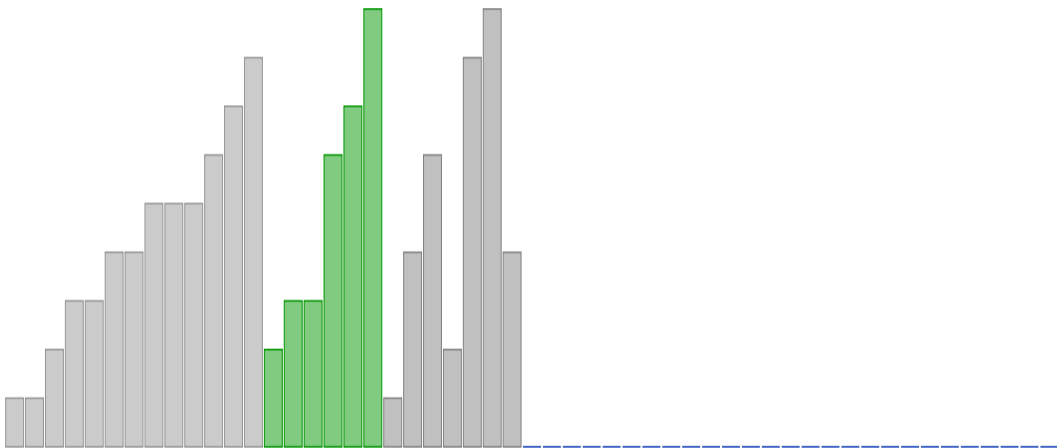


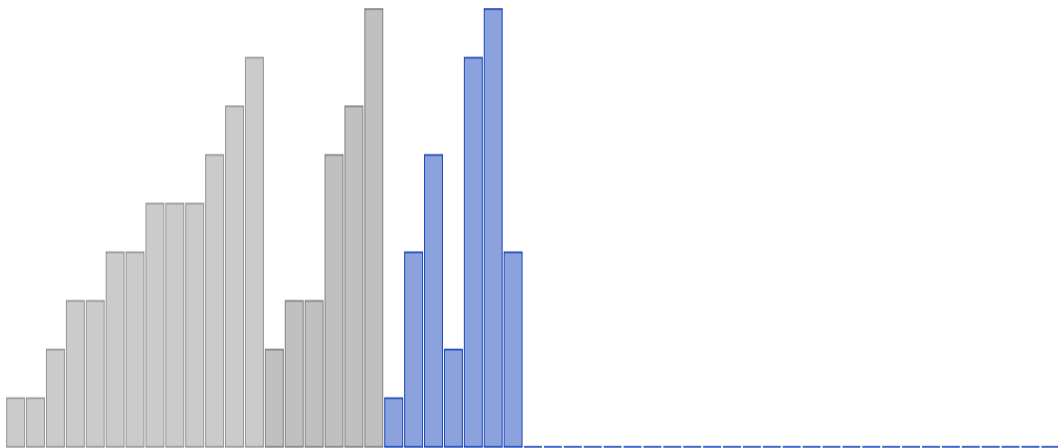


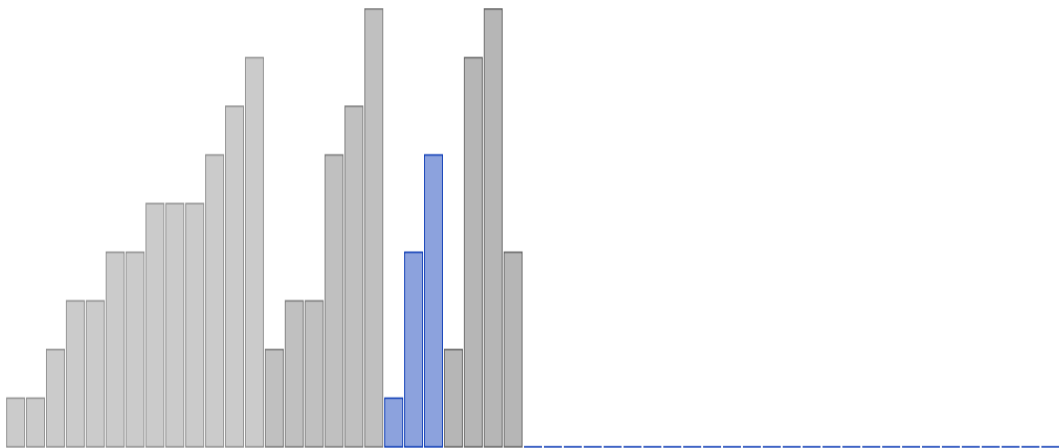


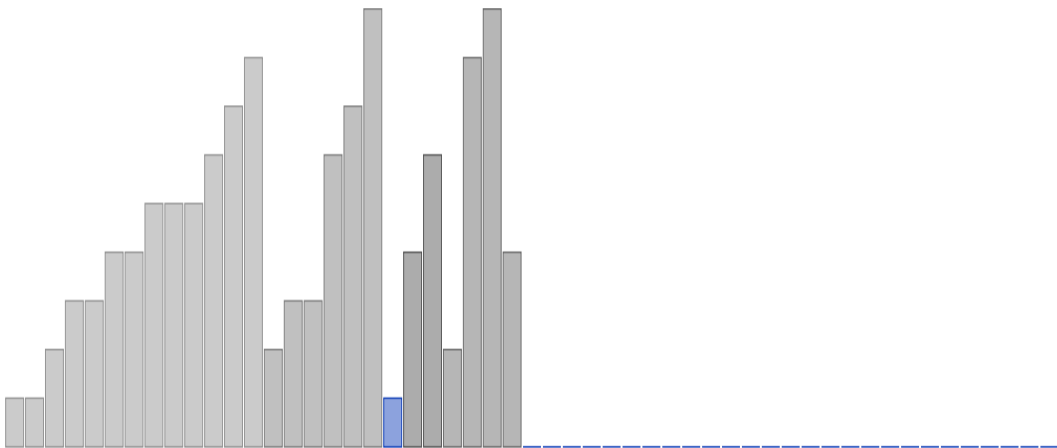


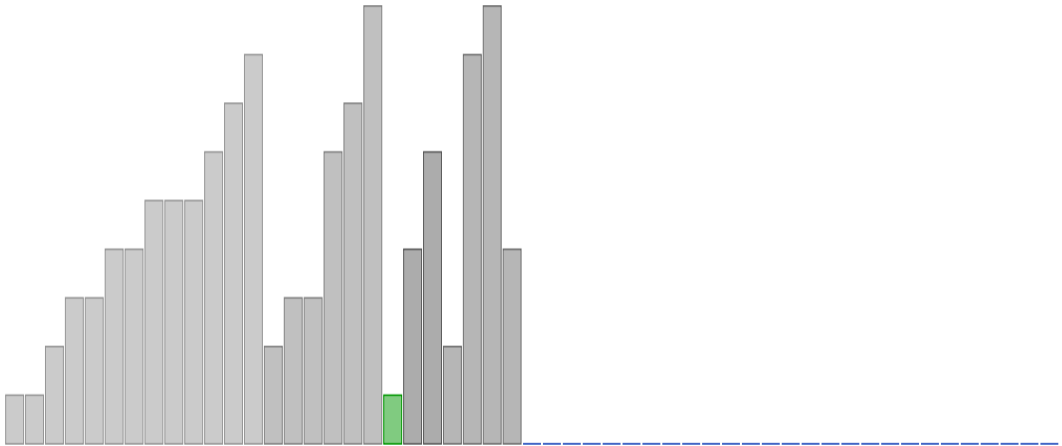


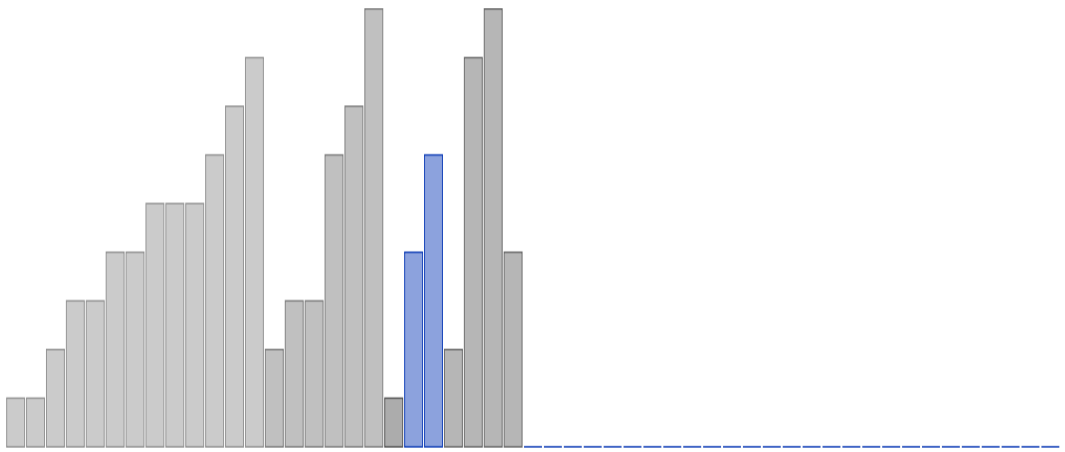


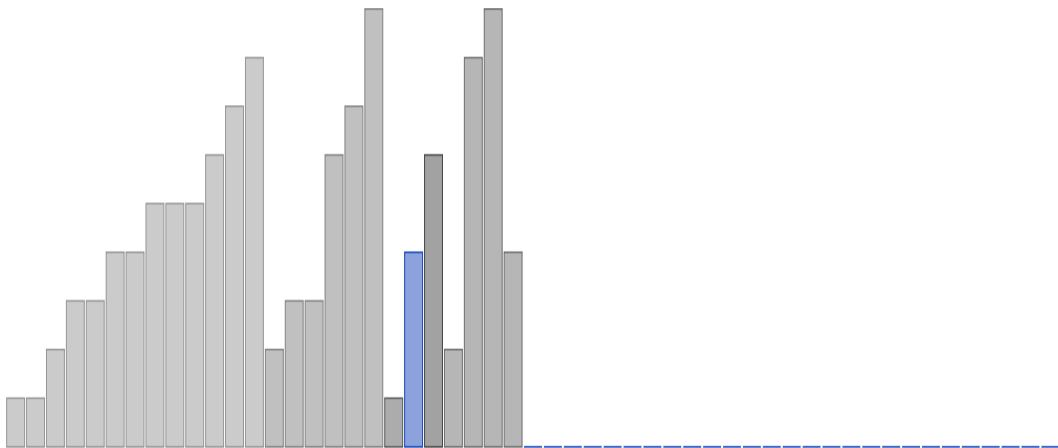


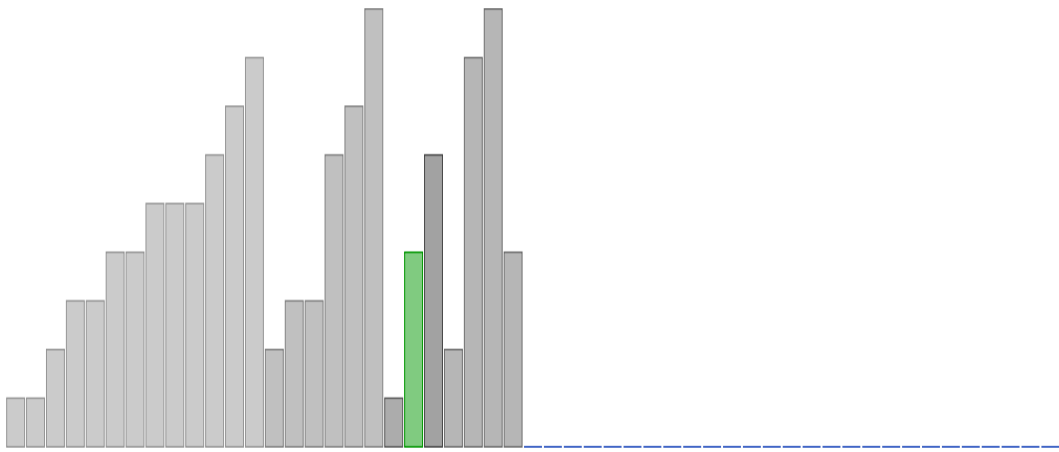


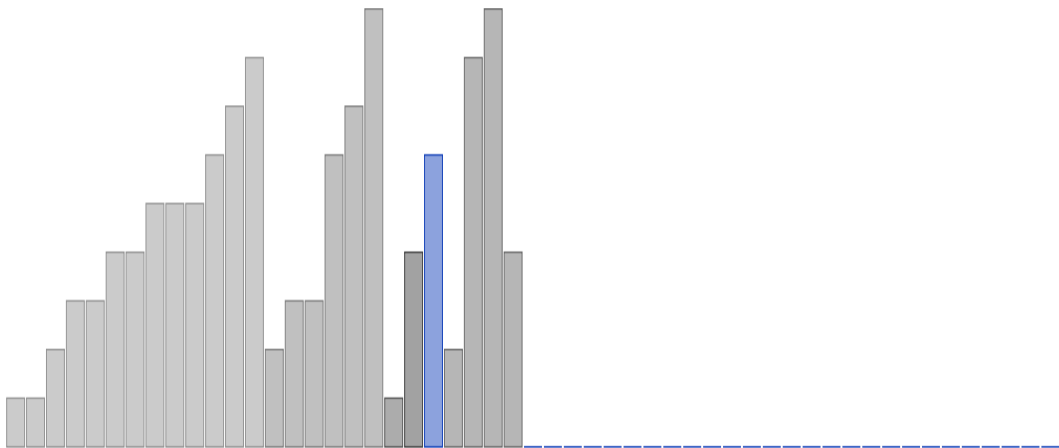


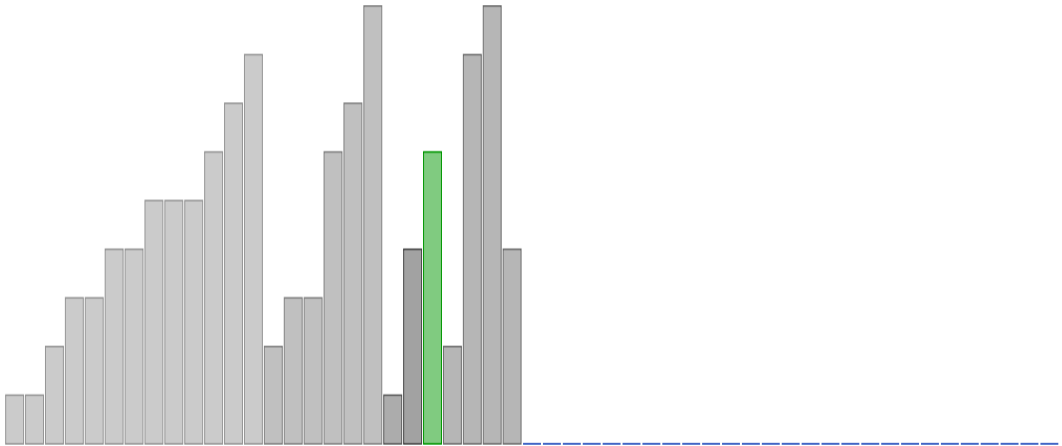


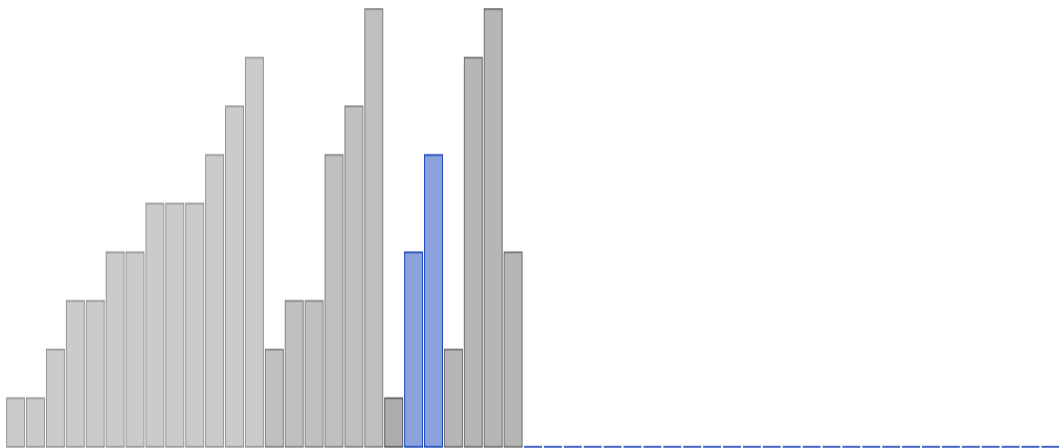


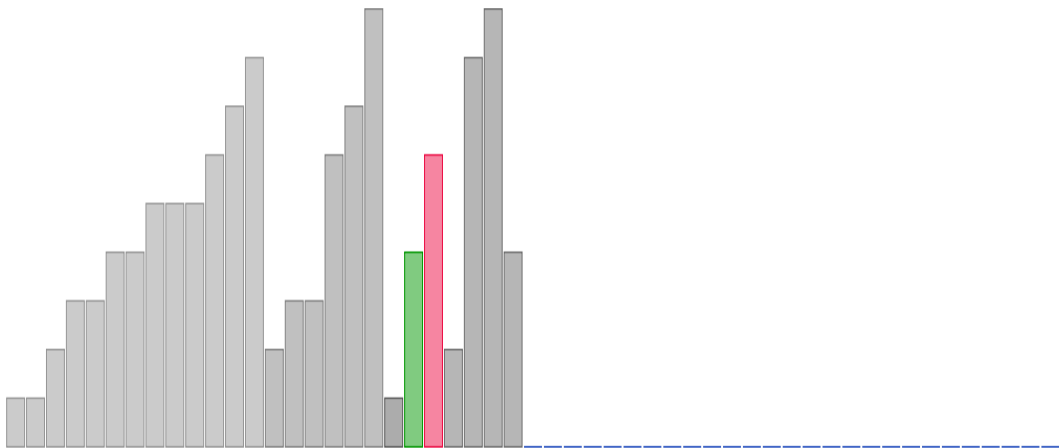


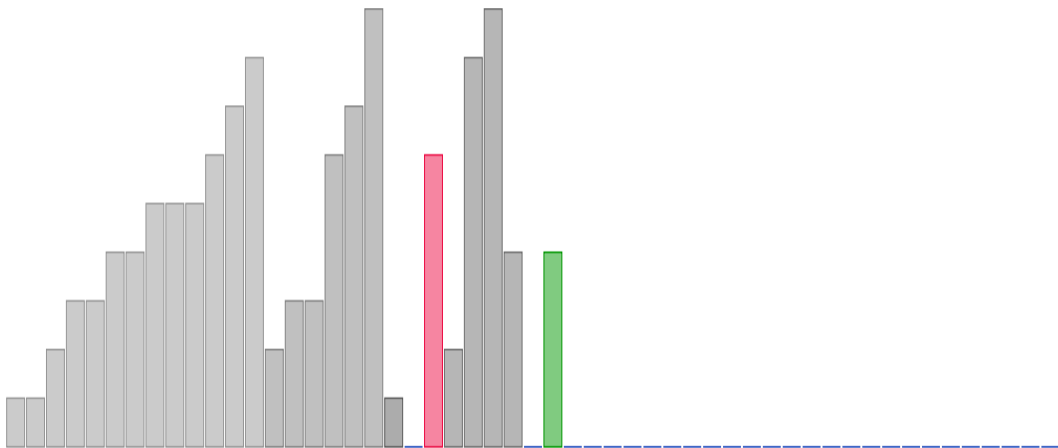


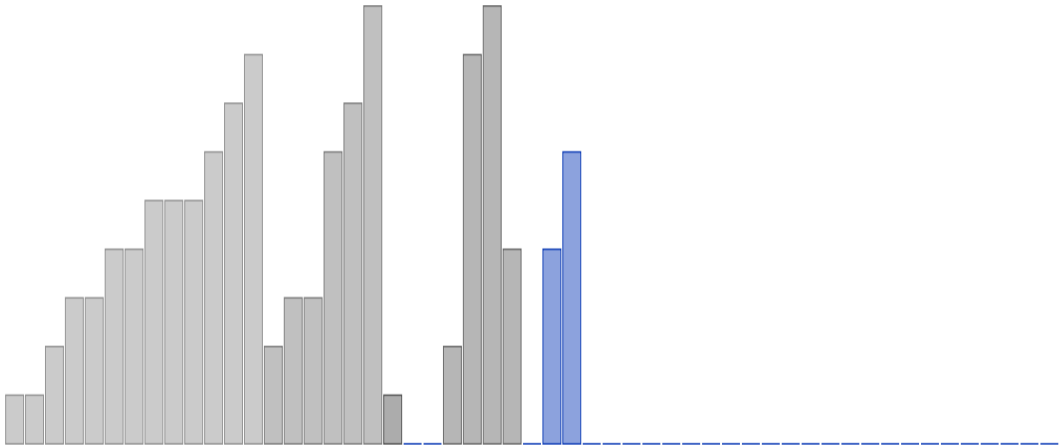


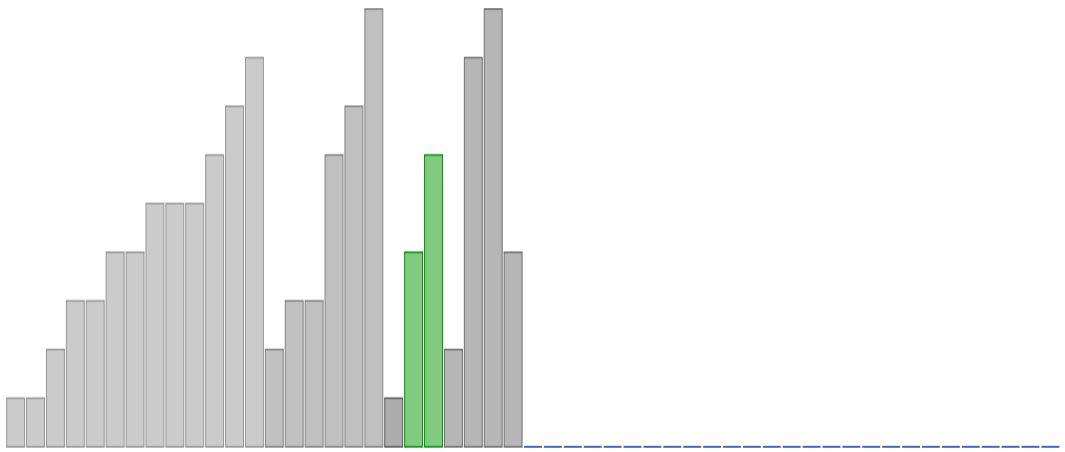


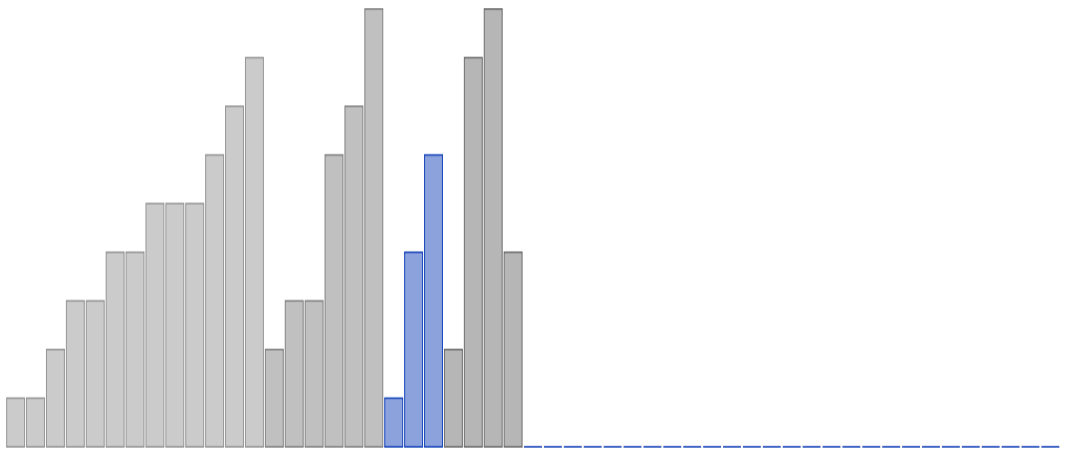


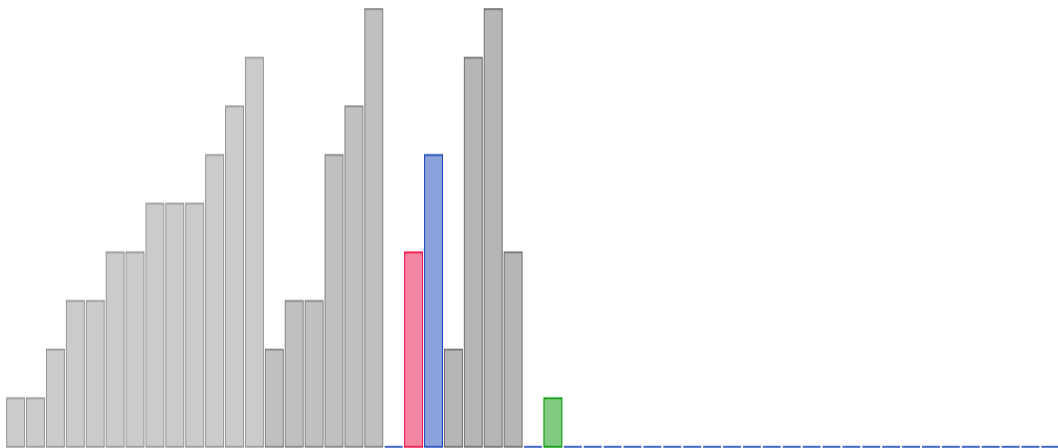


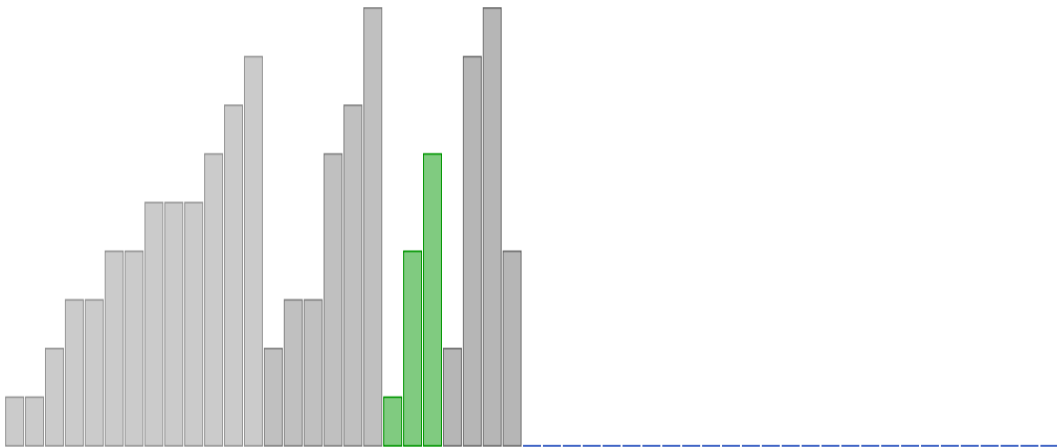


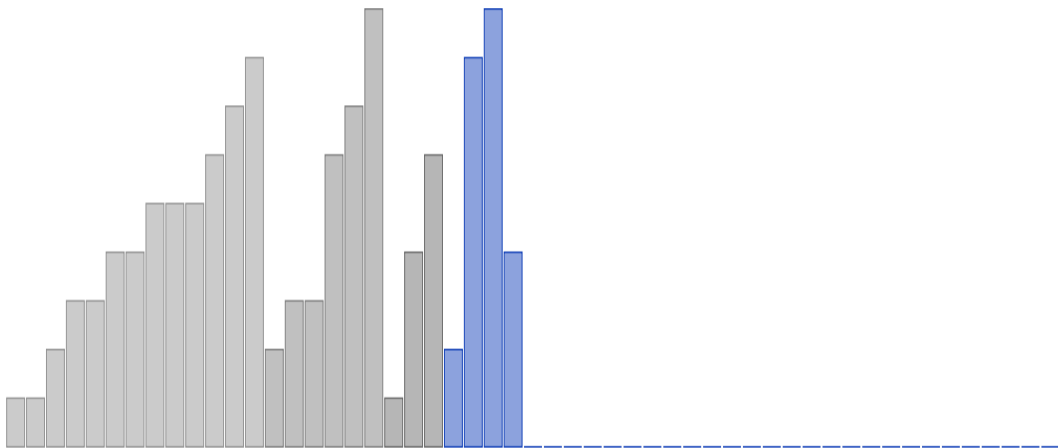


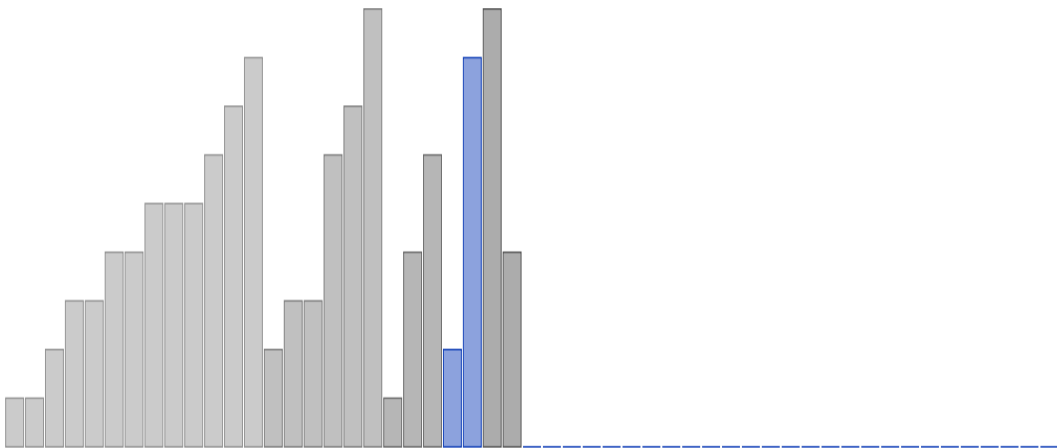


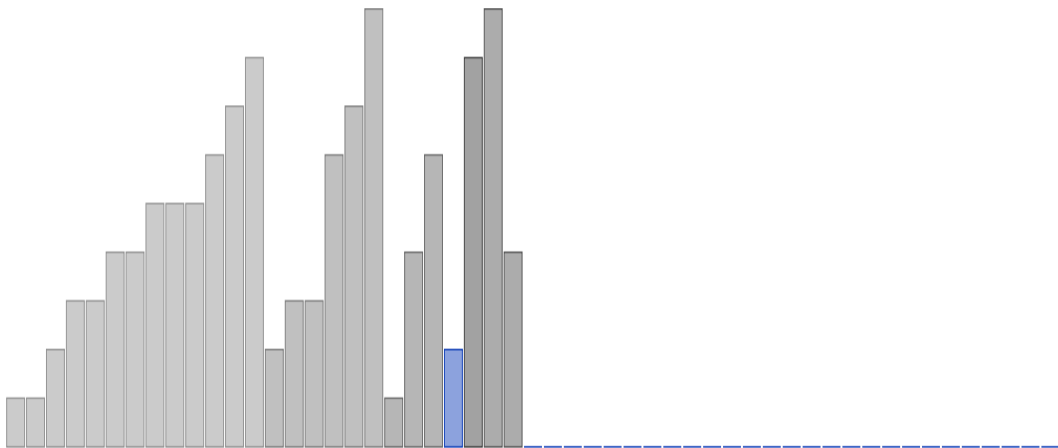


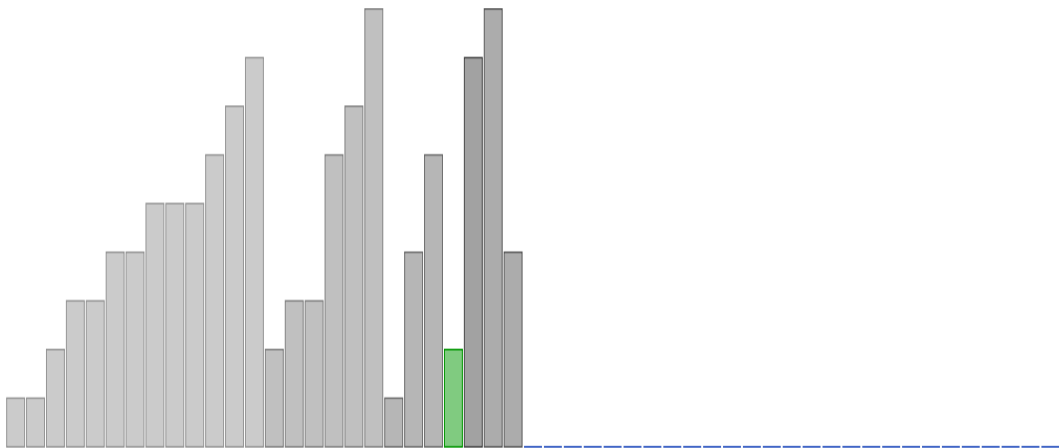


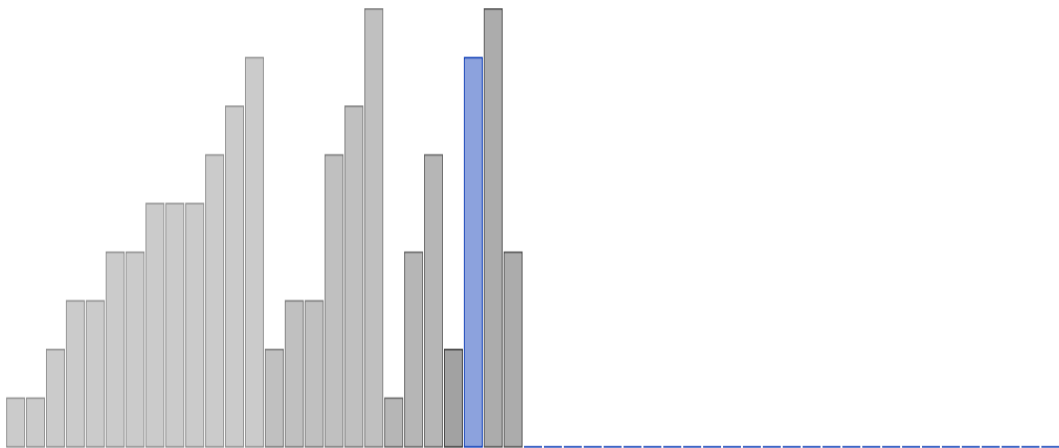


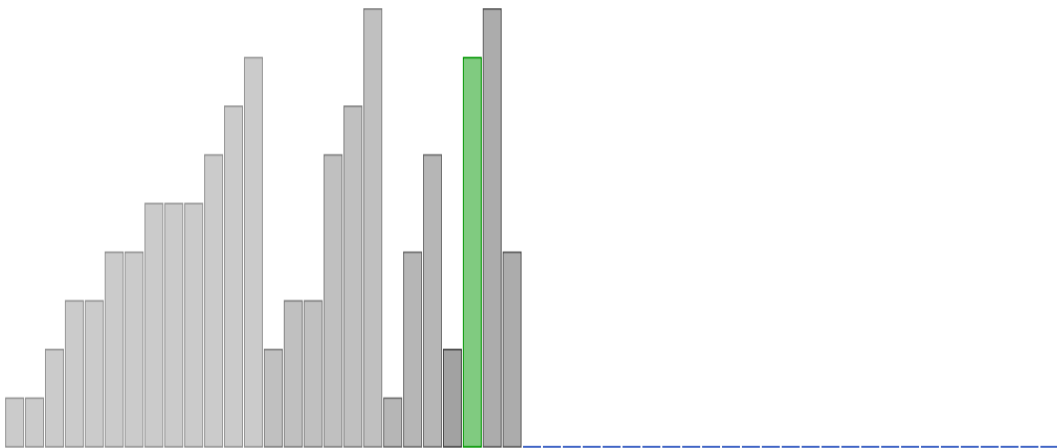


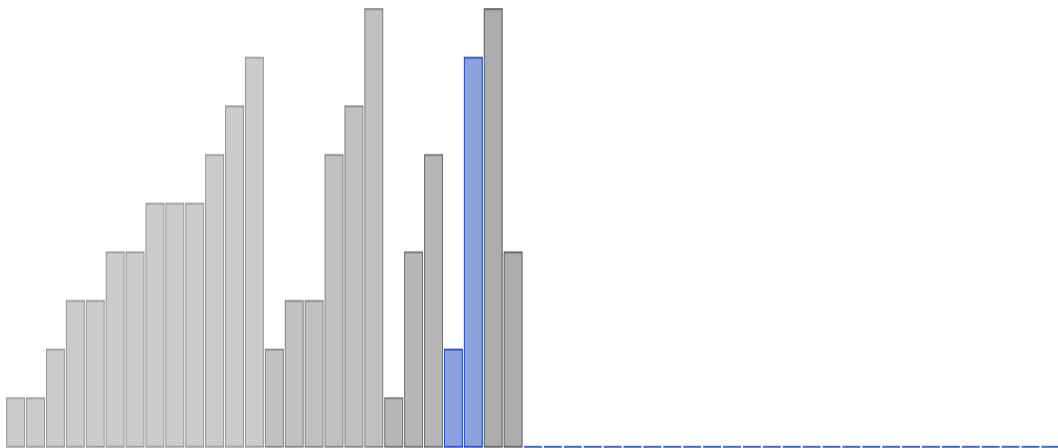


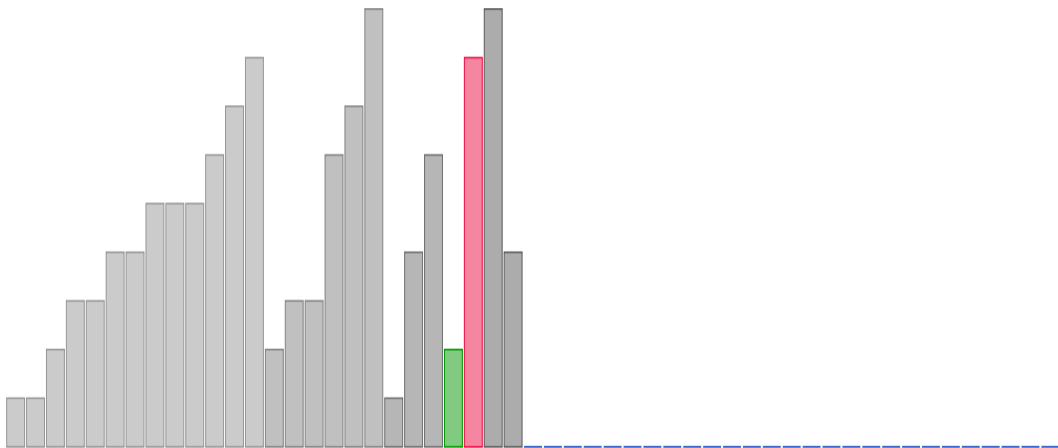


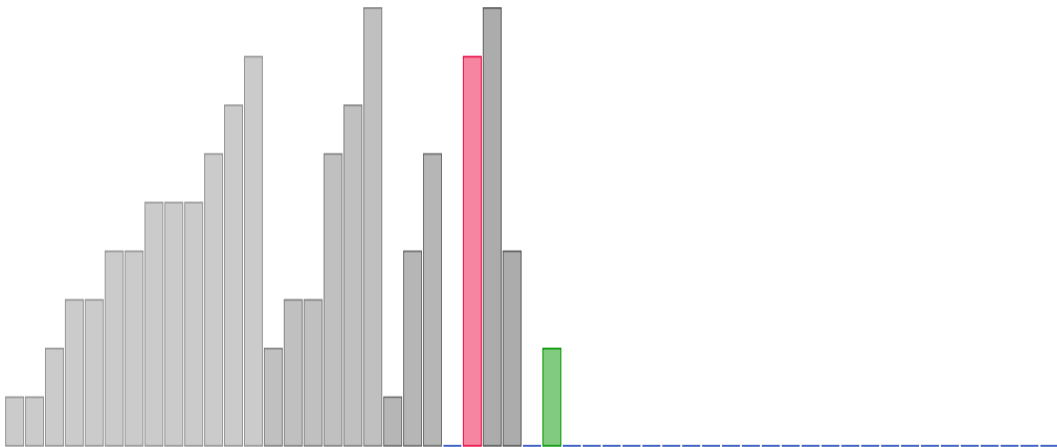


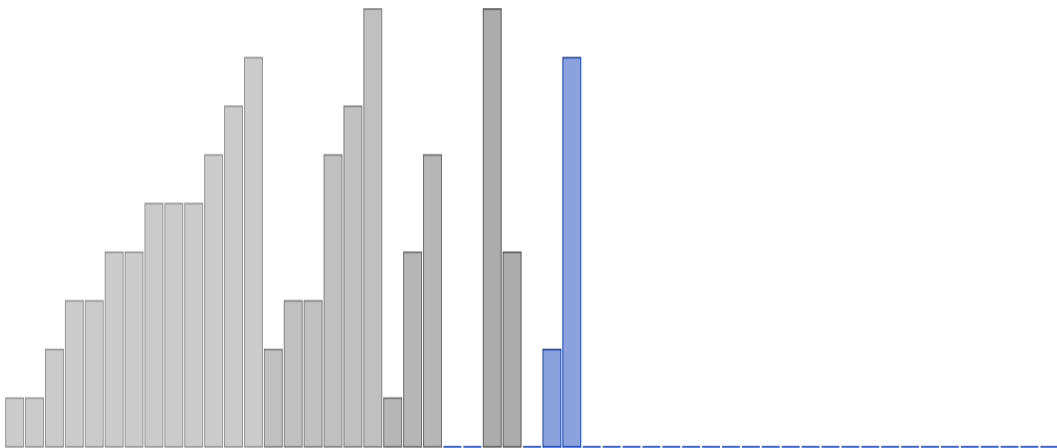


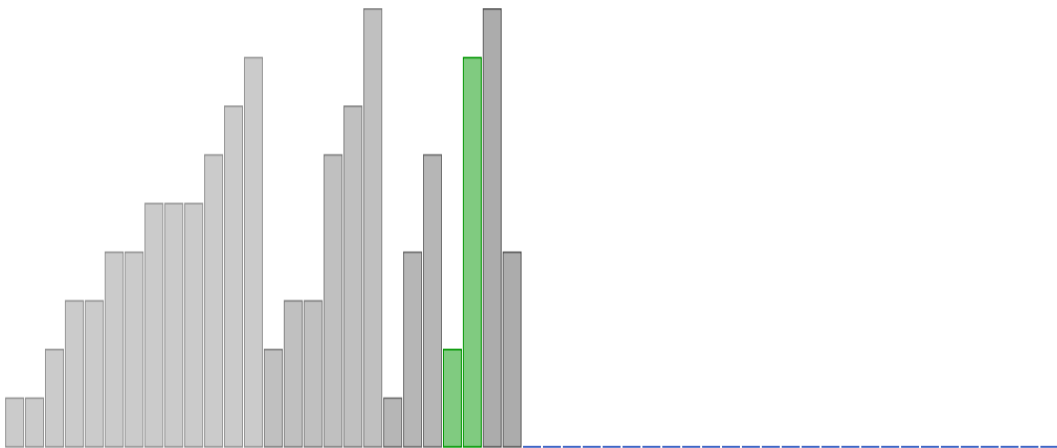


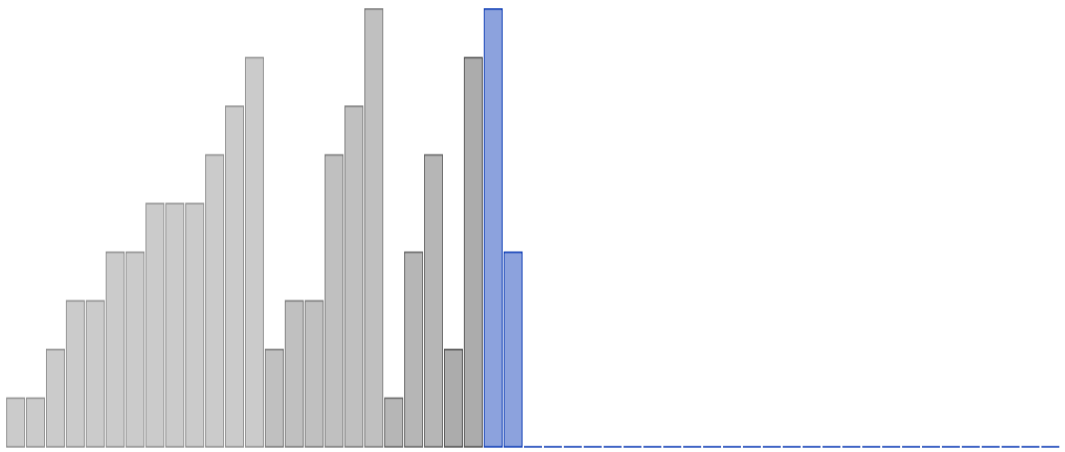


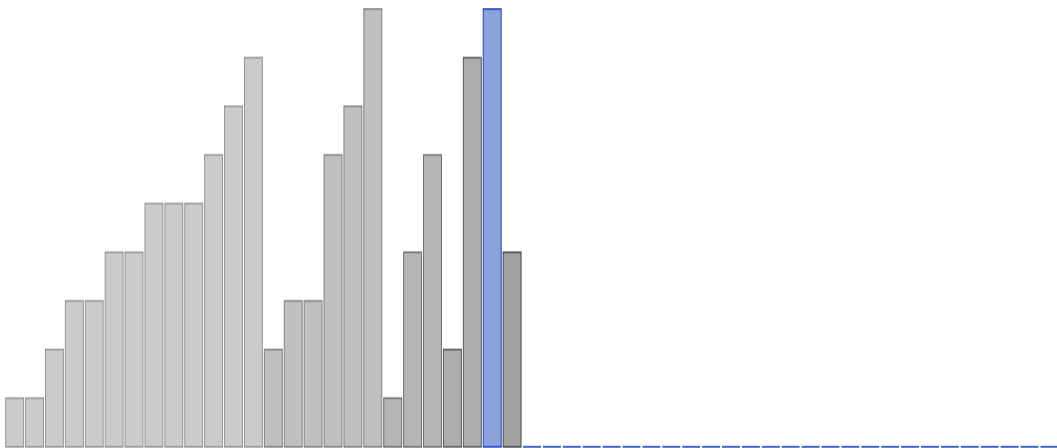


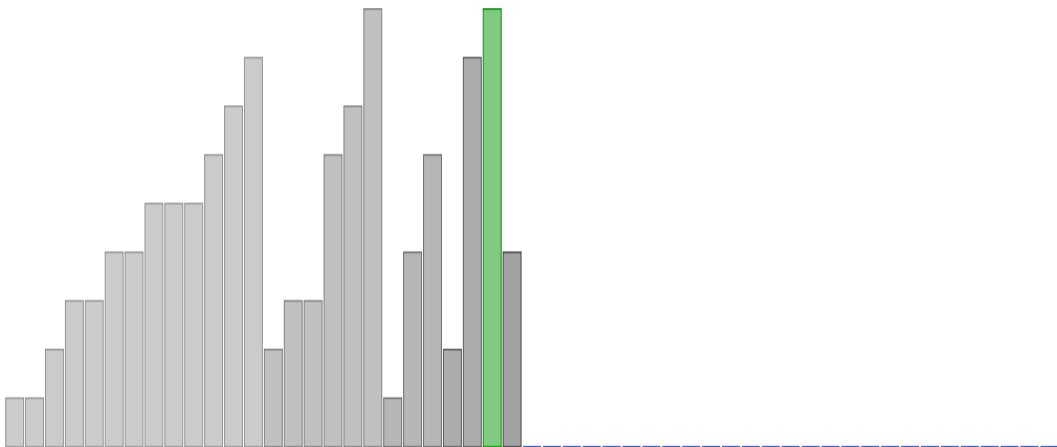


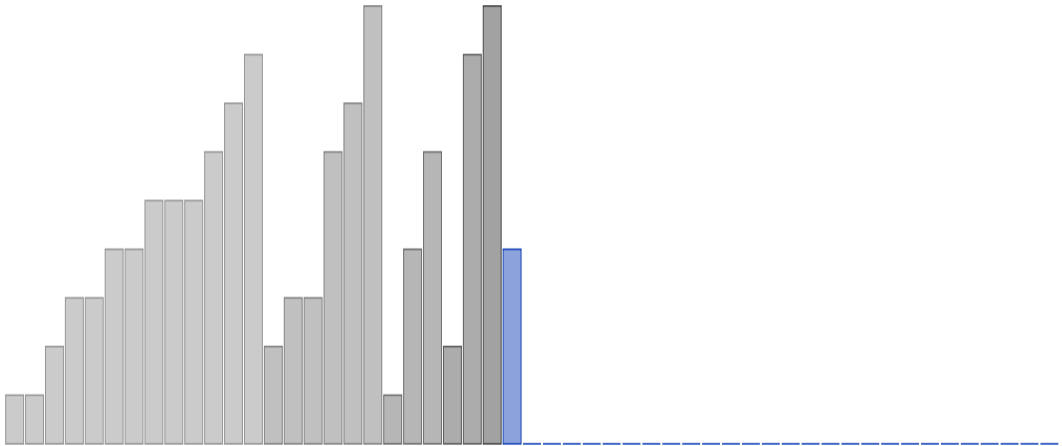


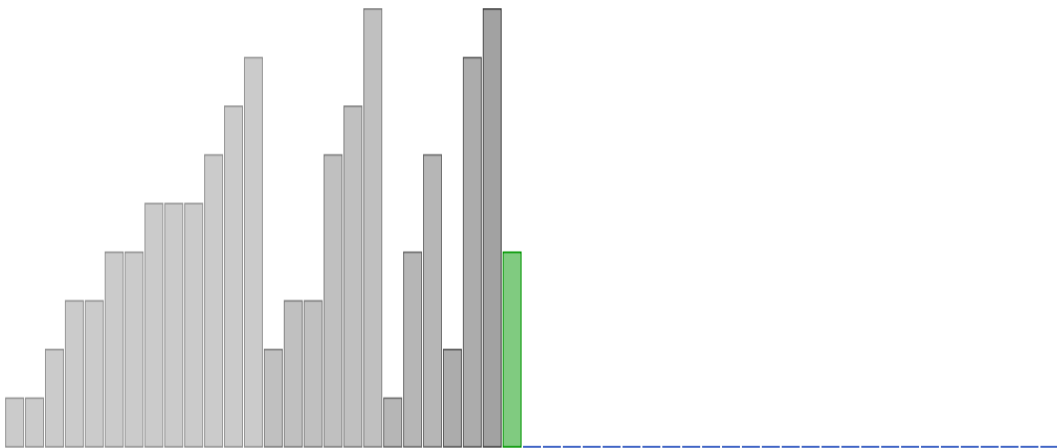


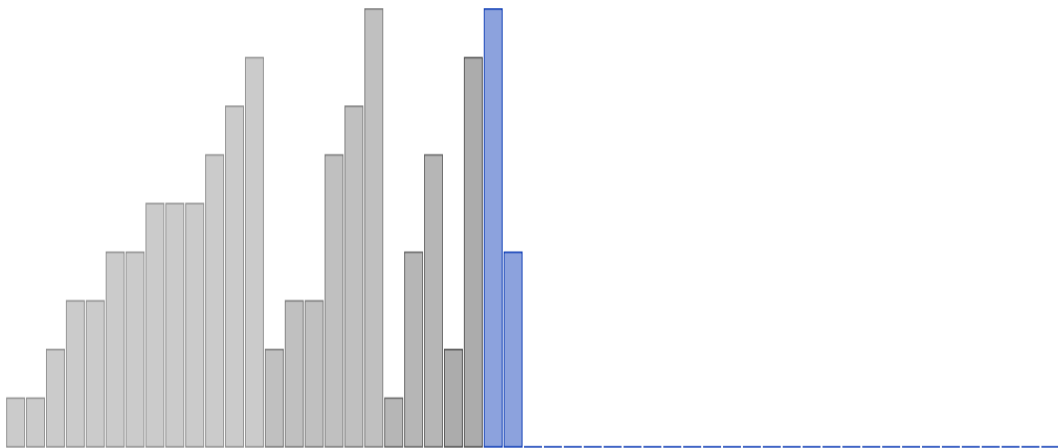


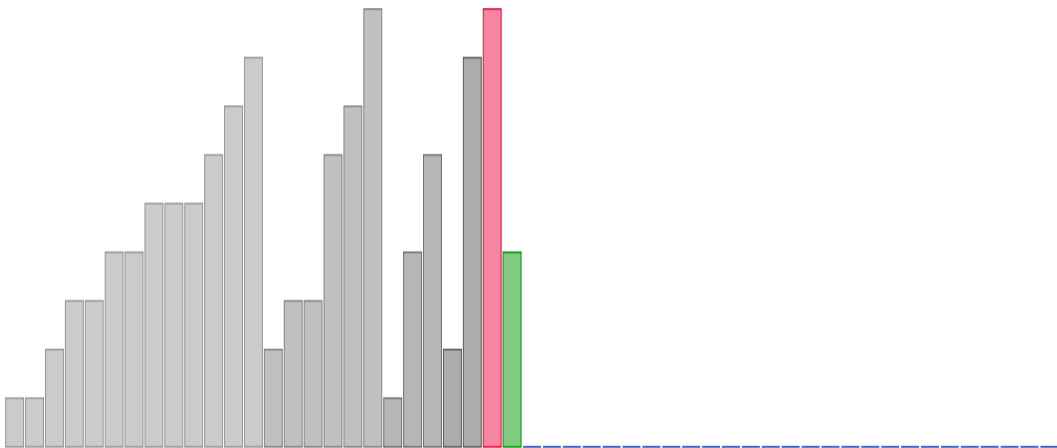


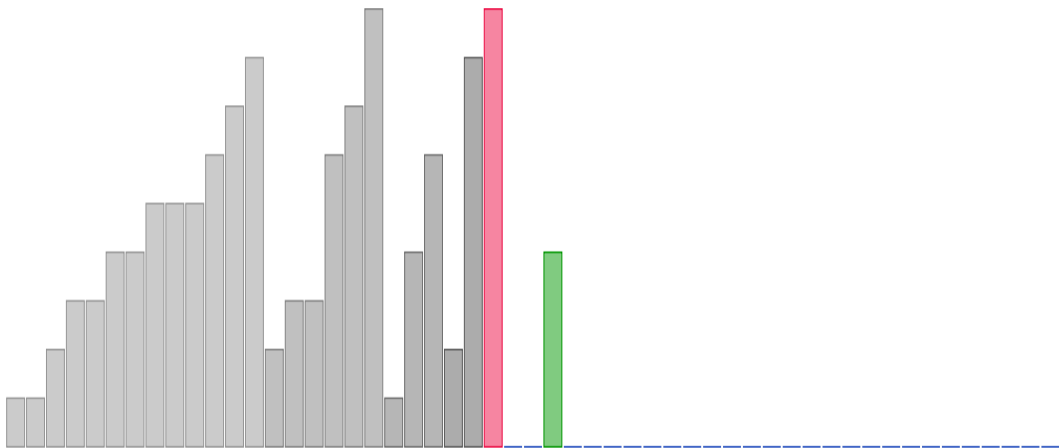


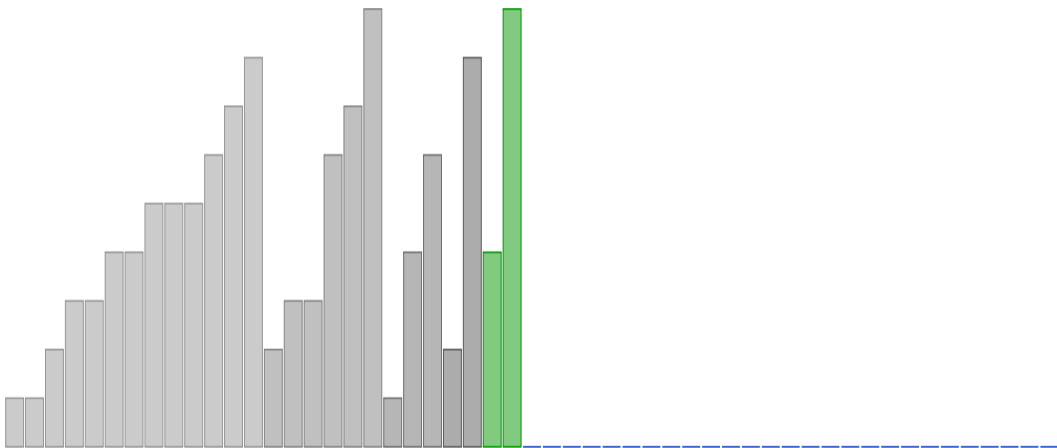


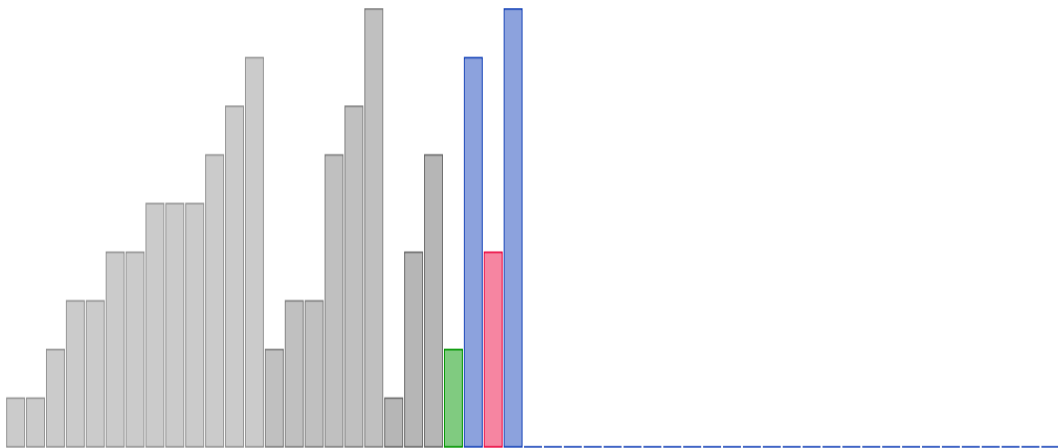


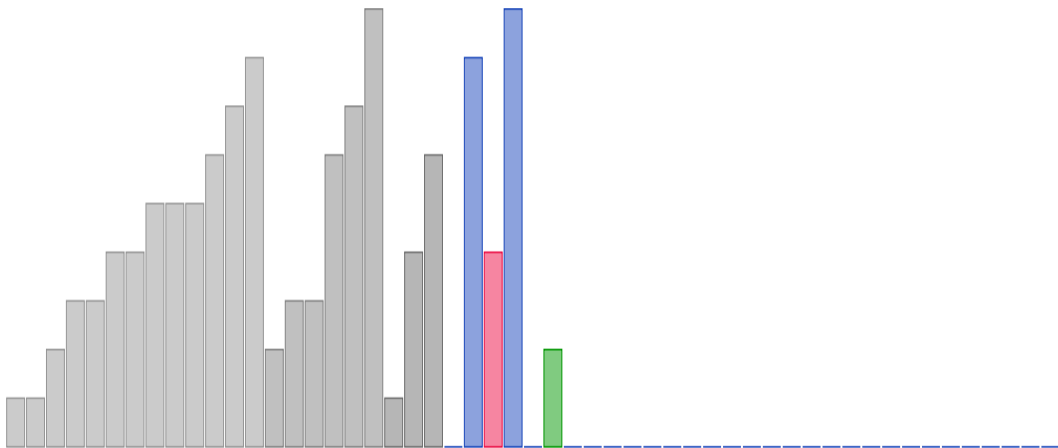


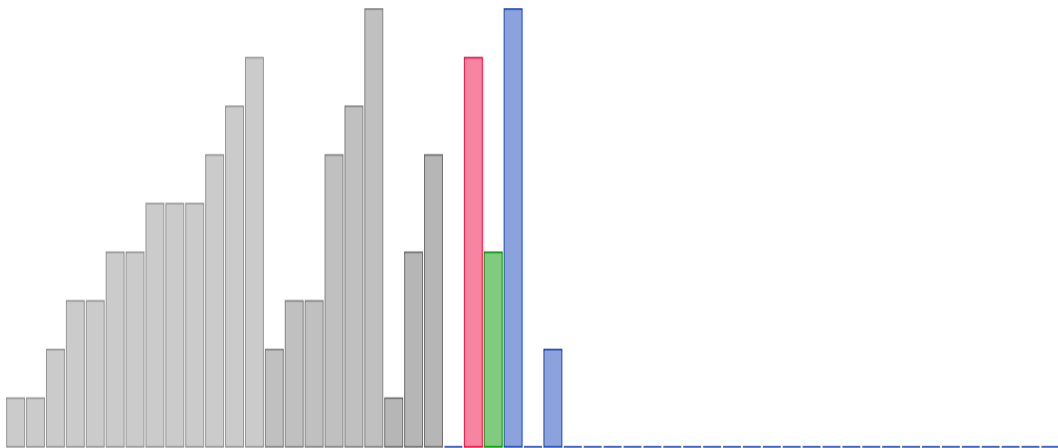


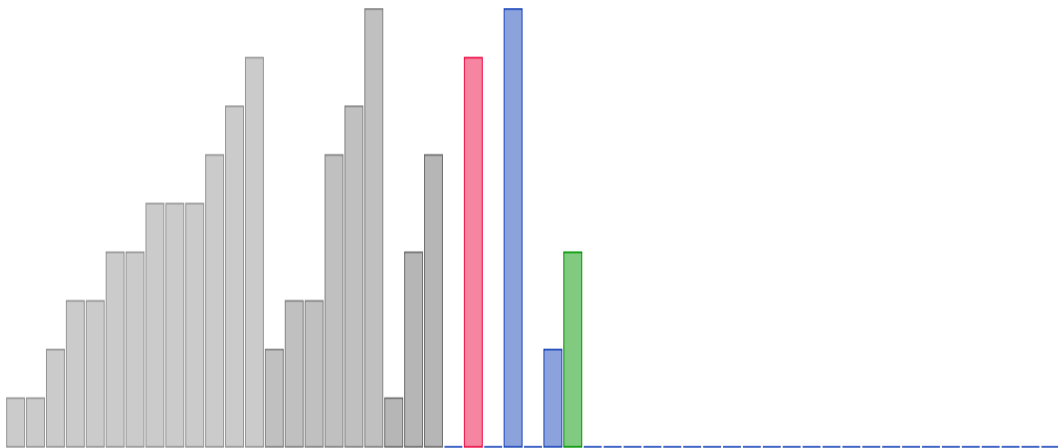


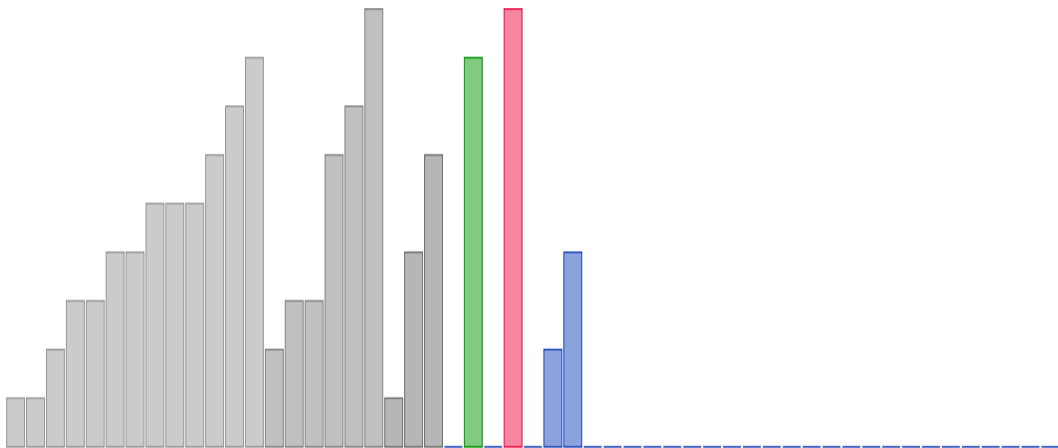


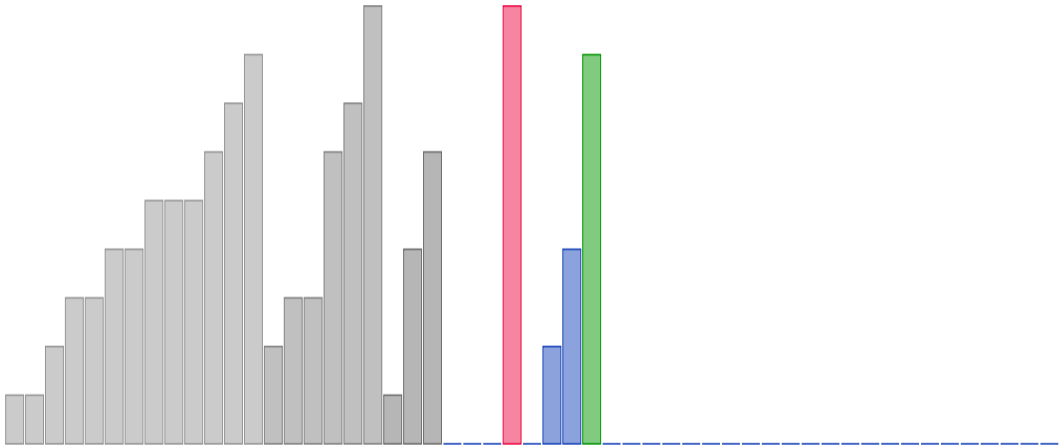


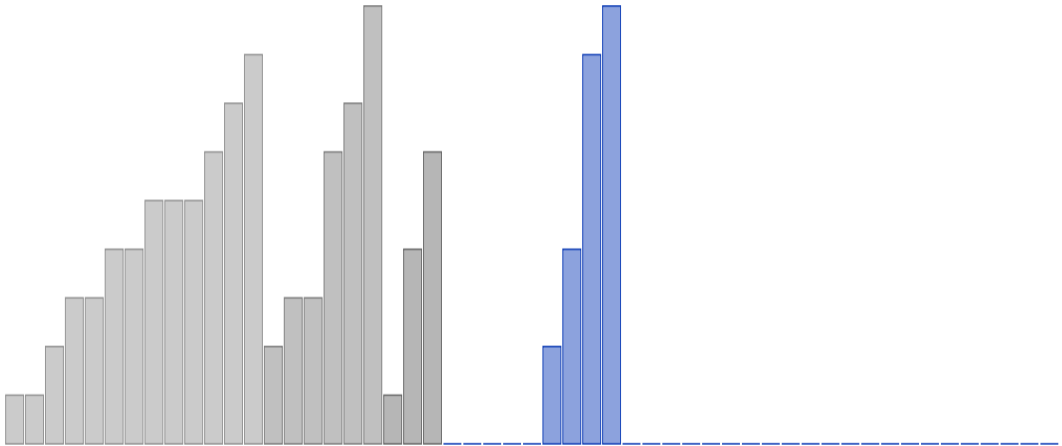


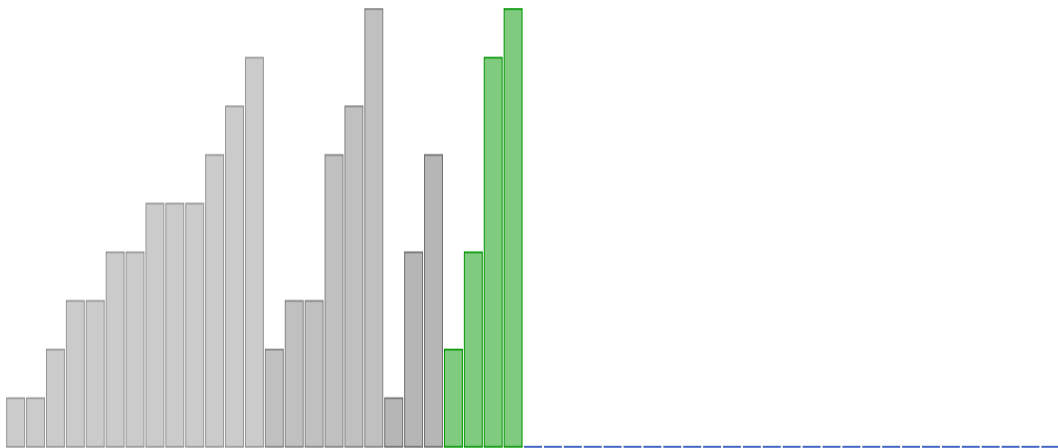


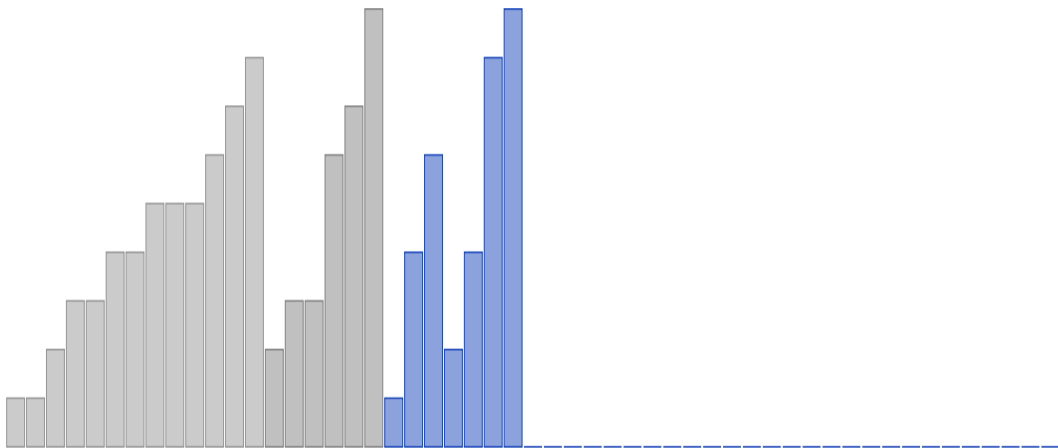


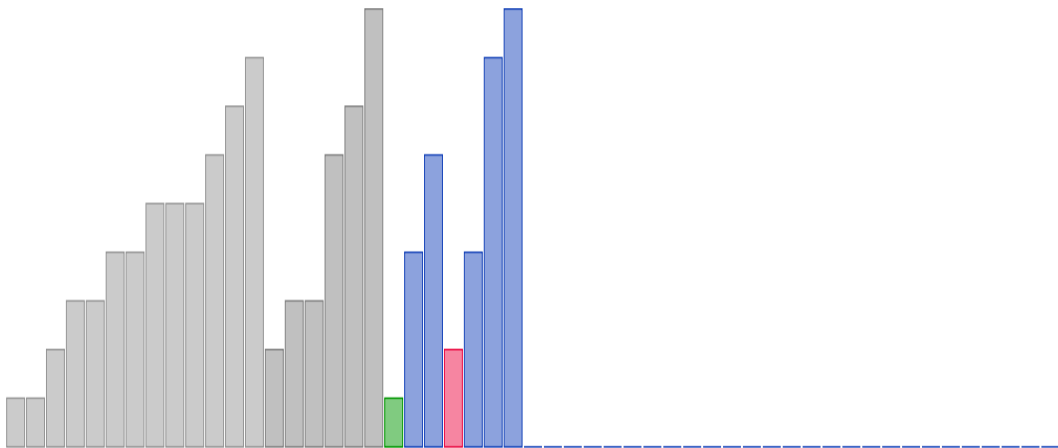


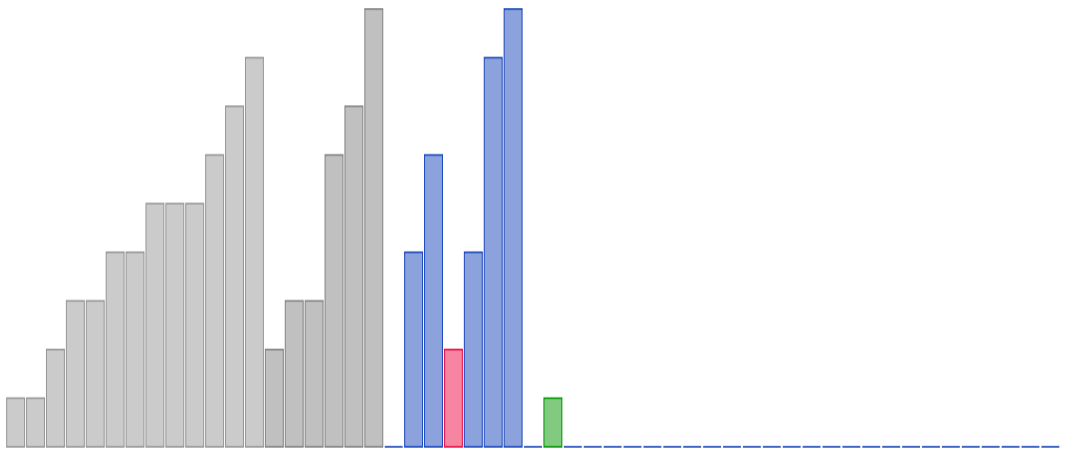


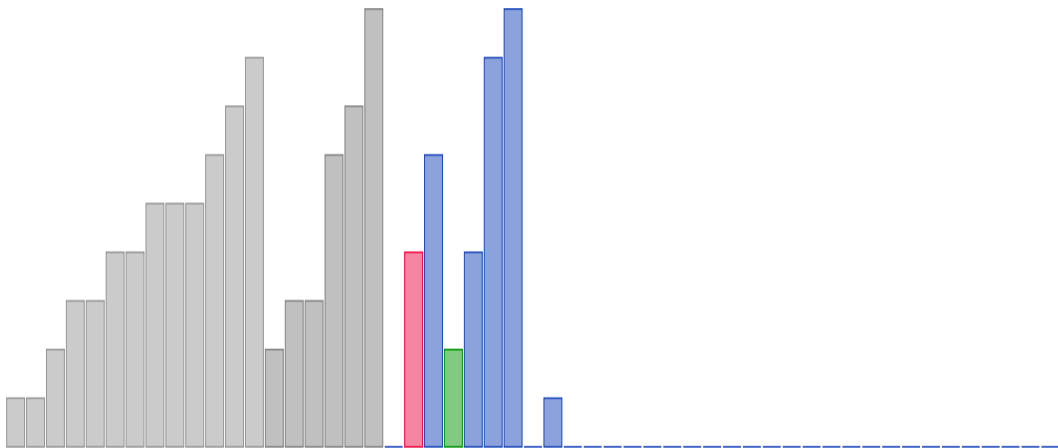


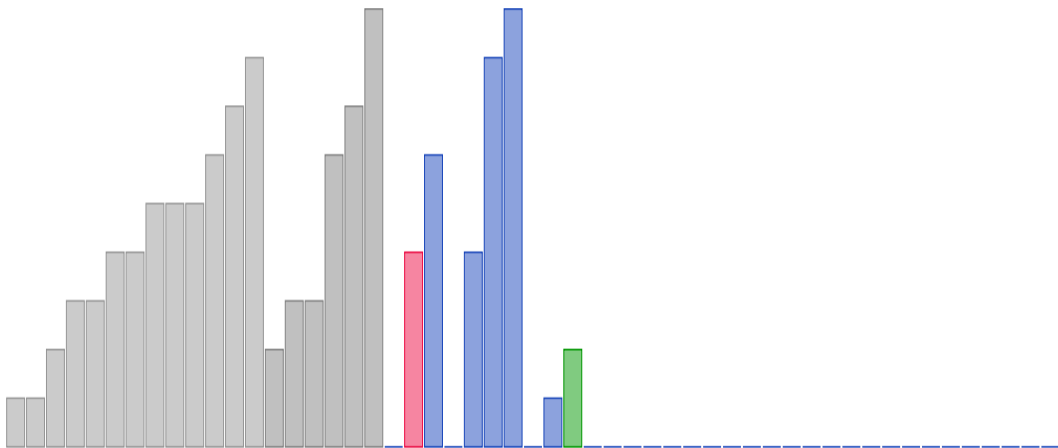


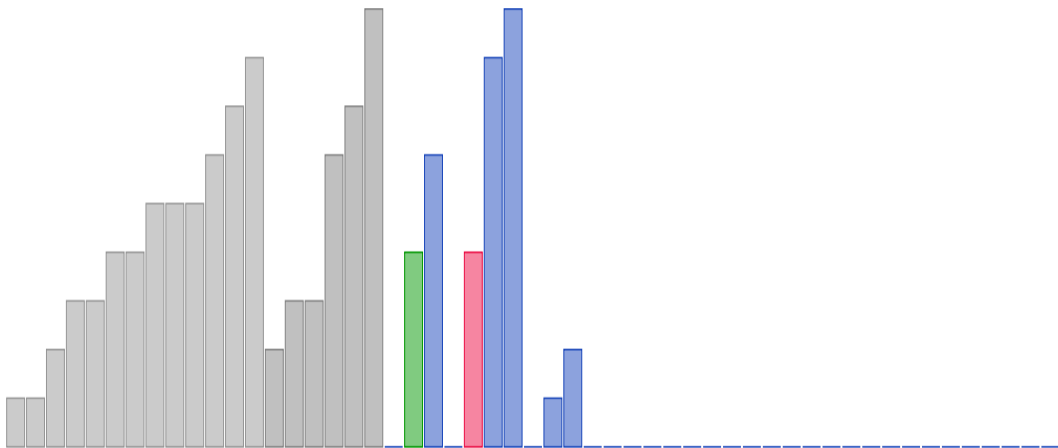


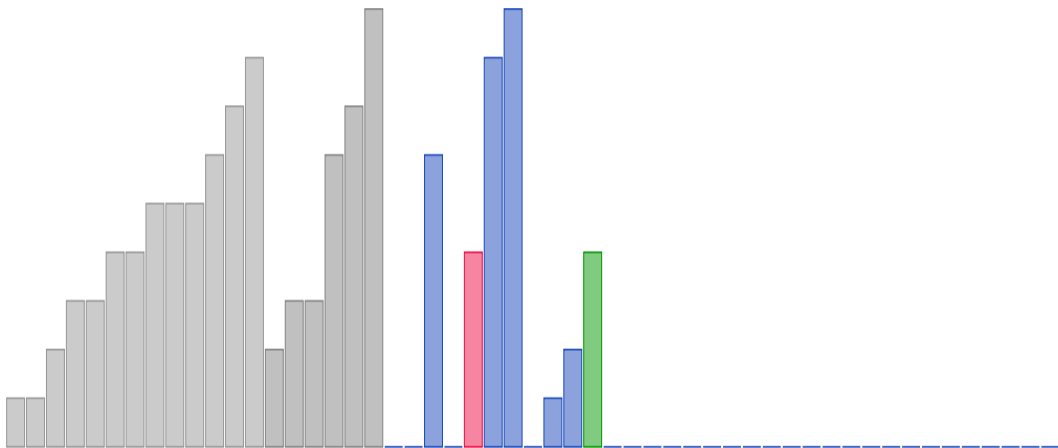


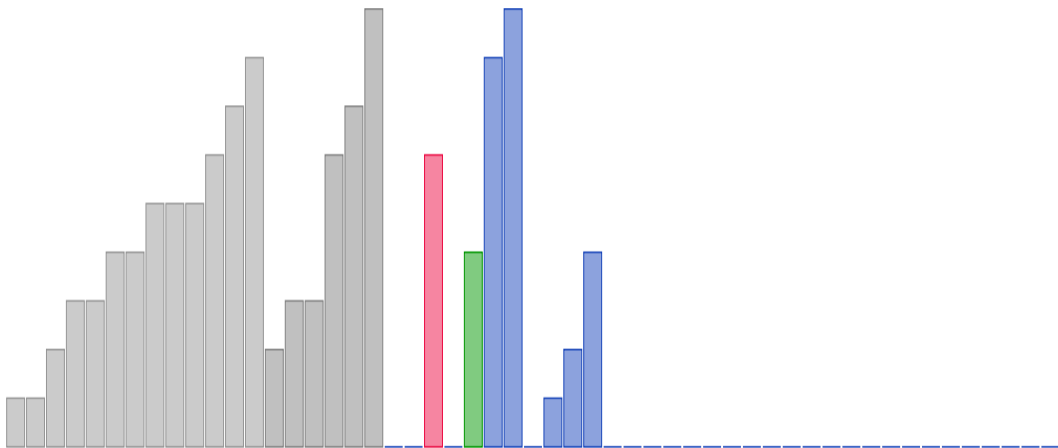


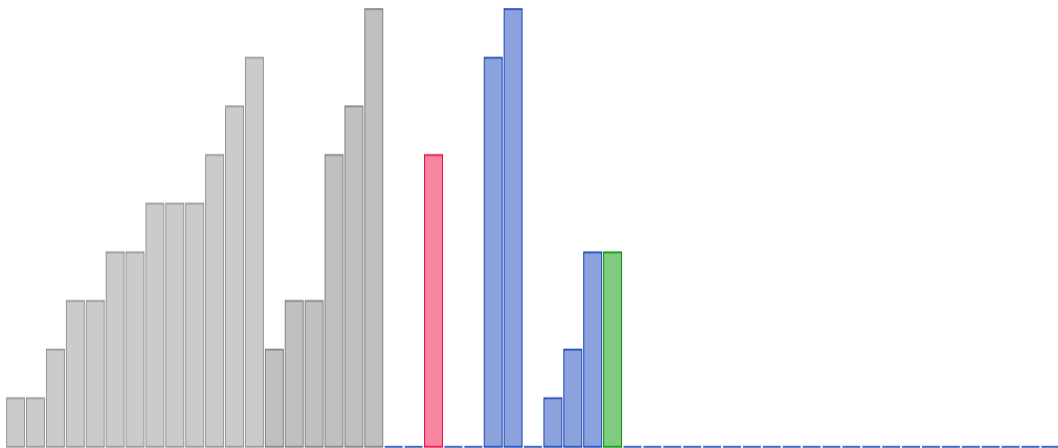


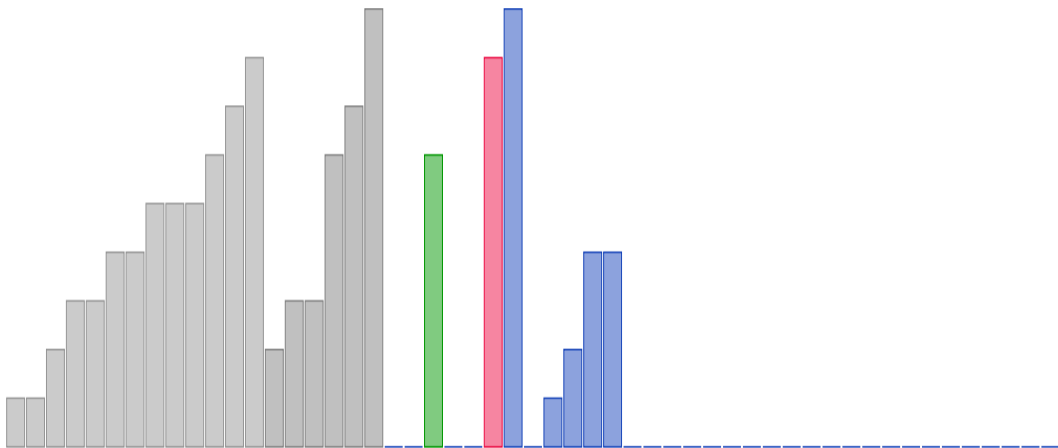


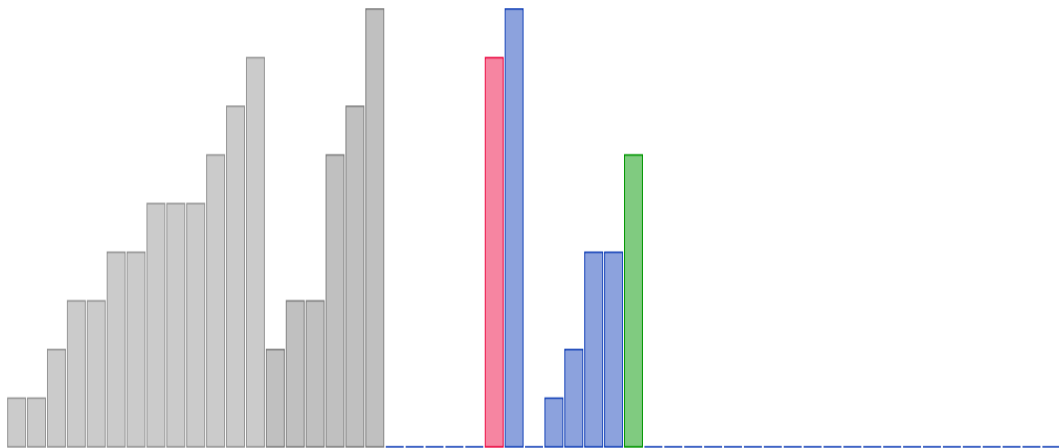


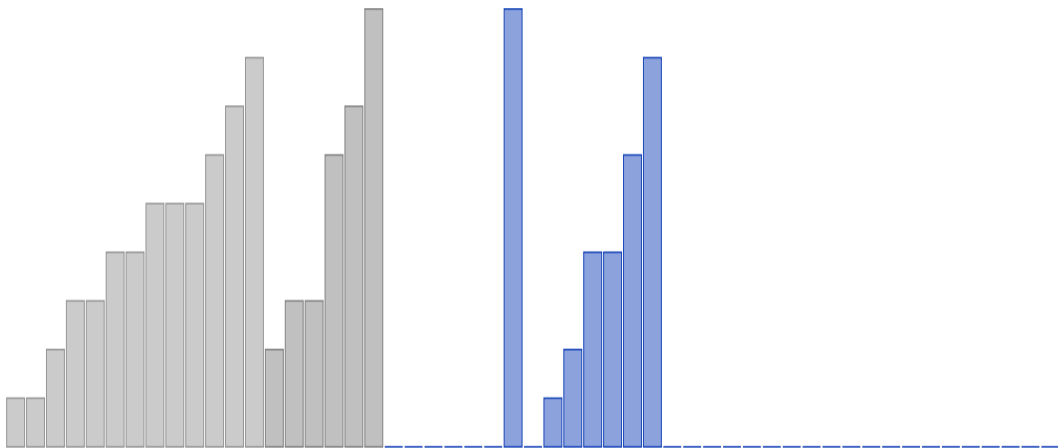


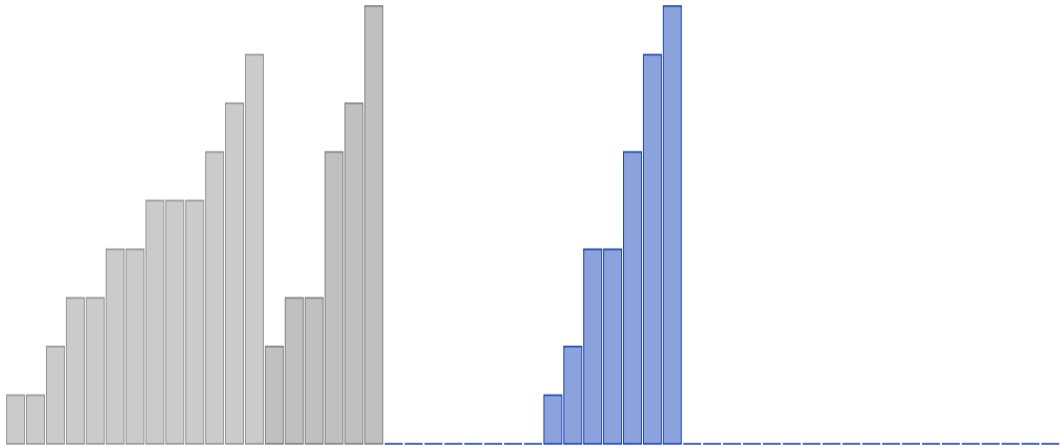


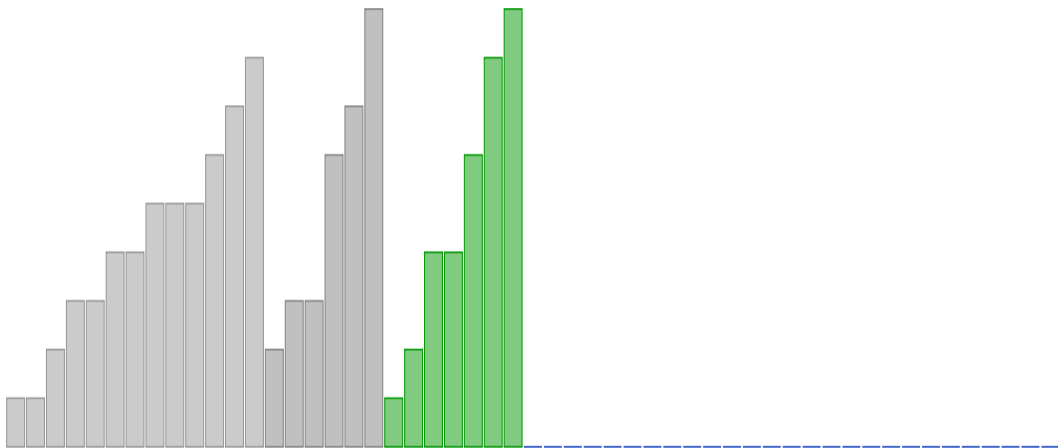


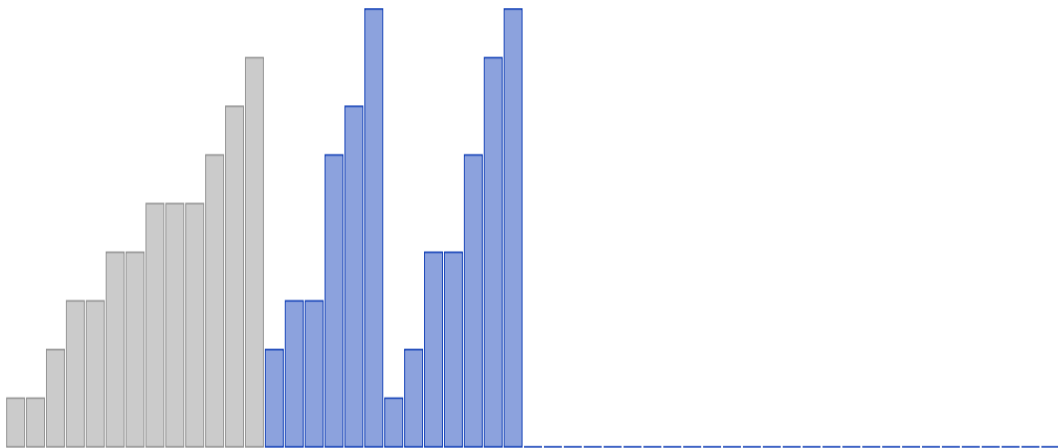


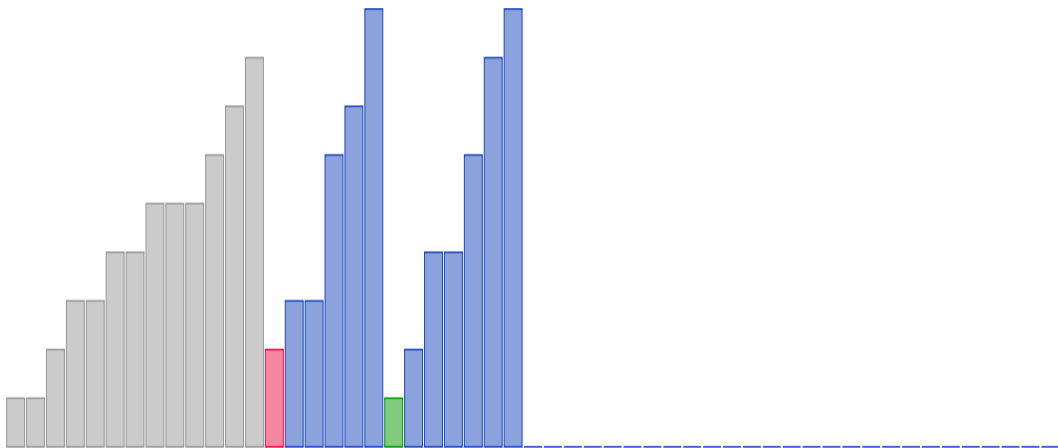


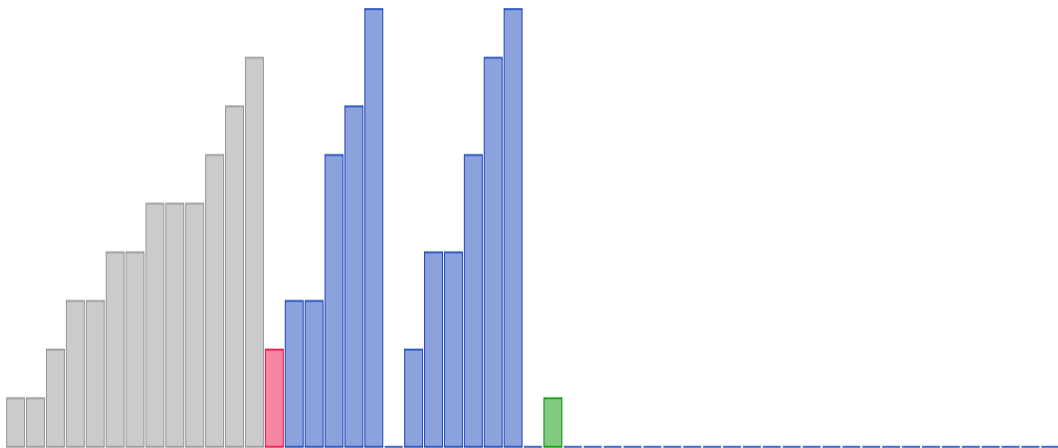


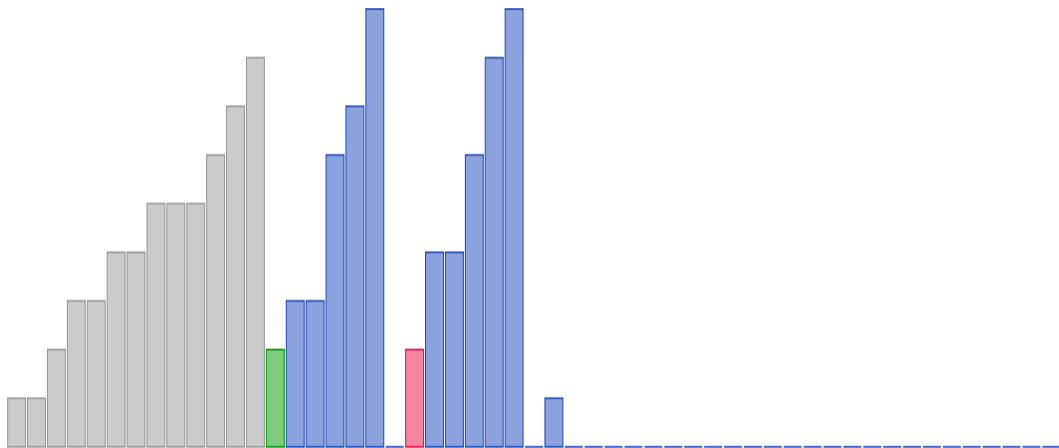


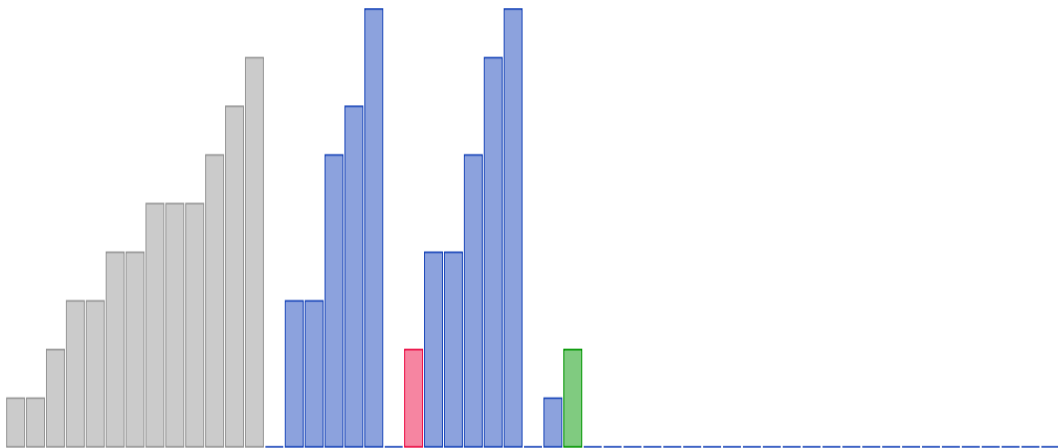


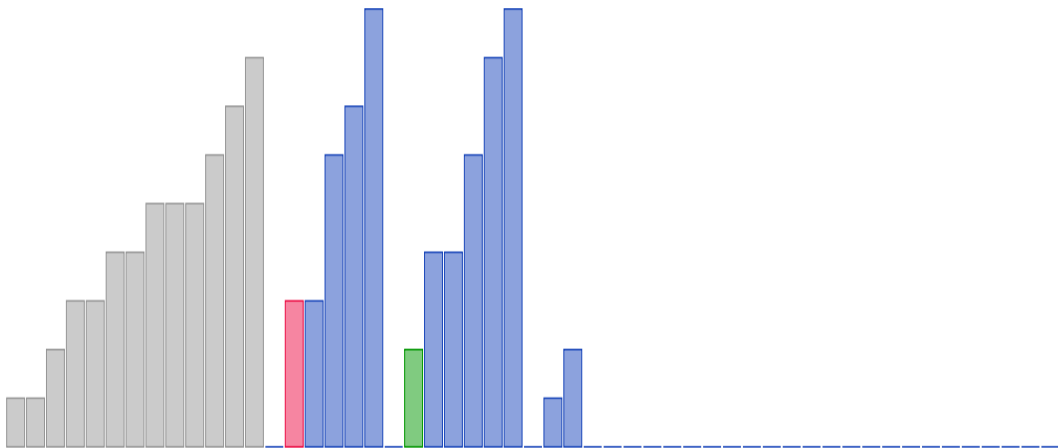


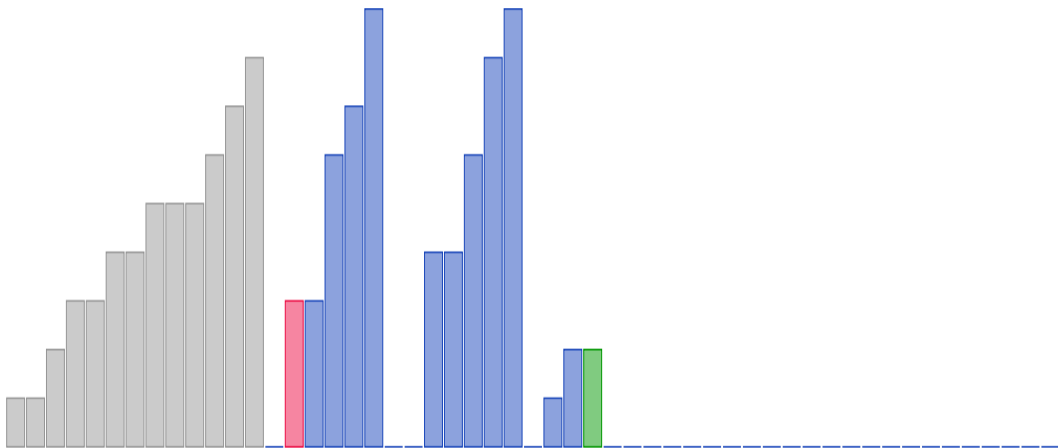


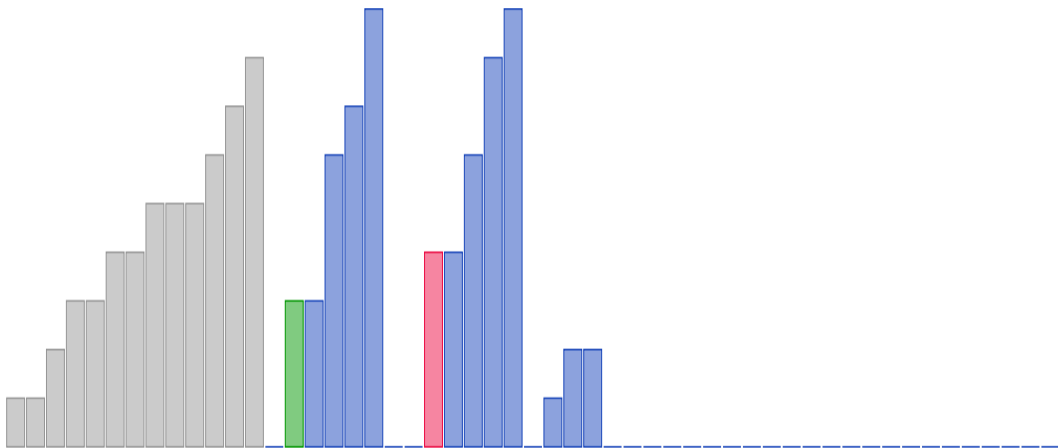


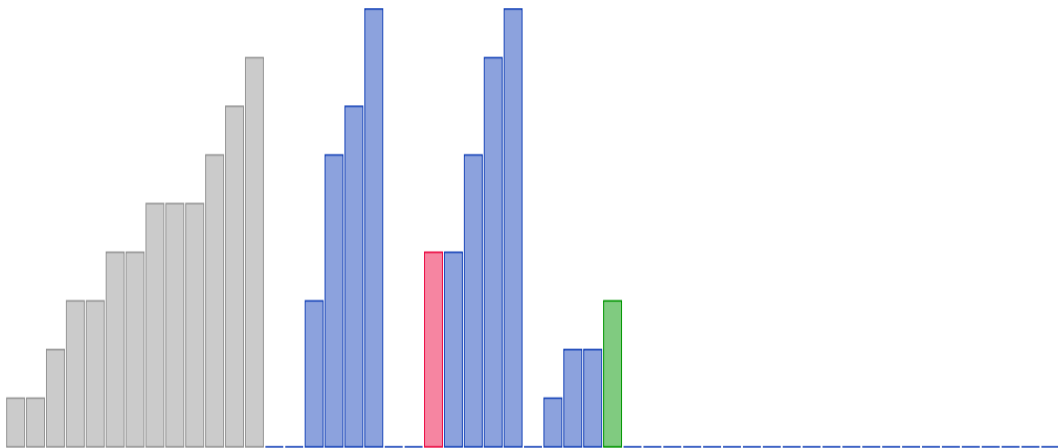


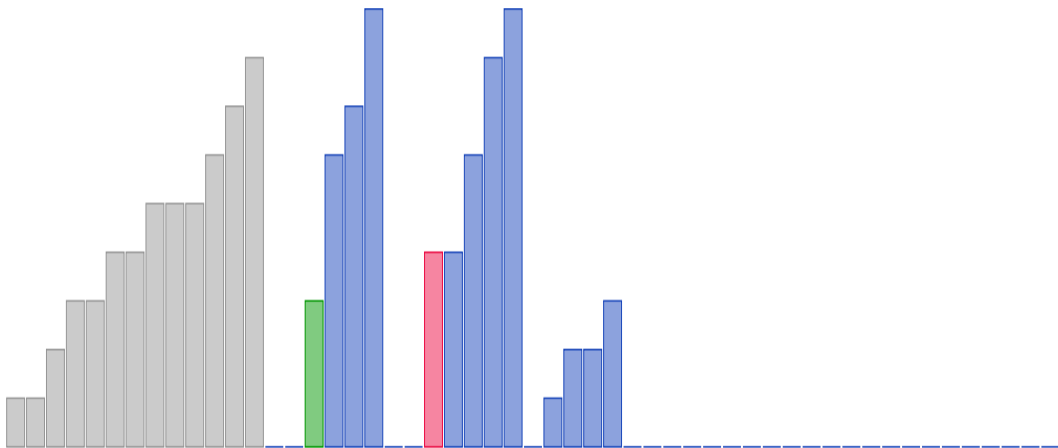


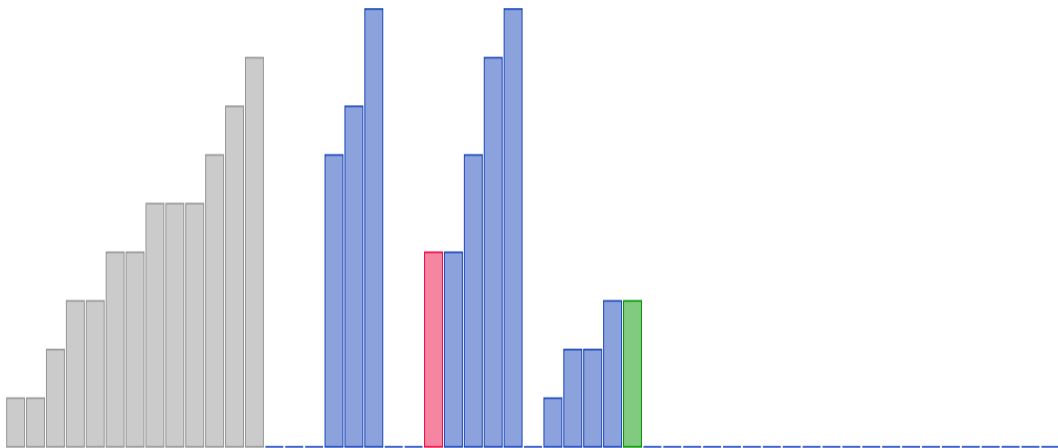


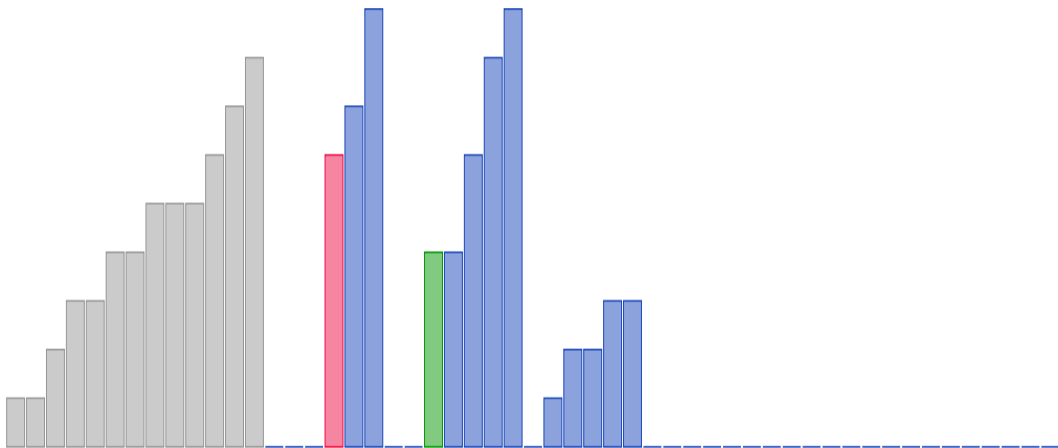


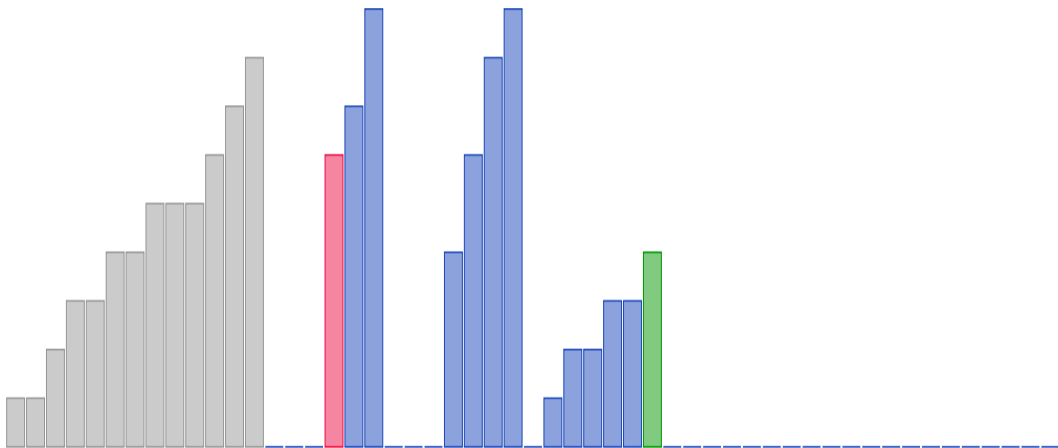


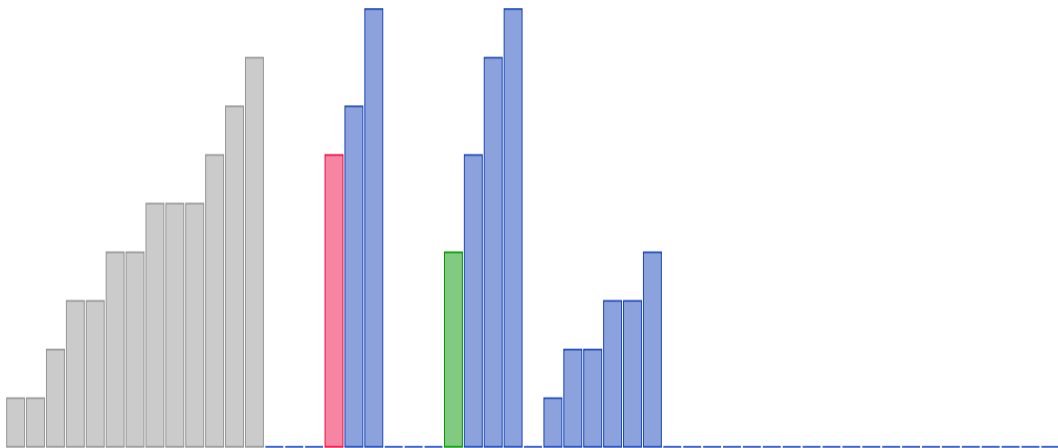


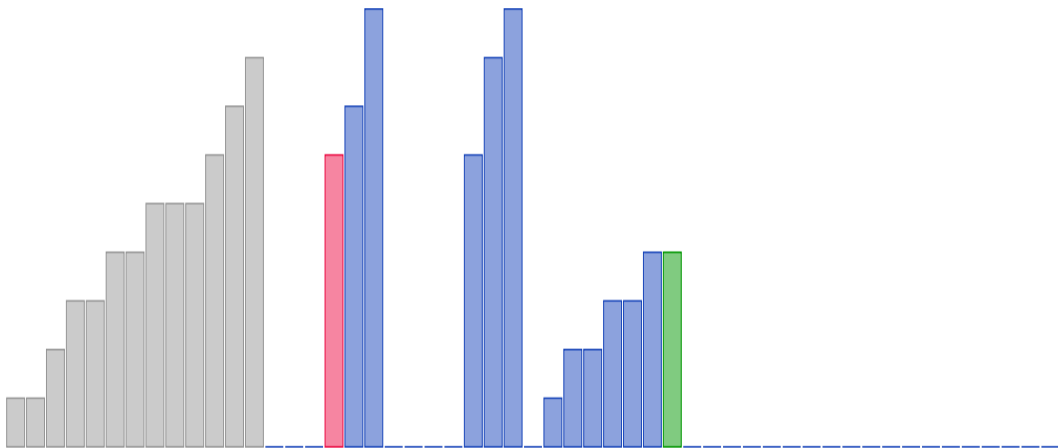


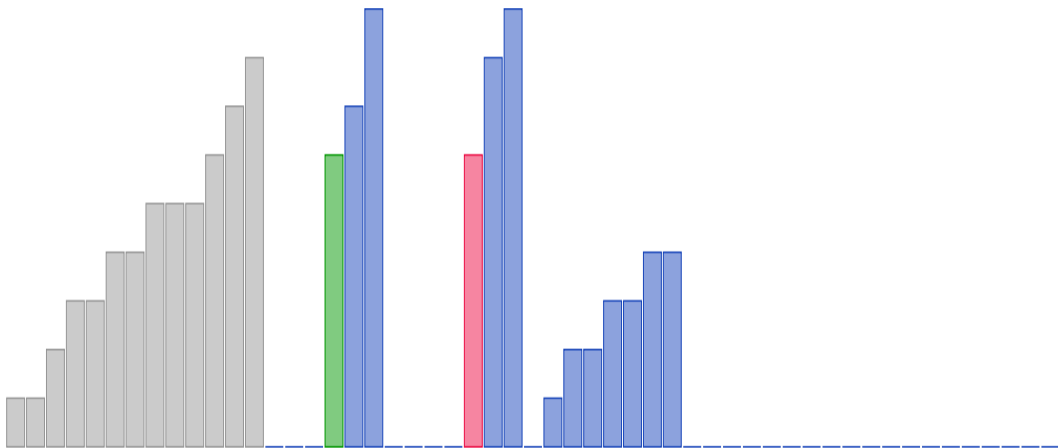


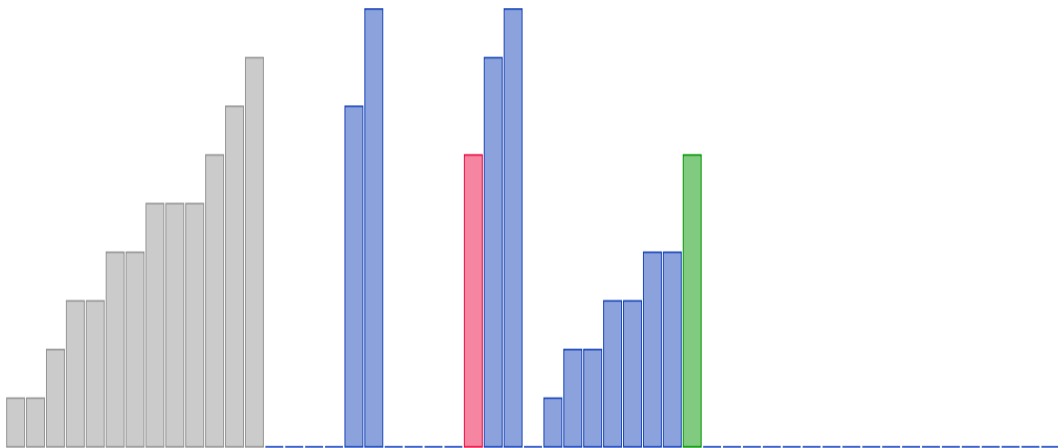


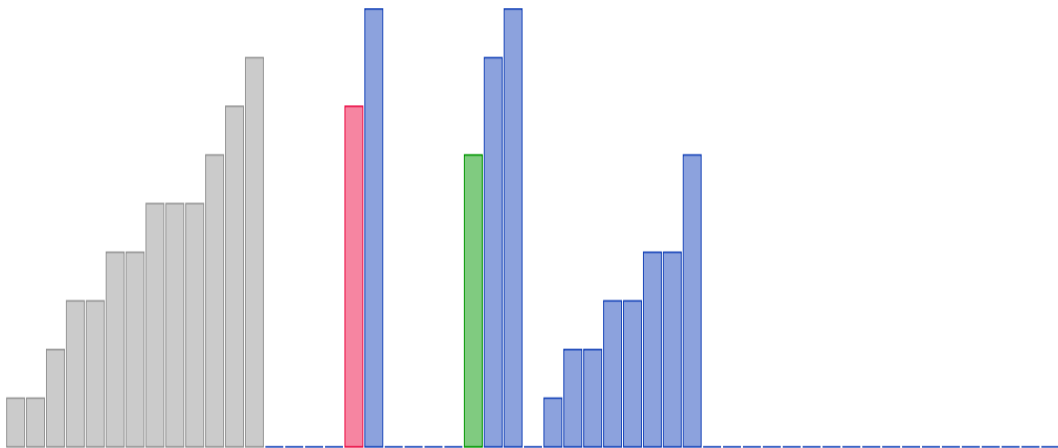


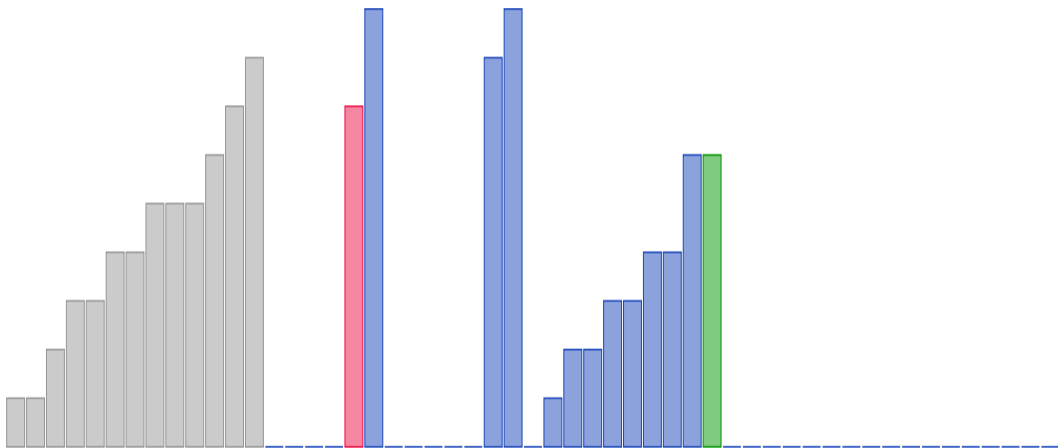


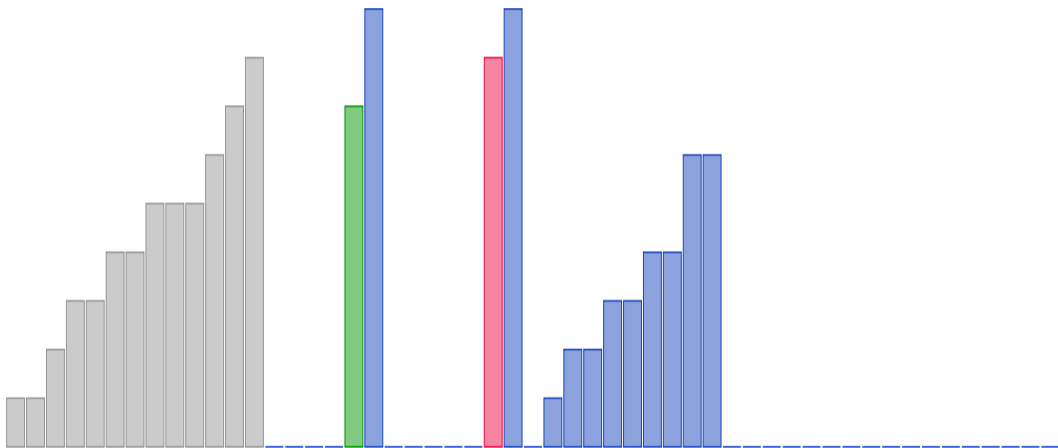


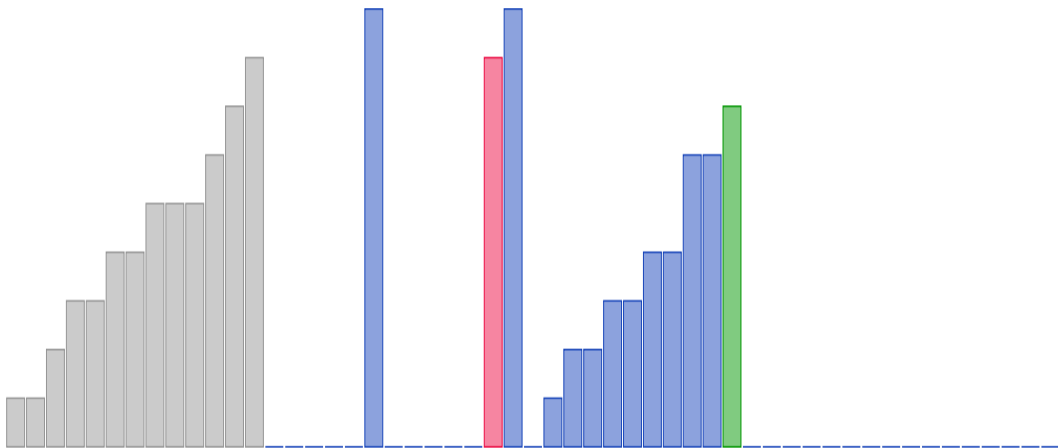


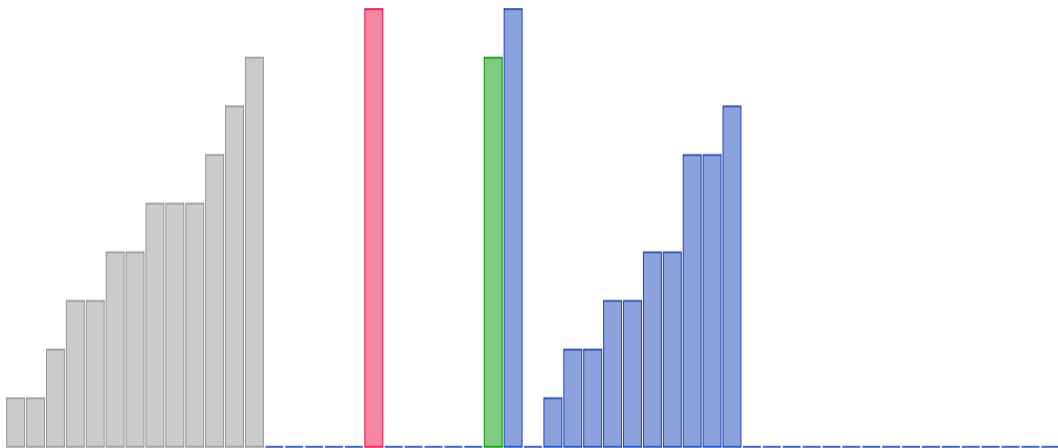


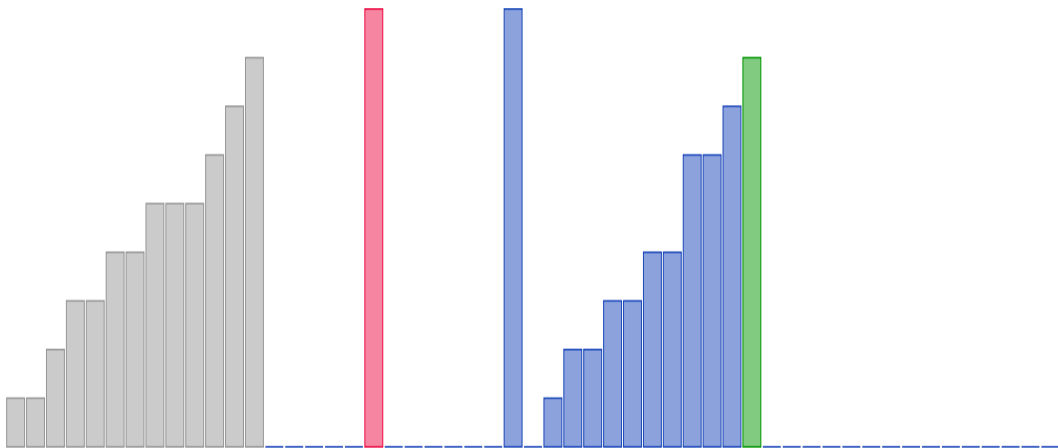


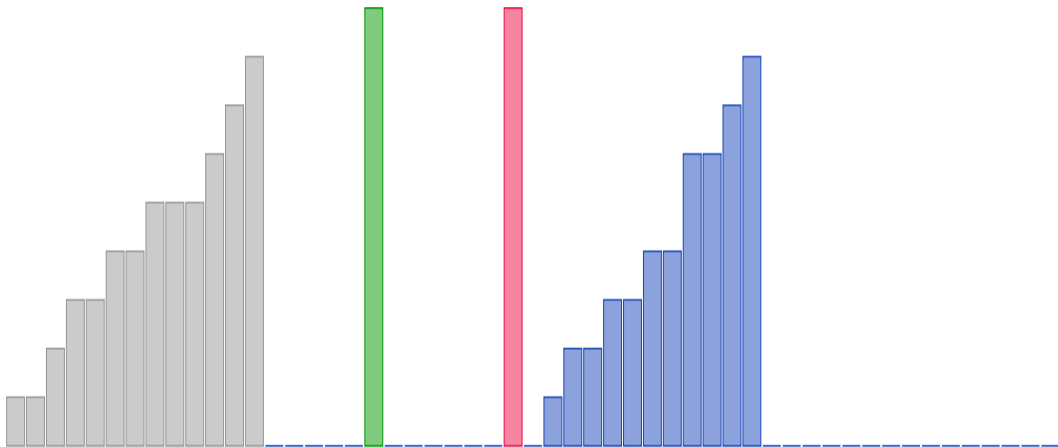


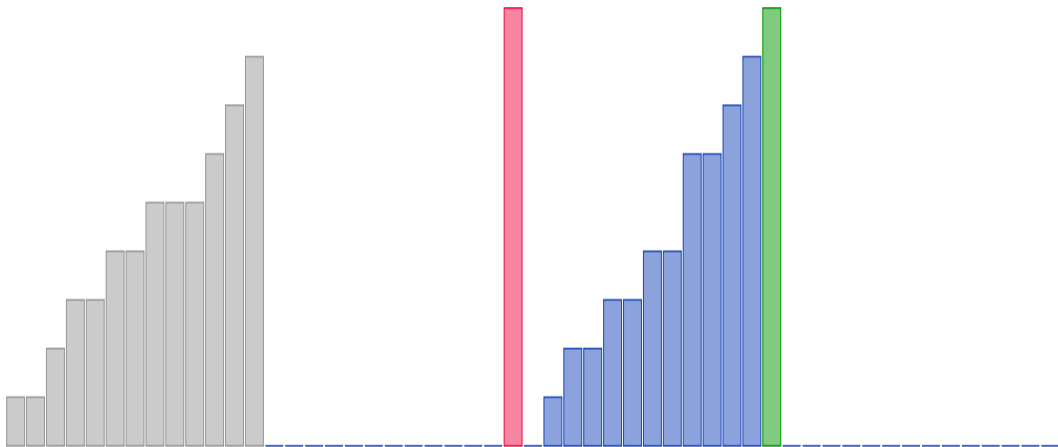


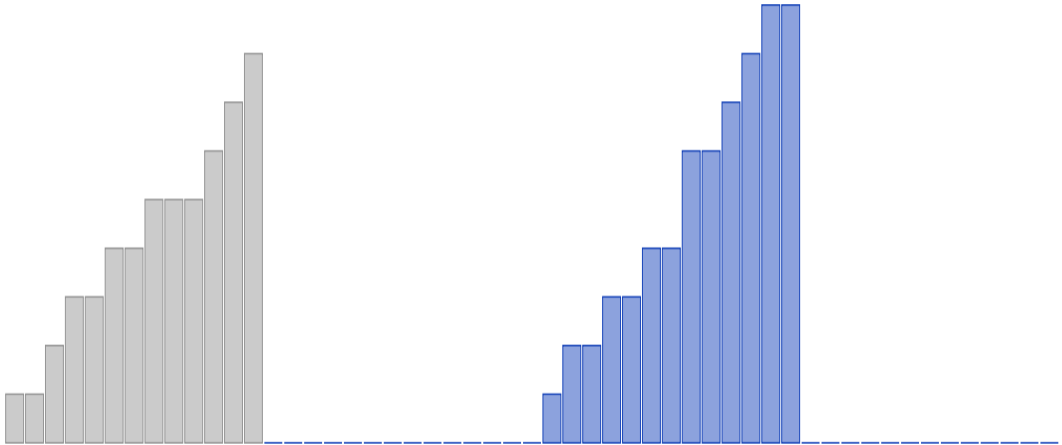


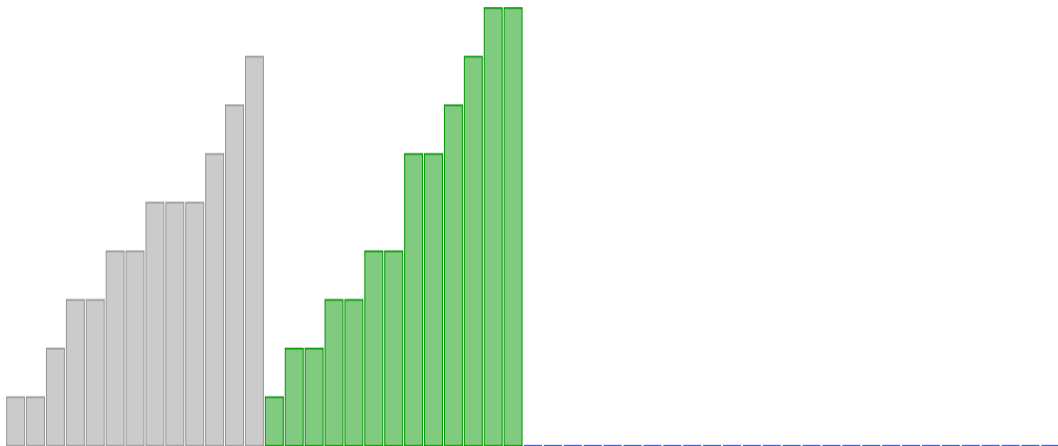


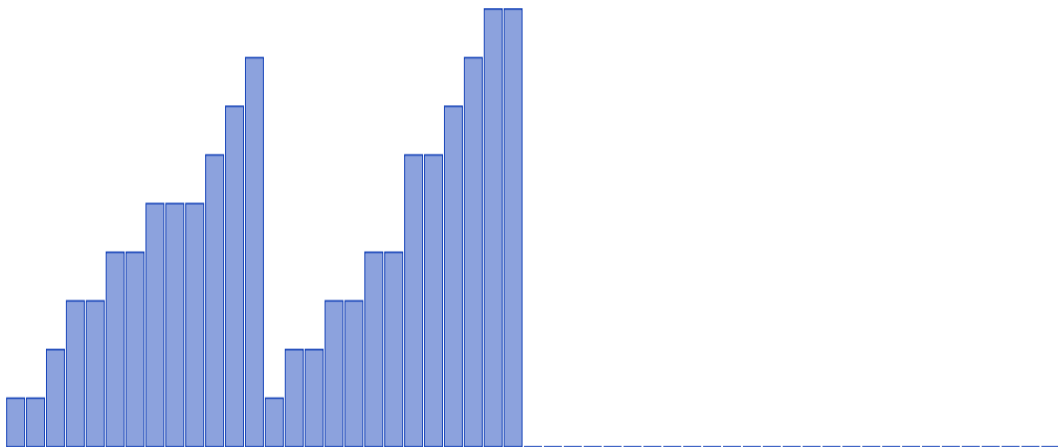


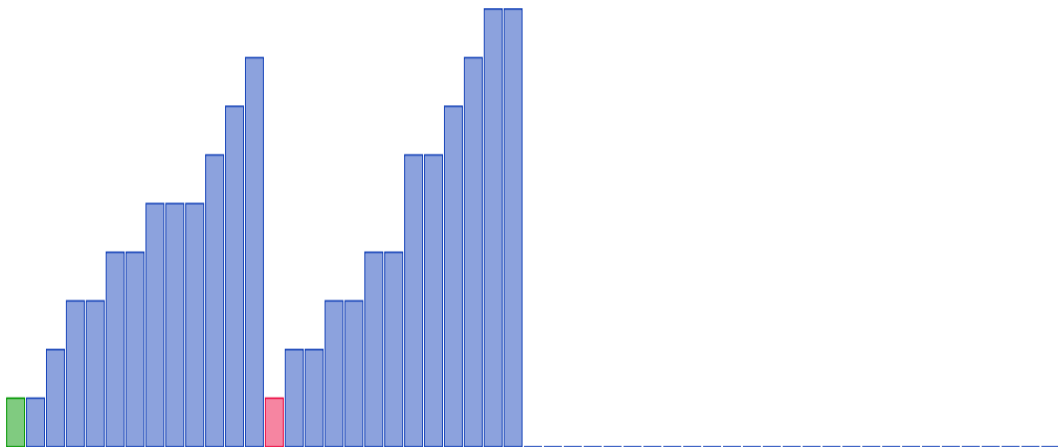


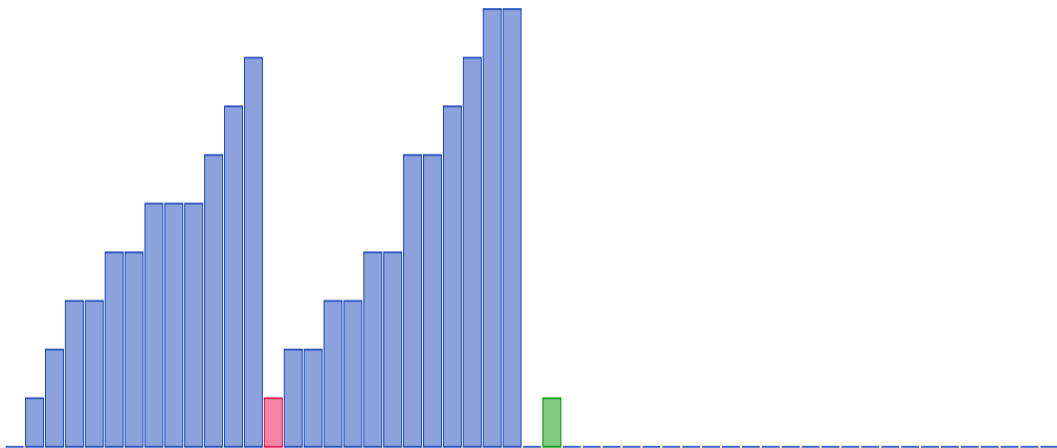


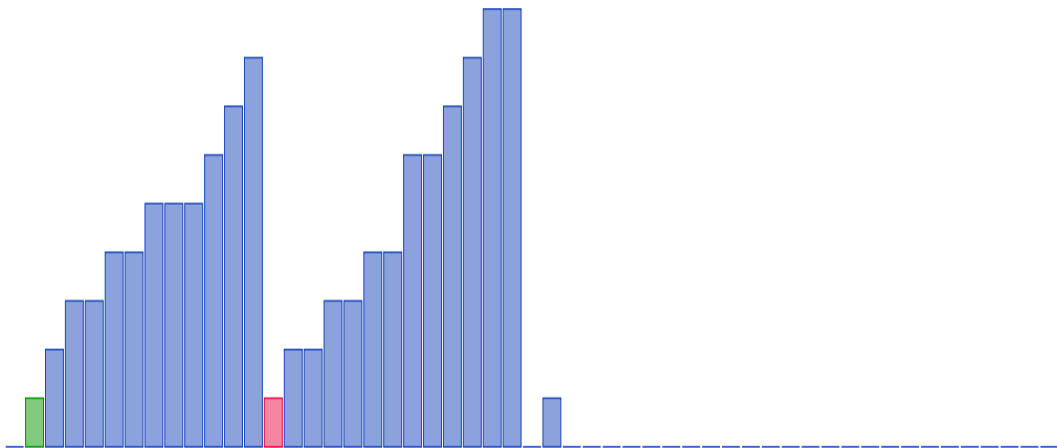


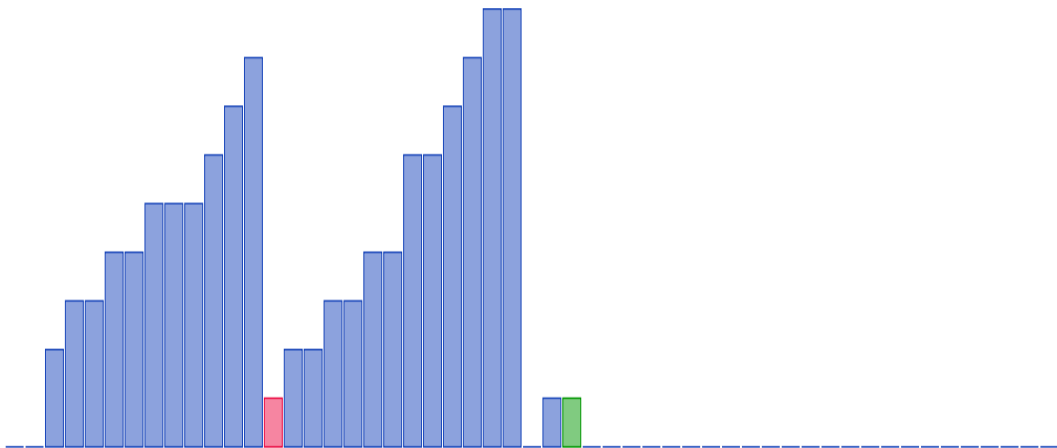


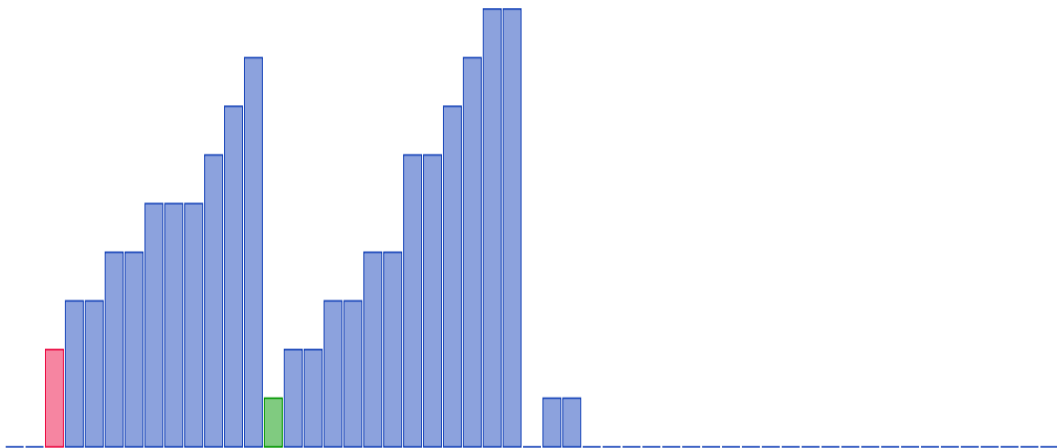


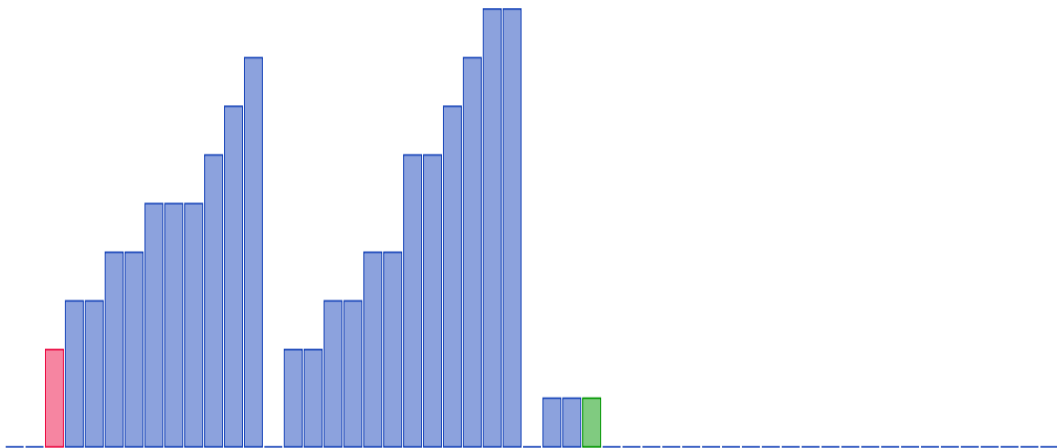


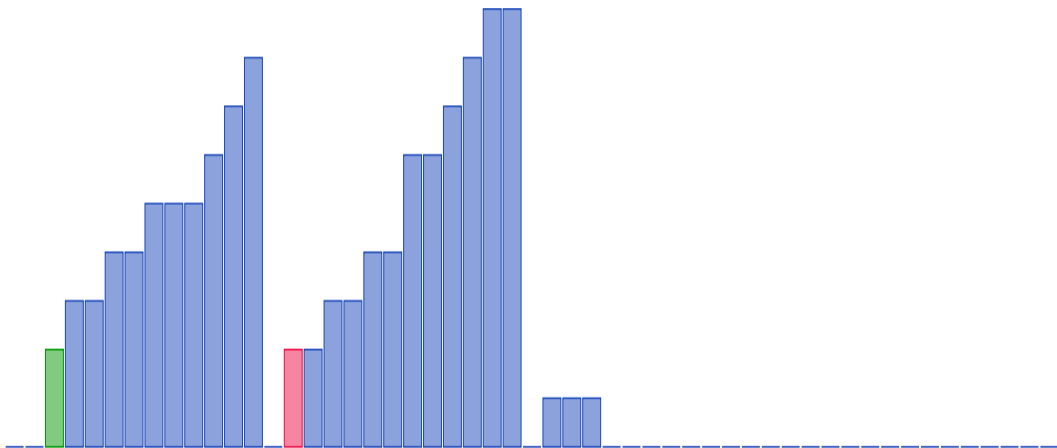


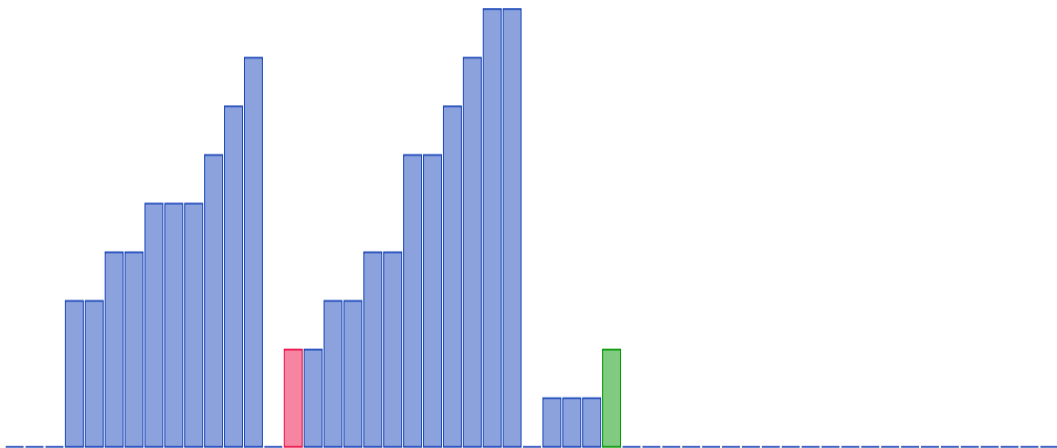


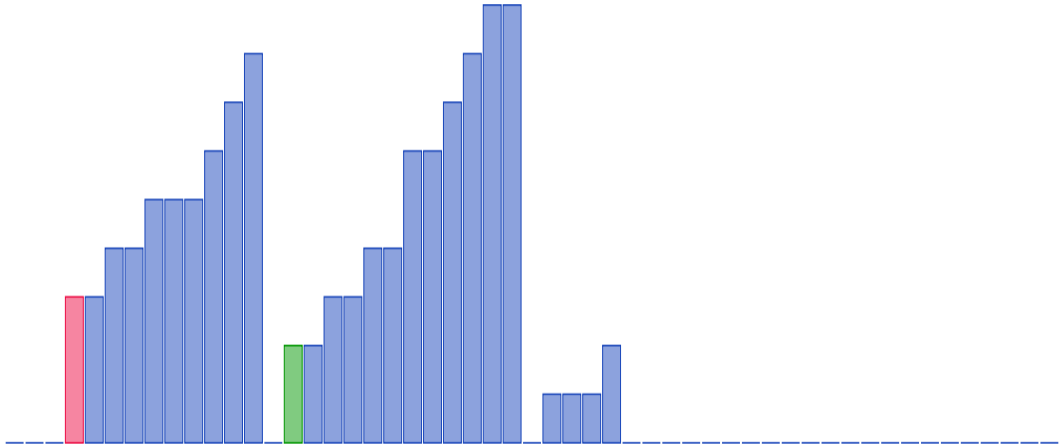


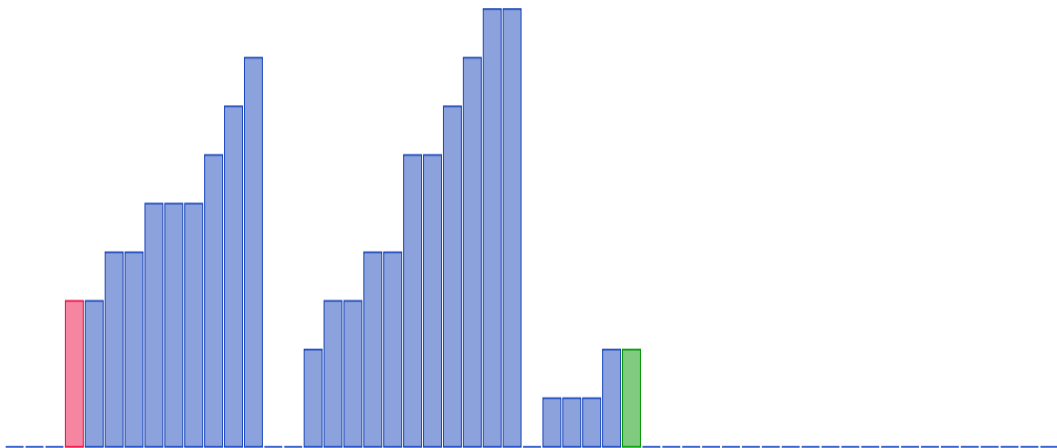


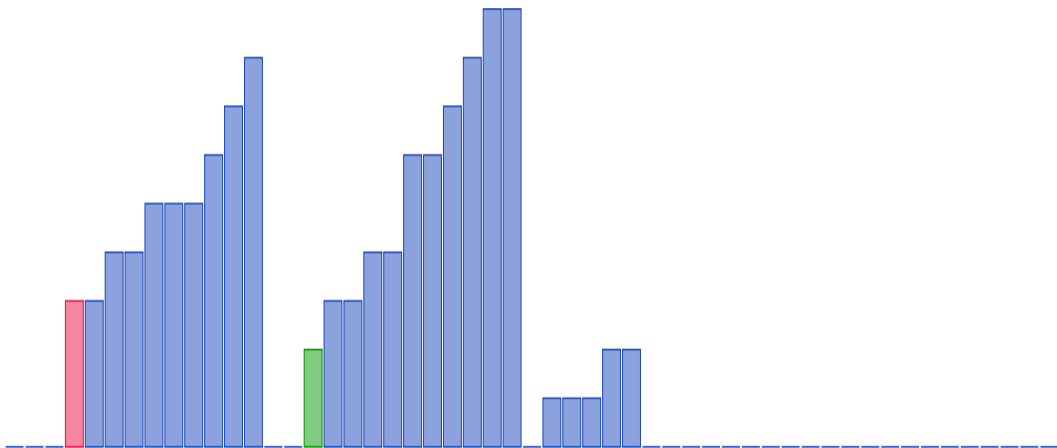


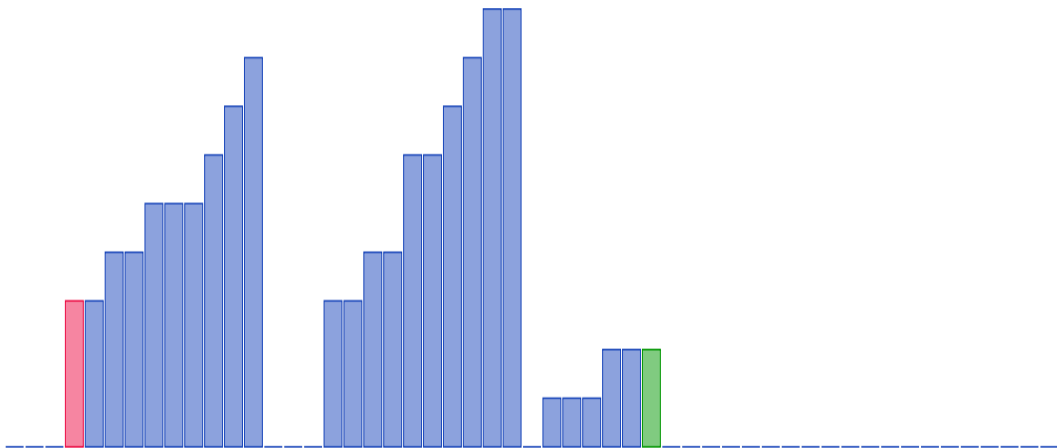


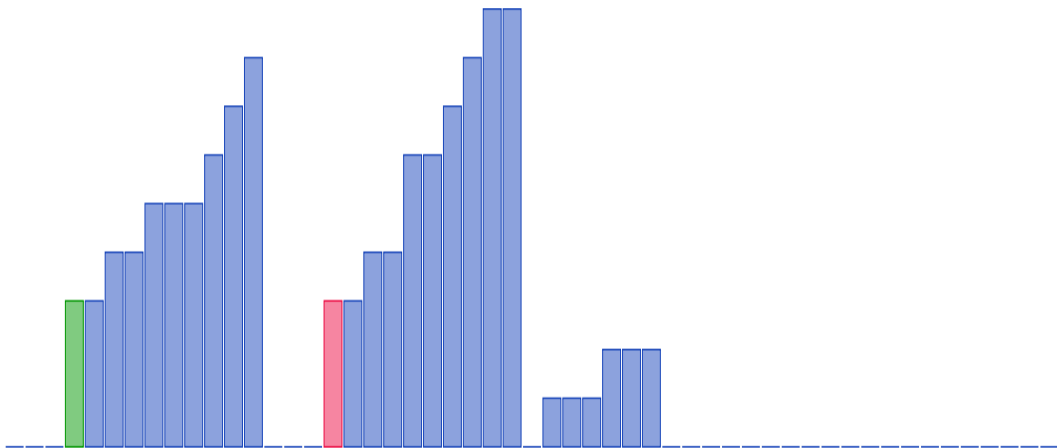


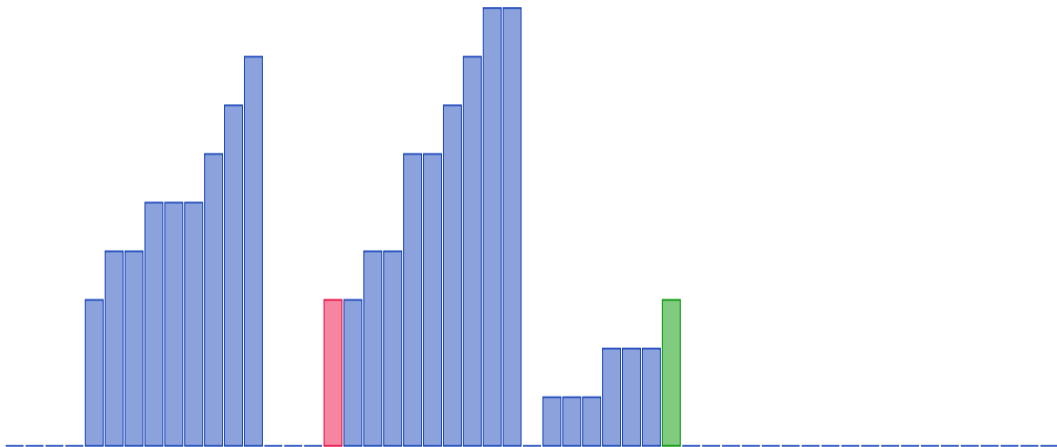


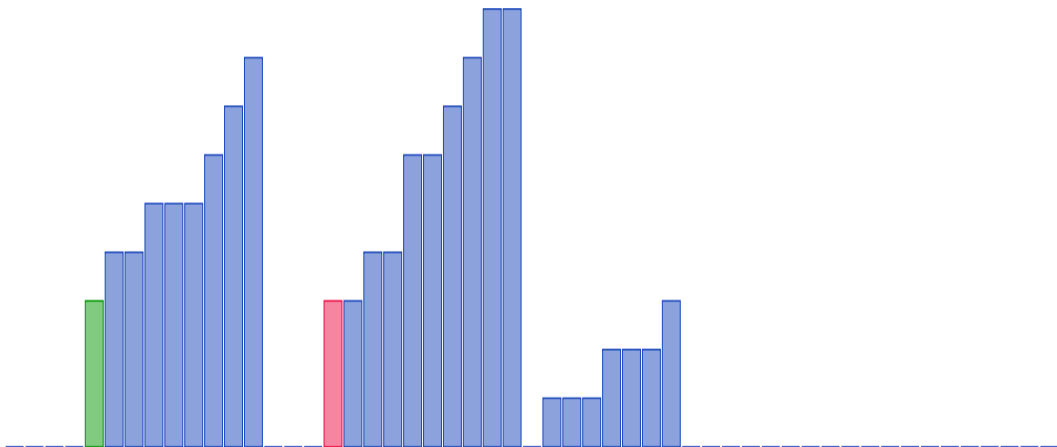


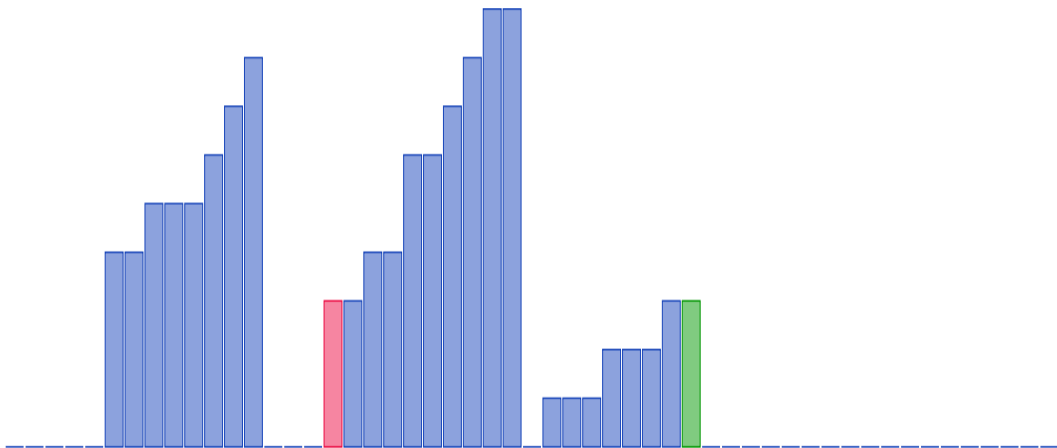


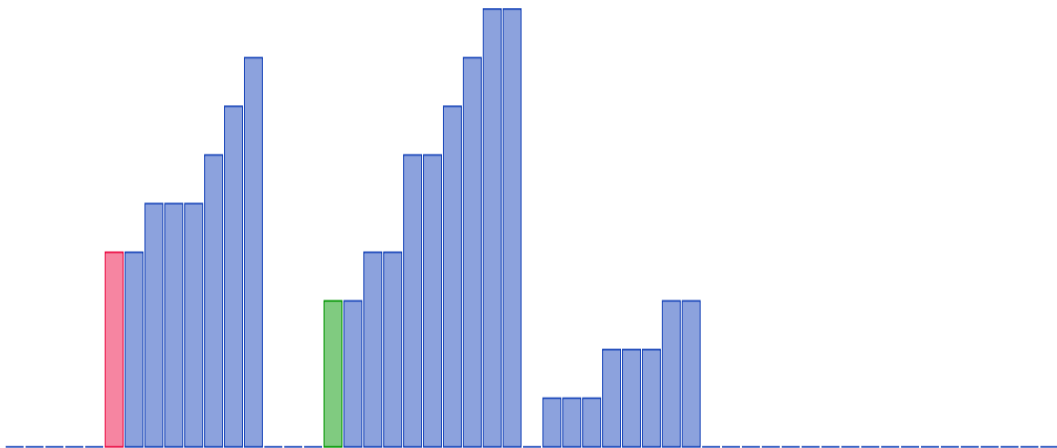


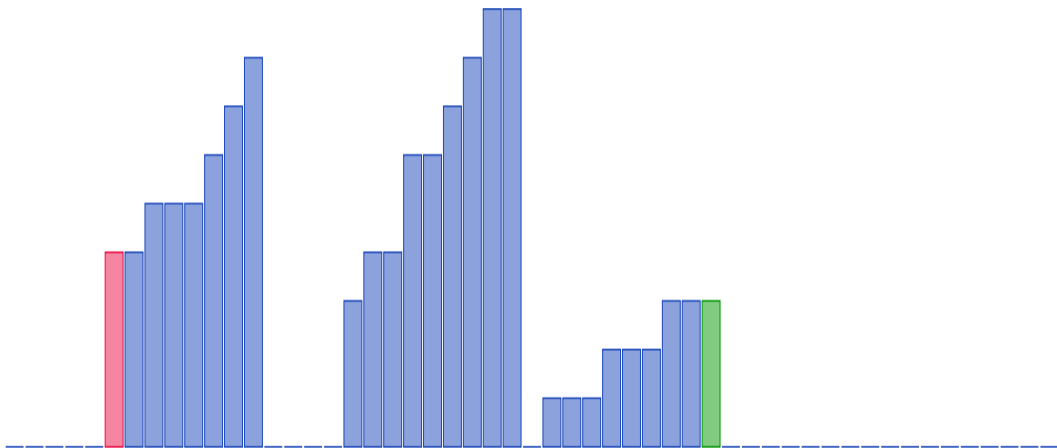


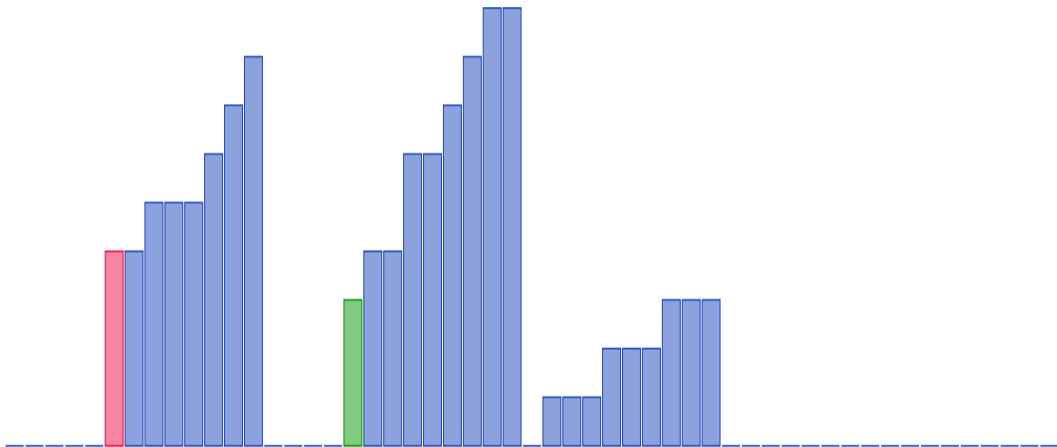


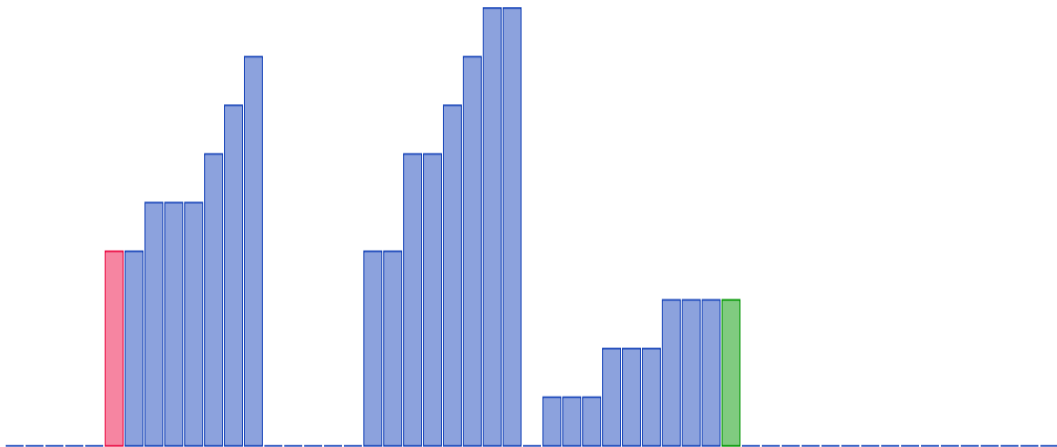


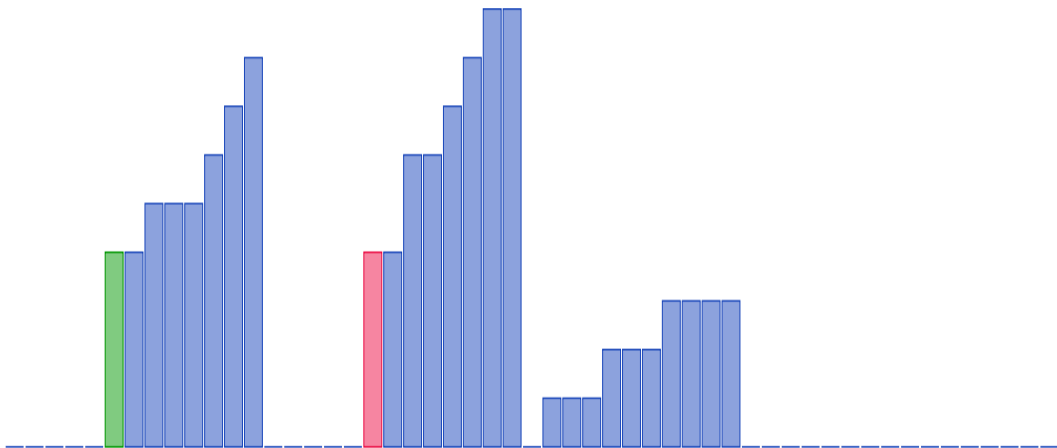


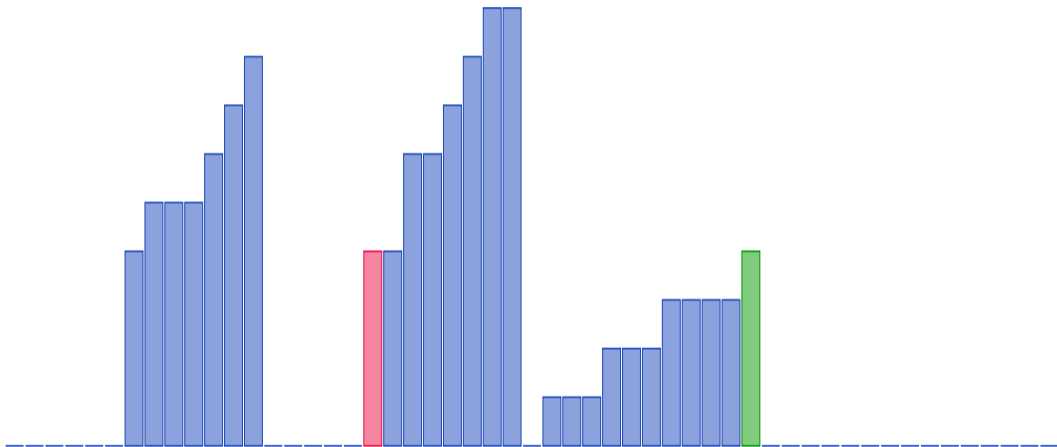


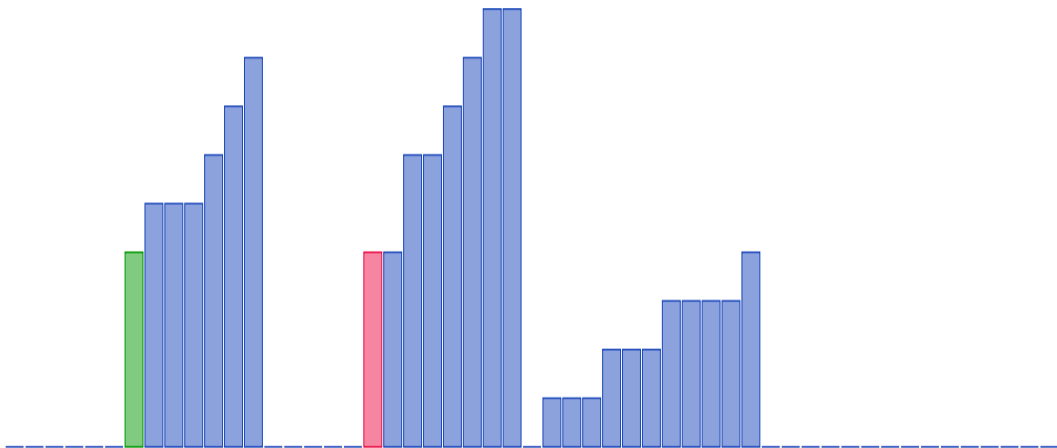


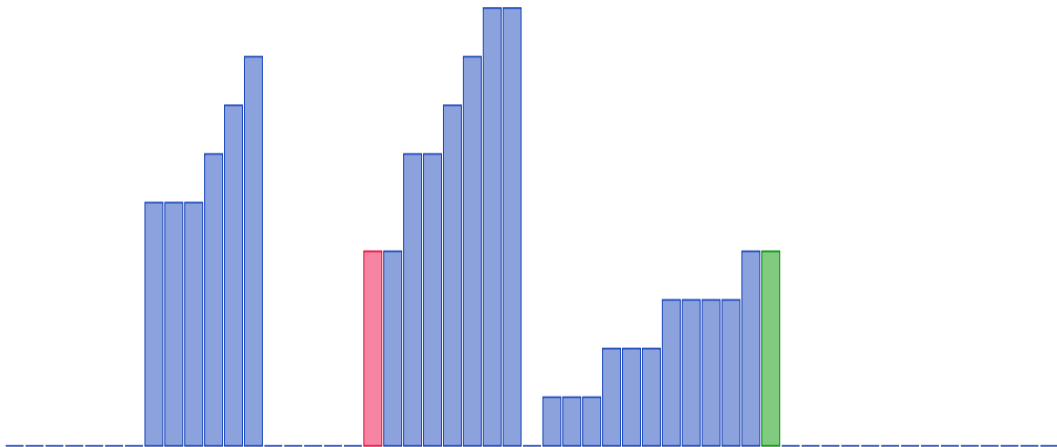


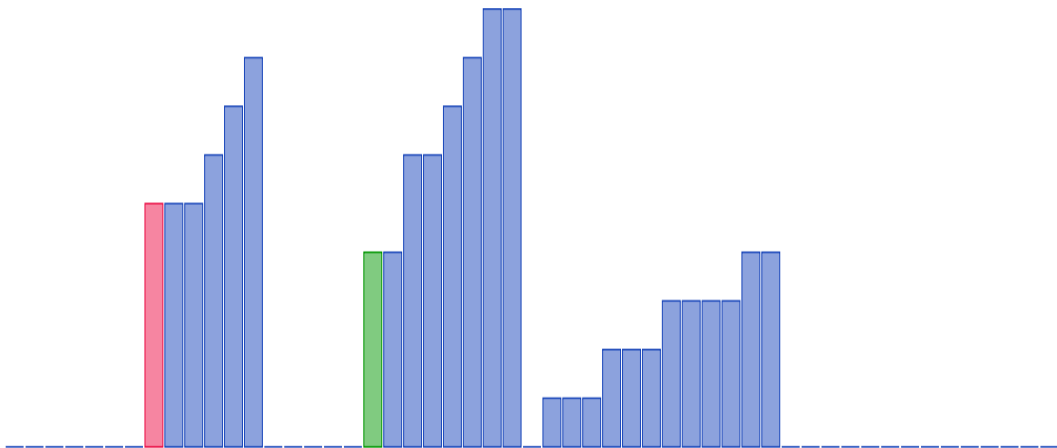


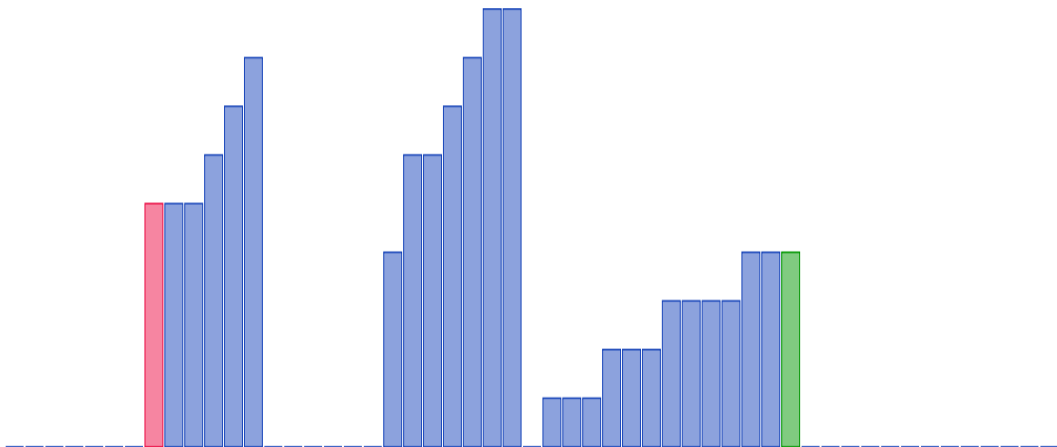


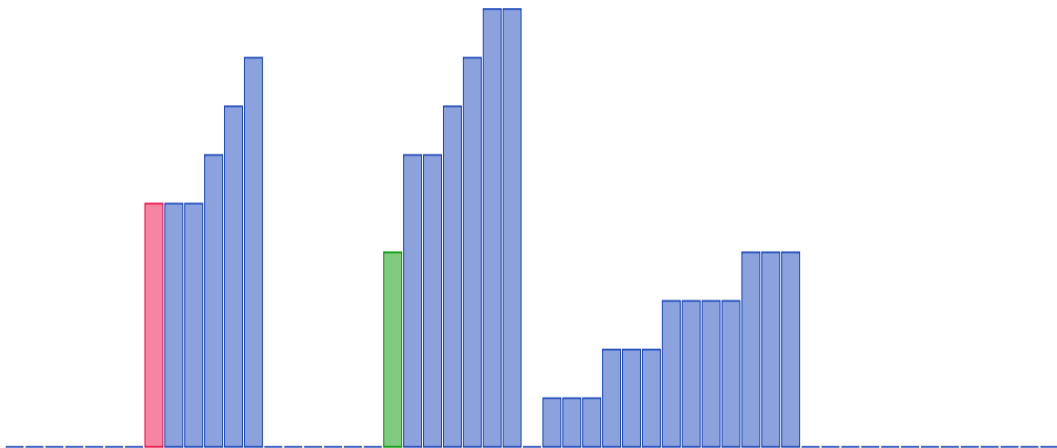


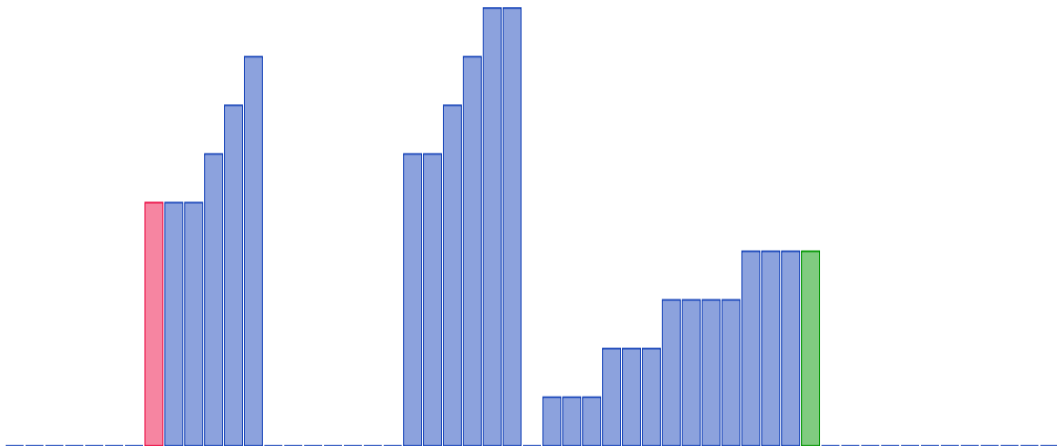


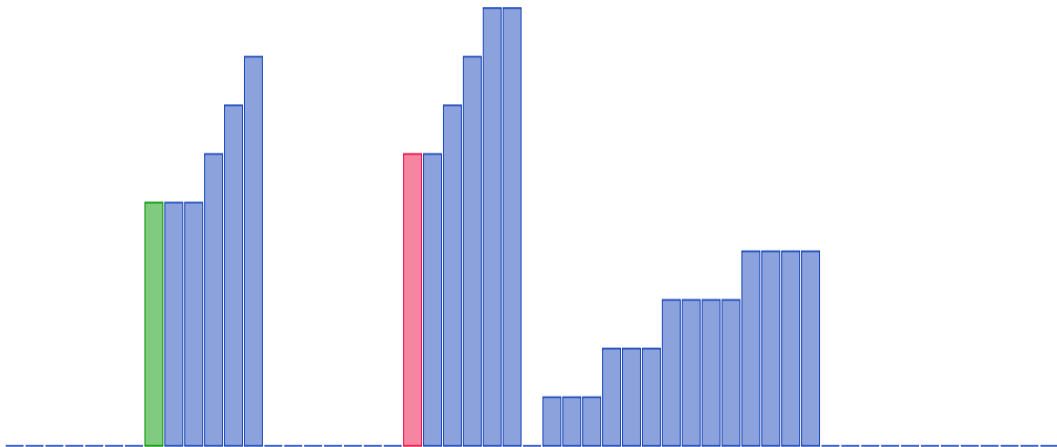


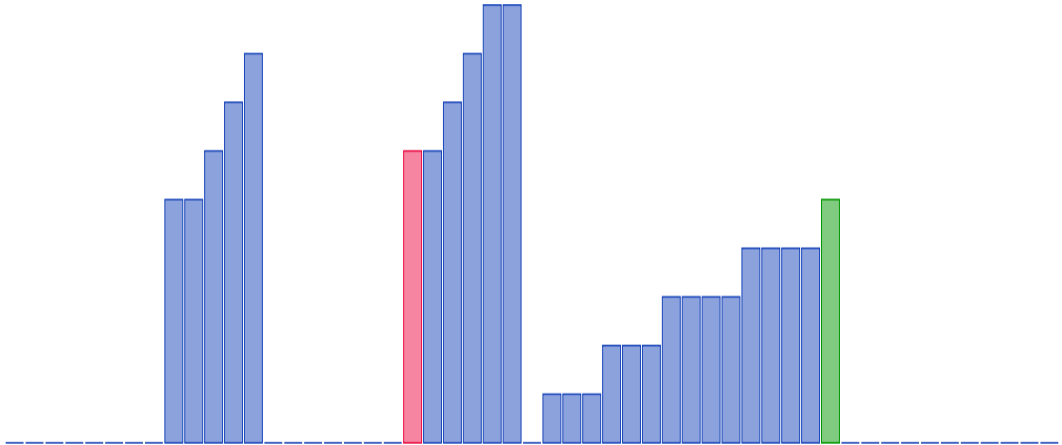


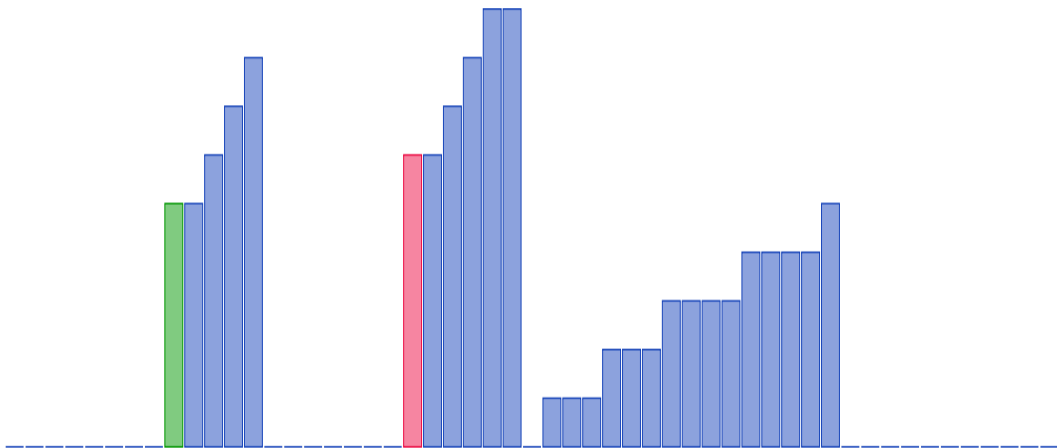


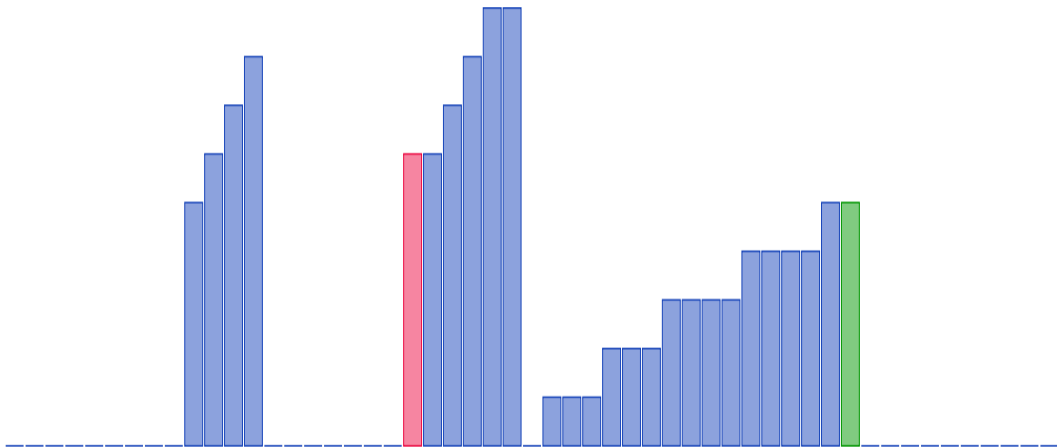


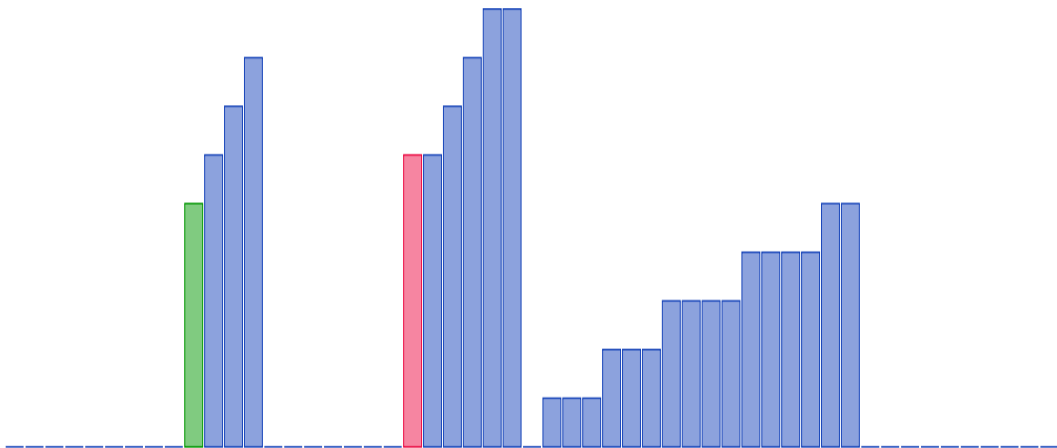


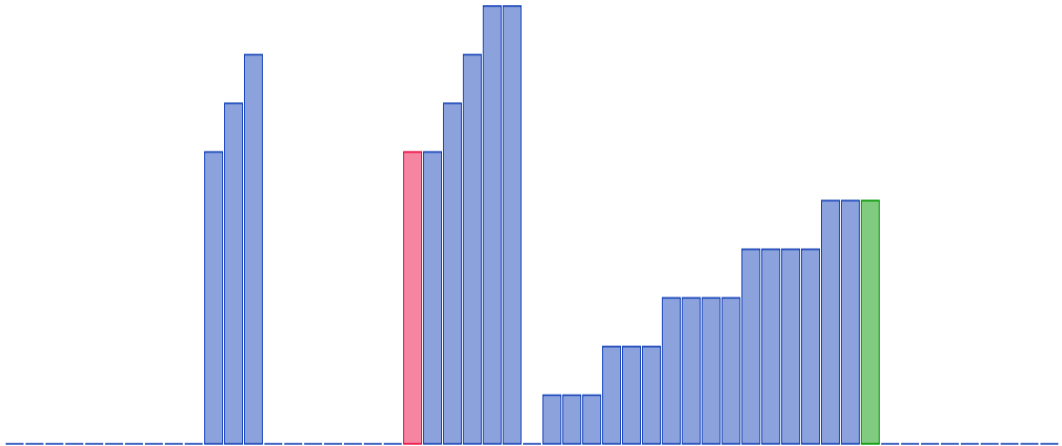


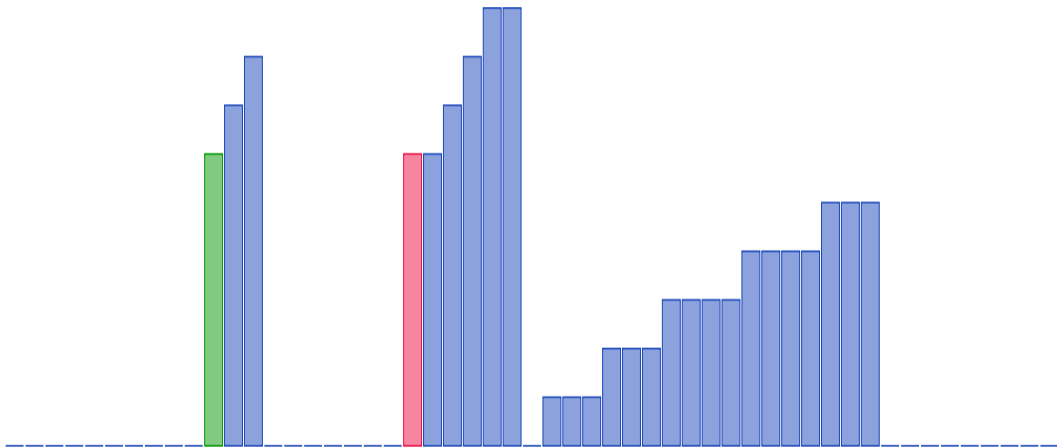


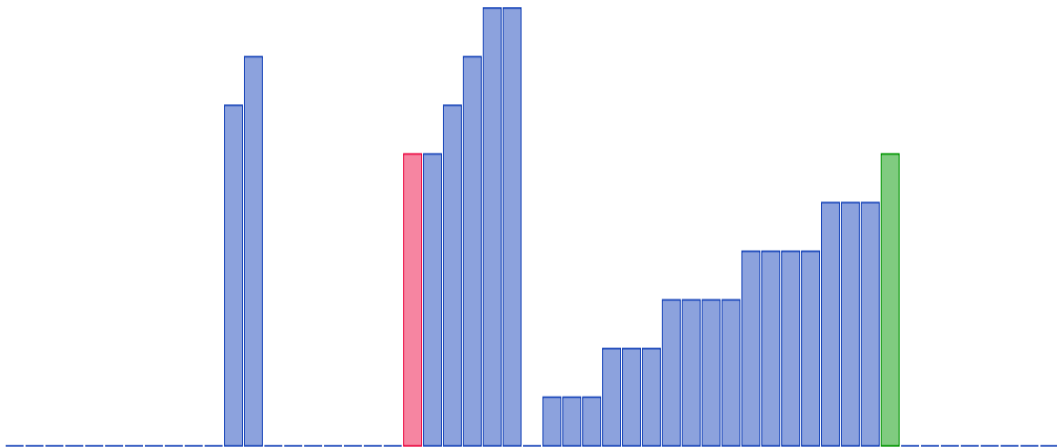


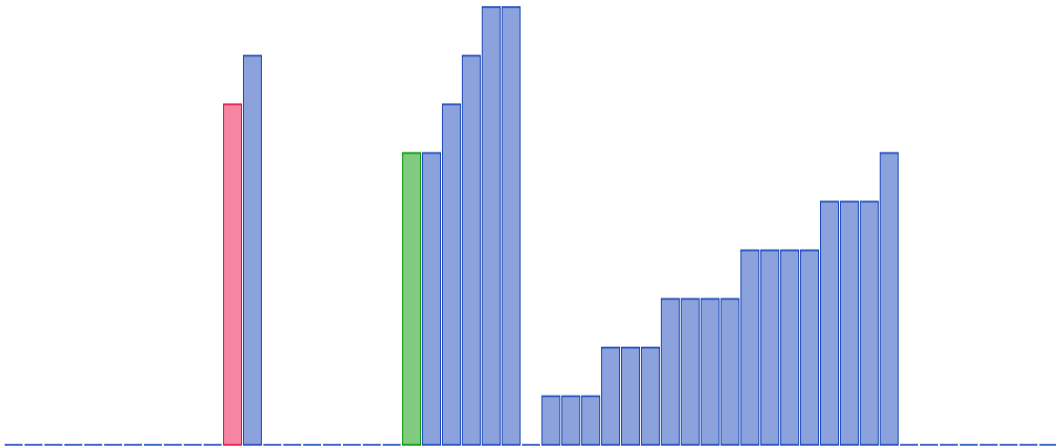


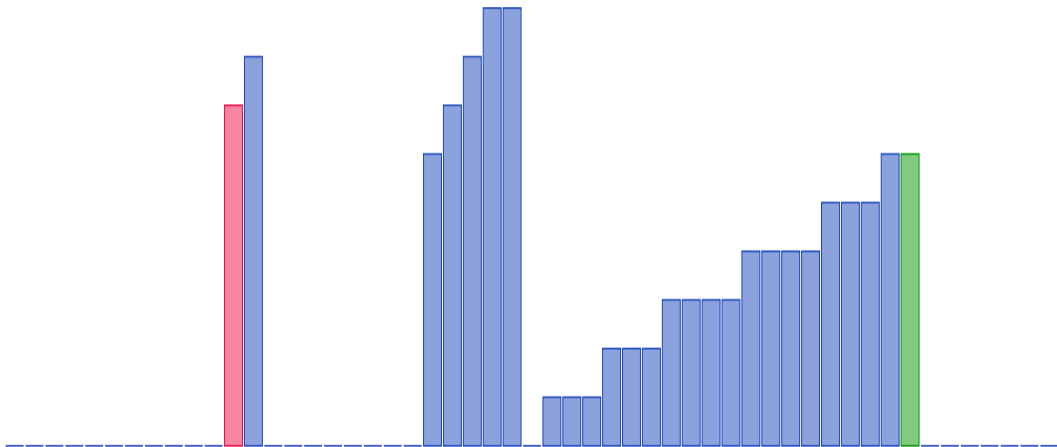


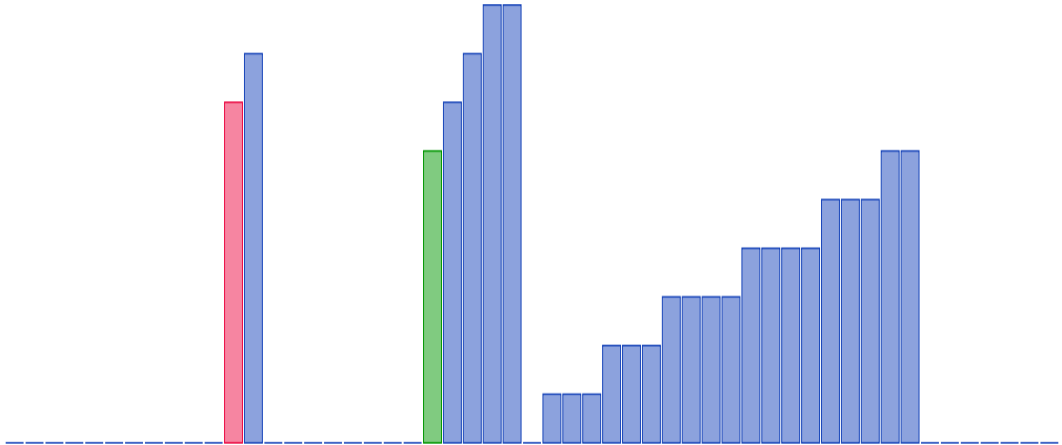


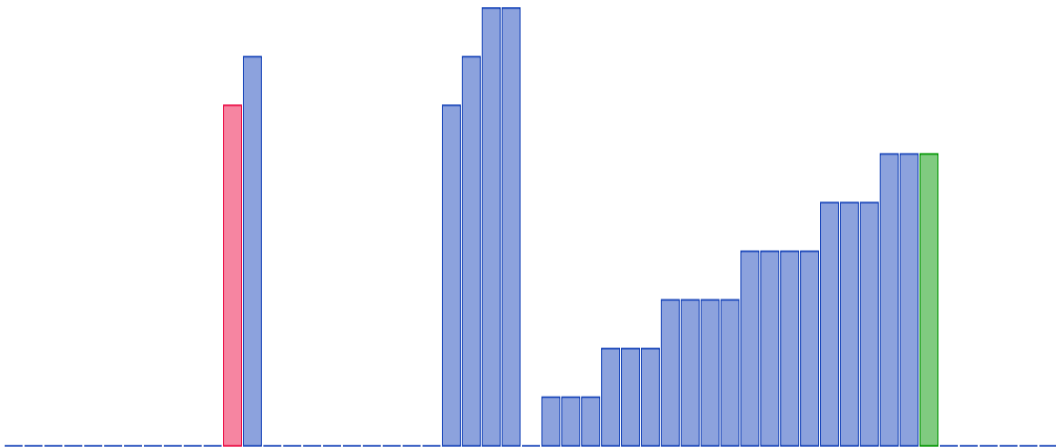


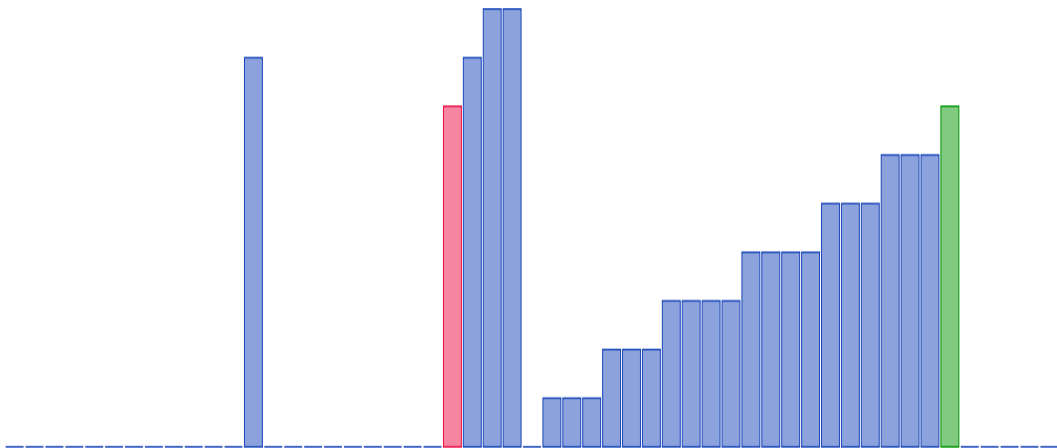


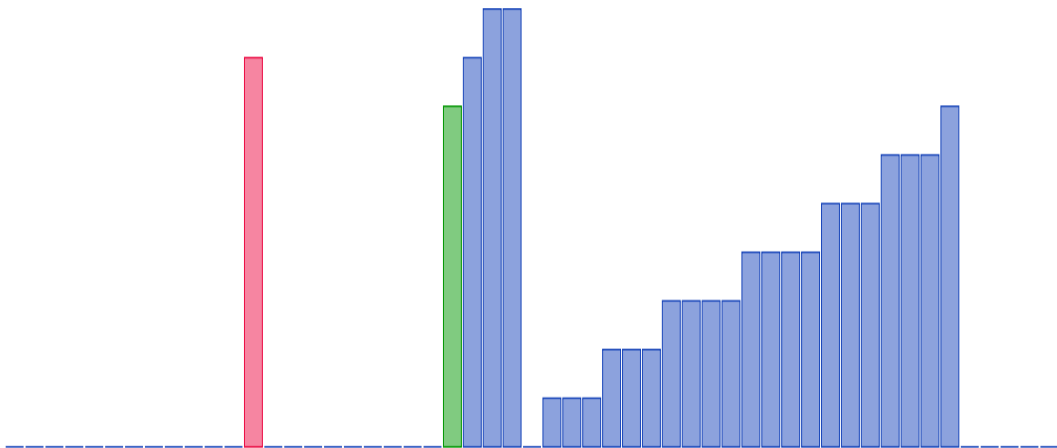


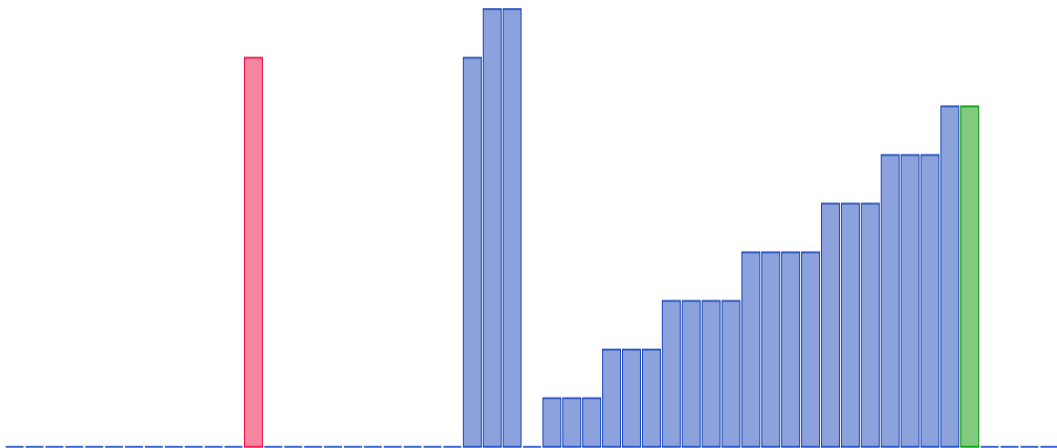


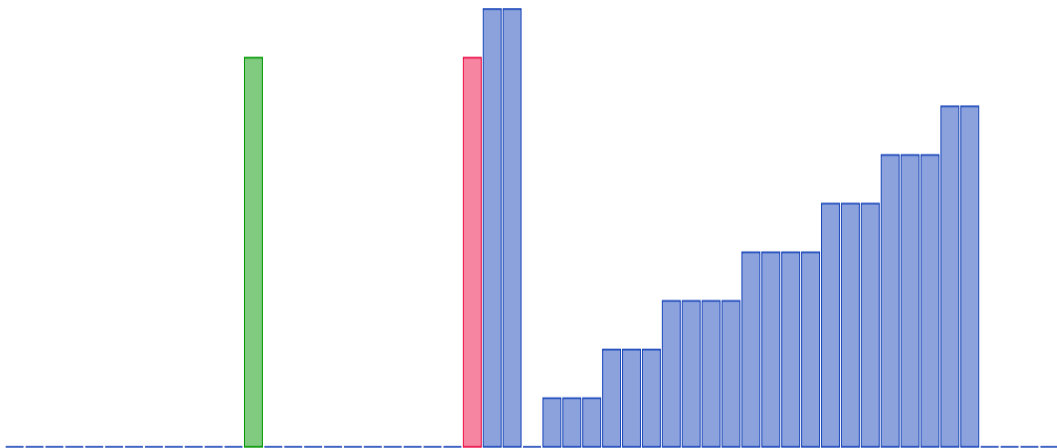


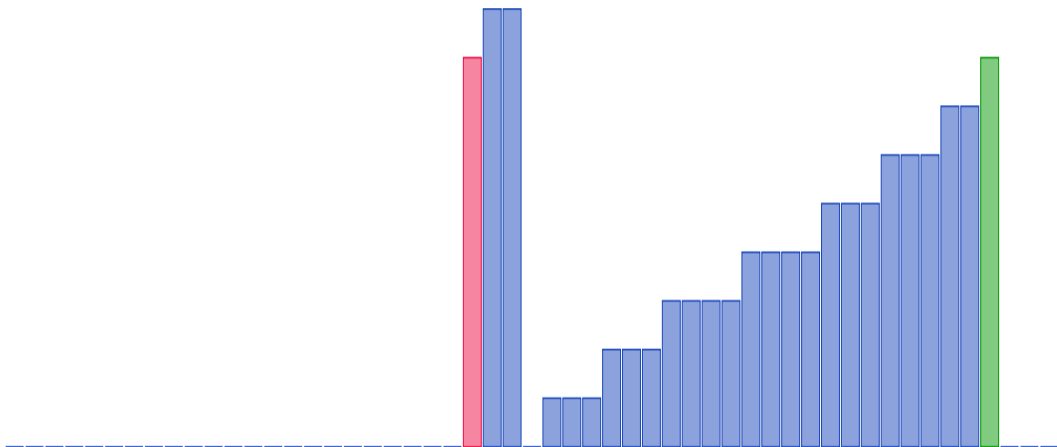


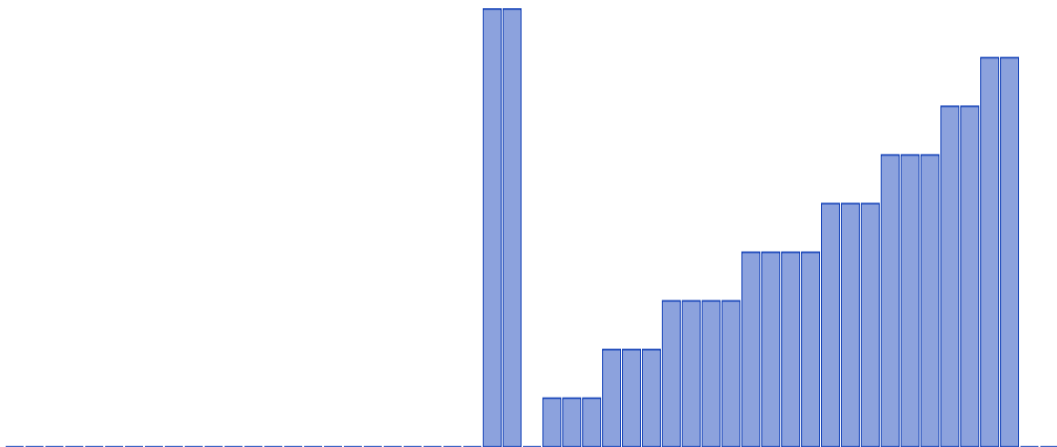


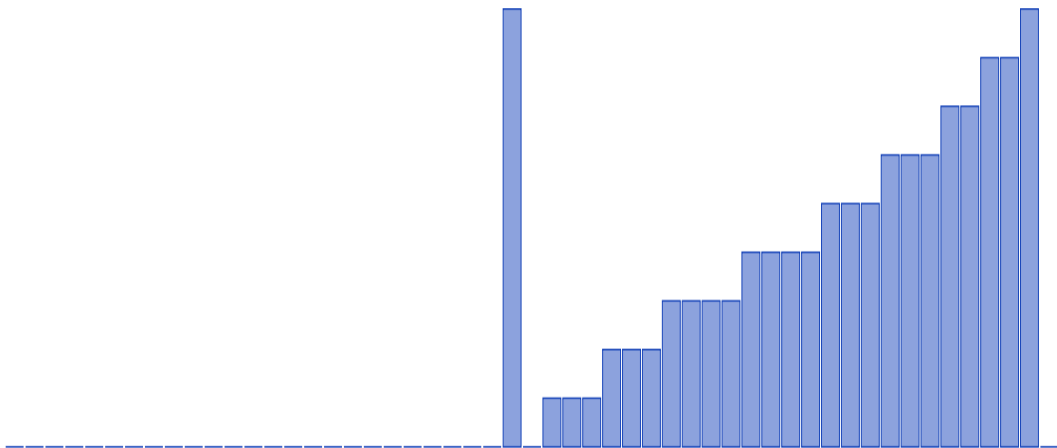


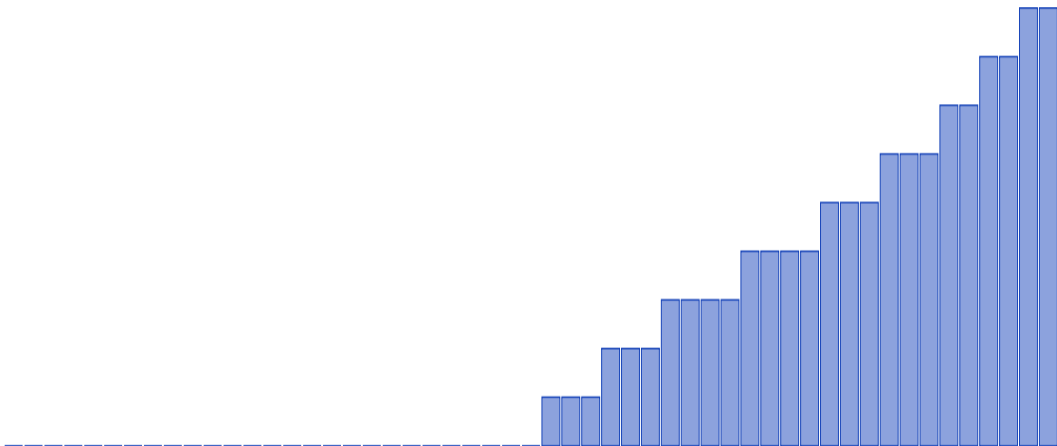


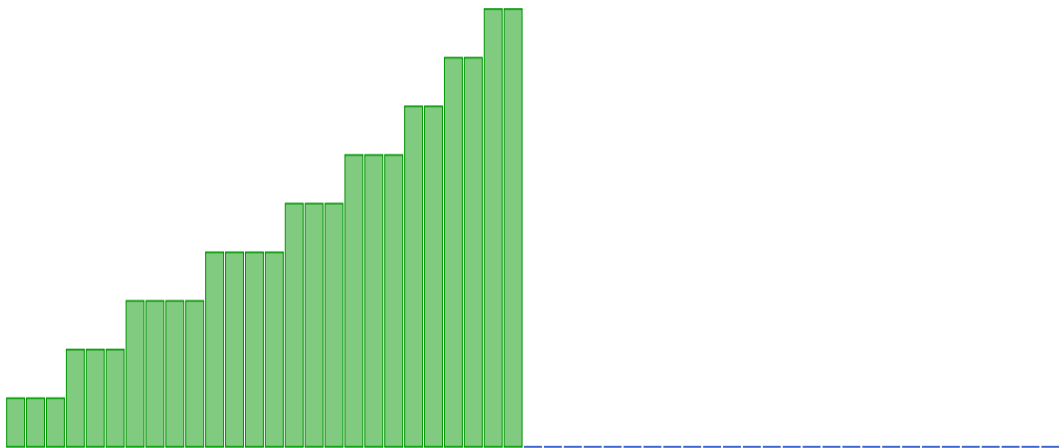












- ▶ What should the merge procedure do?
 - ▶ Given two sorted arrays, construct a sorted array from them

```
i ← s, j ← m, k ← s
while i < m or j < t do
  if j = t or (i < m and  $A[i] \leq A[j]$ ) then
     $W[k] \leftarrow A[i]$ , i ← i + 1, k ← k + 1
  else
     $W[k] \leftarrow A[j]$ , j ← j + 1, k ← k + 1
  end if
end while
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- ▶ What should the merge procedure do?
 - ▶ Given two sorted arrays, construct a sorted array from them
- ▶ Why does it do this?
 - ▶ When $A[i]$ with $s \leq i < m$ is moved?
 - ▶ After all **not greater** elements $s \leq t < i$
 - ▶ After all elements from $[m; e)$ which are **smaller** than $A[i]$

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 $i \leftarrow s, j \leftarrow m, k \leftarrow s$   
while  $i < m$  or  $j < t$  do  
  if  $j = t$  or  $(i < m$  and  $A[i] \leq A[j])$  then  
     $W[k] \leftarrow A[i], i \leftarrow i + 1, k \leftarrow k + 1$   
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 - ▶ When $A[j]$ with $m \leq j < e$ is moved?
 - ▶ After all **not greater** elements $m \leq t < j$
 - ▶ After all elements from $[s; m)$ which are **not greater** than $A[j]$

```
 $i \leftarrow s, j \leftarrow m, k \leftarrow s$   
while  $i < m$  or  $j < e$  do  
  if  $j = e$  or  $(i < m$  and  $A[i] \leq A[j])$  then  
     $W[k] \leftarrow A[i], i \leftarrow i + 1, k \leftarrow k + 1$   
  else  
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 - ▶ So, all elements are moved precisely in the sorted order

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 - ▶ When $A[j]$ with $m \leq j < e$ is moved?
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 - ▶ After all elements from $[s; m)$ which are **not greater** than $A[j]$
 - ▶ So, all elements are moved precisely in the sorted order
- ▶ The overall correctness of mergesort simply follows

```

i ← s, j ← m, k ← s
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  end if
end while

```

Running time is always $\Theta(N \log N)$.

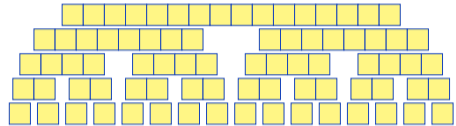
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Proof:

Running time is **always** $\Theta(N \log N)$.

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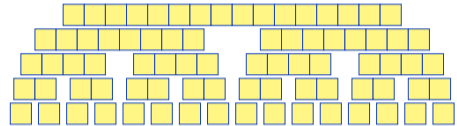
- ▶ Look at the call tree to the right



Running time is **always** $\Theta(N \log N)$.

Proof:

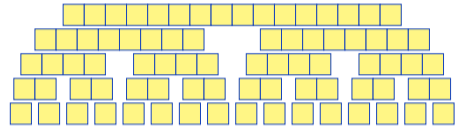
- ▶ Look at the call tree to the right
- ▶ Maximum depth: $\Theta(\log N)$, as every subarray size is **exactly half** of its parent's size



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Proof:

- ▶ Look at the call tree to the right
- ▶ Maximum depth: $\Theta(\log N)$, as every subarray size is **exactly half** of its parent's size
- ▶ All merges at given depth run in $\Theta(N)$



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Proof:

- ▶ Look at the call tree to the right
- ▶ Maximum depth: $\Theta(\log N)$, as every subarray size is **exactly half** of its parent's size
- ▶ All merges at given depth run in $\Theta(N)$
- ▶ Overall running time: $\Theta(N \log N)$

