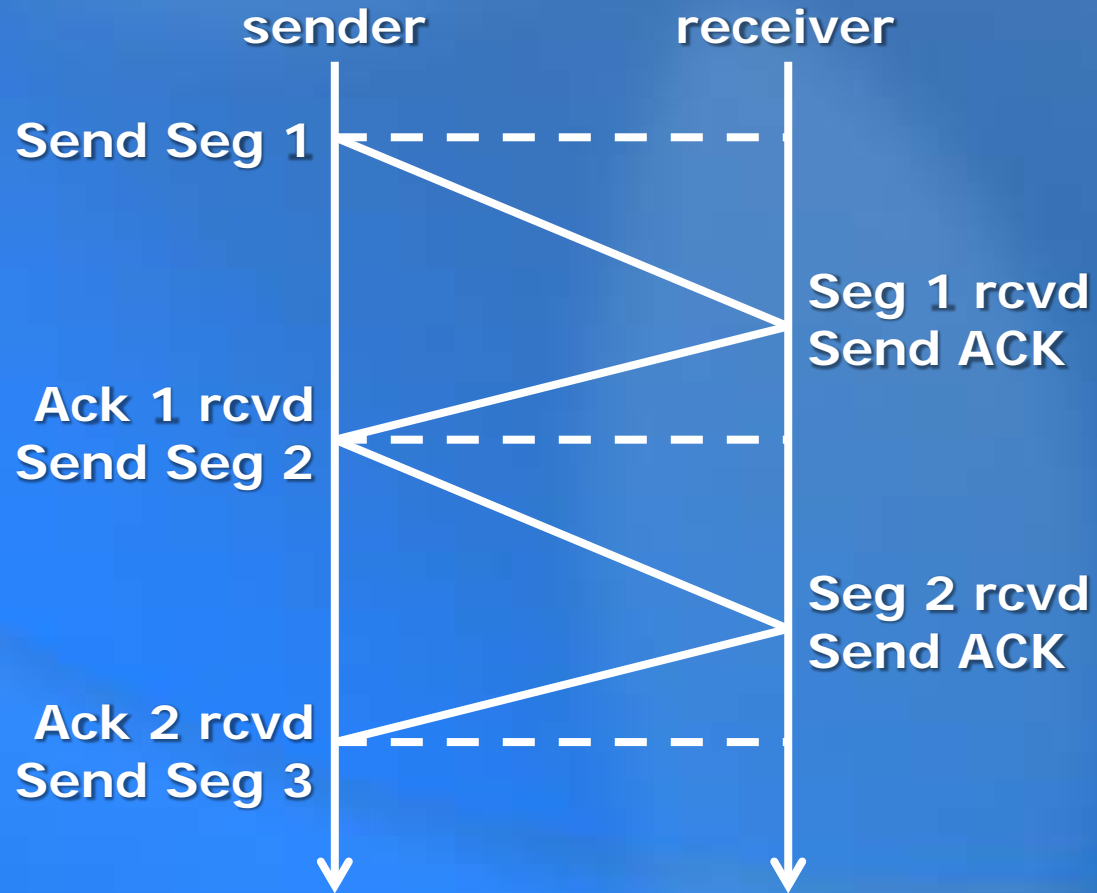


# Stop-and-wait Protocol

# Stop-and-wait (no losses)



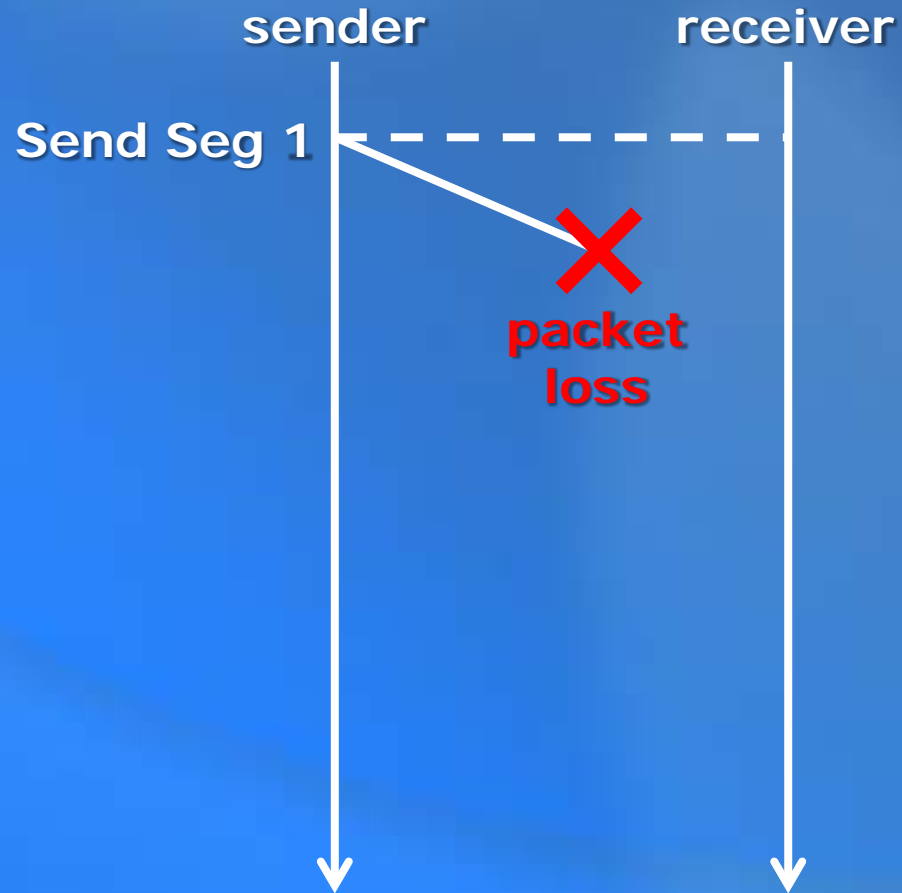
## Sender

1. Send segment  $n$
2. Wait for ACK for segment  $n$
3.  $n = n + 1$ , go to 1

## Receiver

1. Wait for segment
2. When receive, send ACK with segment number
3. Deliver packet to application
4. Go to 1

# Effect of packet loss



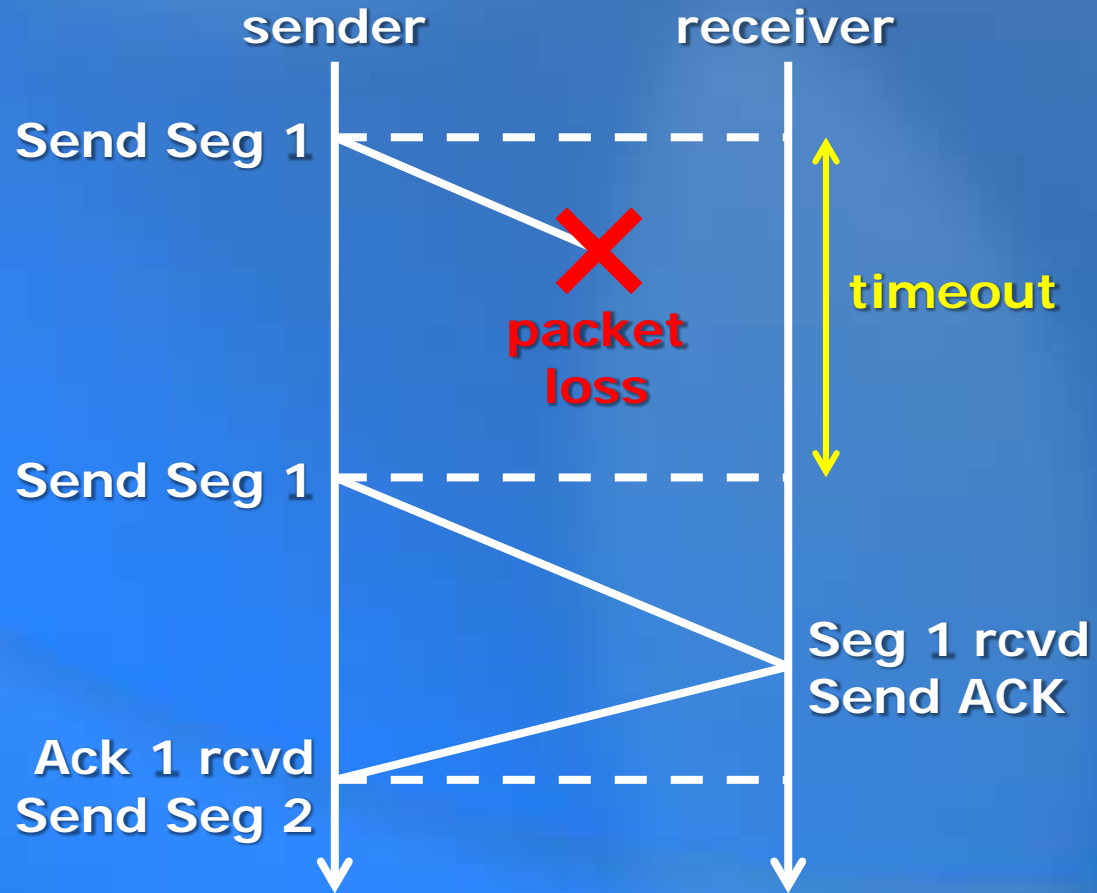
## Sender

1. Send segment  $n$
2. Wait for ACK for segment  $n$
3.  $n = n + 1$ , go to 1

## Receiver

1. Wait for segment
2. When receive, send ACK with segment number
3. Deliver packet to application
4. Go to 1

# Revised protocol



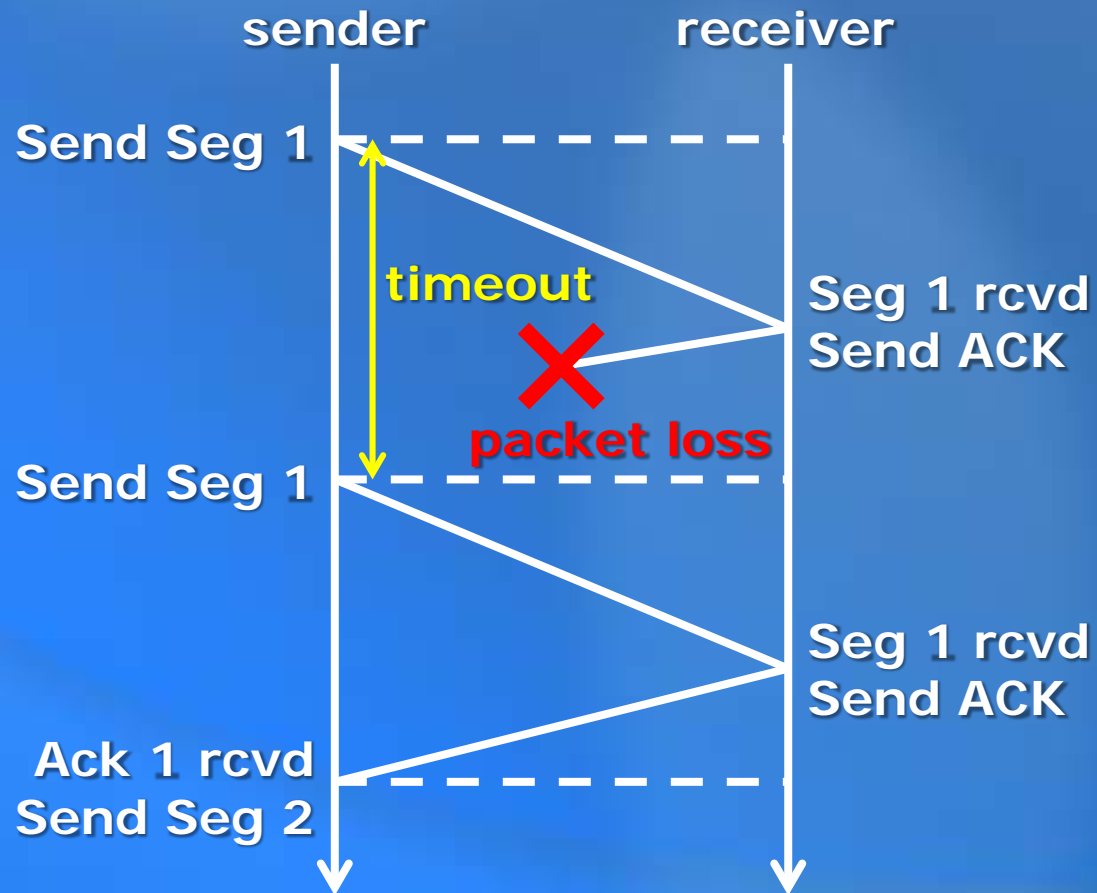
## Sender

1. Send segment  $n$
2. Wait for ACK for segment  $n$ 
  1. If no ACK received before timeout, leave  $n$  unchanged
  2. If ACK received, set  $n = n+1$
3. Go to 1

## Receiver

1. Wait for segment
2. When receive, send ACK for segment number
3. Deliver packet to application
4. Go to 1

# Effect of ACK loss



## Sender

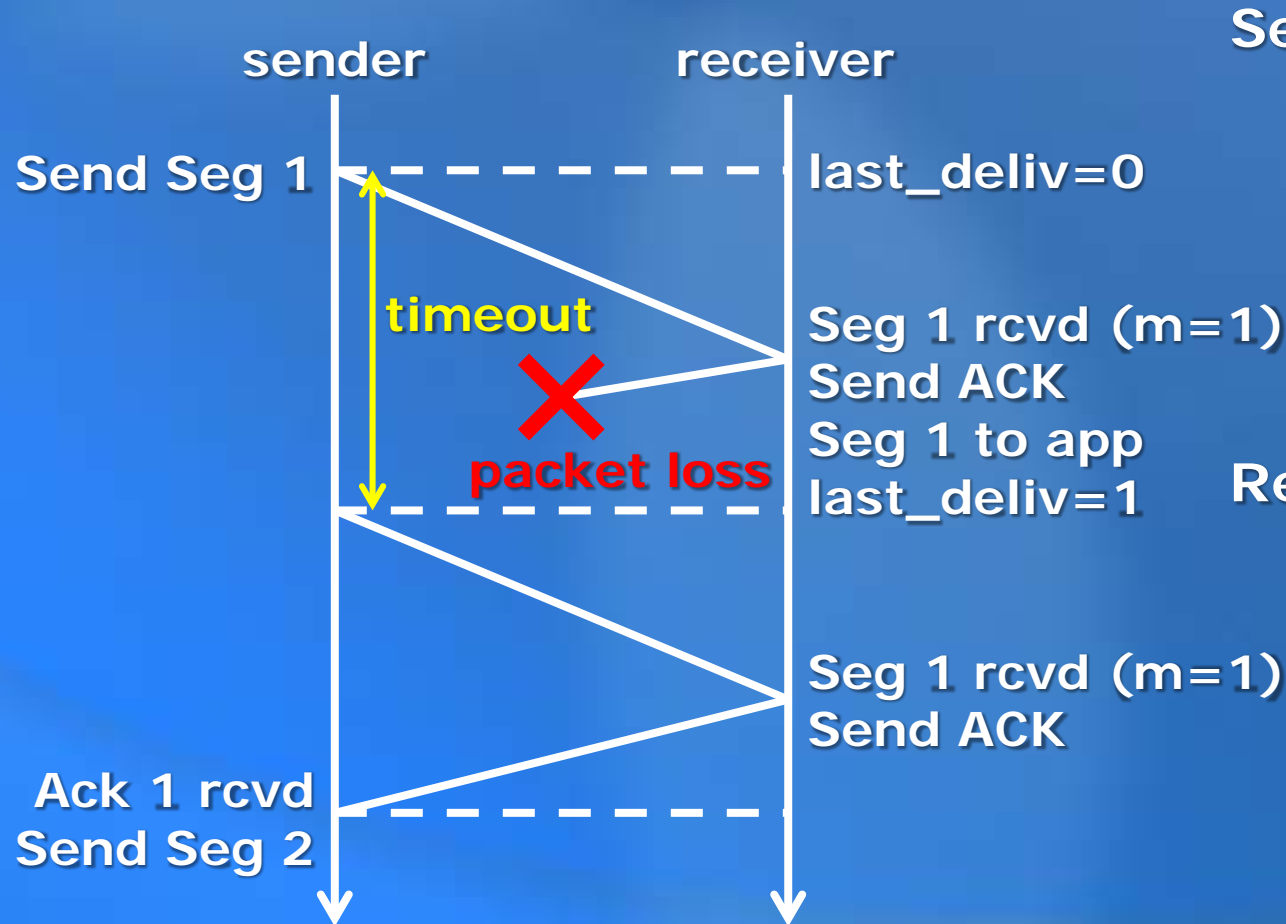
1. Send segment  $n$
2. Wait for ACK for segment  $n$ 
  1. If no ACK received before timeout, leave  $n$  unchanged
  2. If ACK received, set  $n = n + 1$
3. Go to 1

## Receiver

1. Wait for segment
2. When receive, send ACK for segment number
3. Deliver packet to application
4. Go to 1

wanted "exactly once",  
got "at least once"

# Revised Protocol



## Sender

1. Send segment  $n$
2. Wait for ACK for segment  $n$ 
  1. If no ACK received before timeout, leave  $n$  unchanged
  2. If ACK received, set  $n = n + 1$
3. Go to 1

## Receiver

1. Wait for segment
2. When receive, check segment number ( $m$ )
  1. send ACK  $m$
  2. If  $m = \text{last\_deliv} + 1$ , deliver segment to application and set  $\text{last\_deliv} = m$
  3. Otherwise, discard segment
3. Go to 1